Concrete Semantics

Tobias Nipkow & Gerwin Klein

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Abstract

This document presents formalizations of the semantics of a simple imperative programming language together with a number of applications: a compiler, type systems, various program analyses and abstract interpreters. These theories form the basis of the book *Concrete Semantics with Isabelle/HOL* by Nipkow and Klein [2].

Contents

1 Arithmetic and Boolean Expressions

1.1 Arithmetic Expressions

theory *AExp* **imports** *Main* **begin**

```
type_synonym vname = string
type_synonym val = int
type_synonym state = vname ⇒ val
```

```
datatype aexp = N int | V vname | Plus aexp aexp
```

```
fun aval :: aexp ⇒ state ⇒ val where
aval (N n) s = n |
aval (V x) s = s x |
aval (Plus a1 a2) s = aval a1 s + aval a2 s
```

```
value aval (Plus (V ''x'') (N 5)) (λx. if x = ''x'' then 7 else 0)
```

The same state more concisely:

```
value aval (Plus (V ''x'') (N 5)) ((λx. 0) (''x'':= 7))
```

A little syntax magic to write larger states compactly:

```
definition null_state (<>) where
  null_state ≡ λx. 0
syntax
  _State :: updbinds => 'a (<_>)
translations
  _State ms == _Update <> ms
  _State (_updbinds b bs) <= _Update (_State b) bs
```

We can now write a series of updates to the function $\lambda x.\ 0$ compactly:

```
lemma <a := 1, b := 2> = (<> (a := 1)) (b := (2::int))
  by (rule refl)
```

```
value aval (Plus (V ''x'') (N 5)) <''x'' := 7>
```

In the <*a* := *b*> syntax, variables that are not mentioned are 0 by default:

```
value aval (Plus (V ''x'') (N 5)) <''y'' := 7>
```

Note that this <...> syntax works for any function space $\tau_1 \Rightarrow \tau_2$ where τ_2 has a *0*.

1.2 Constant Folding

Evaluate constant subsexpressions:

```
fun asimp_const :: aexp ⇒ aexp where
asimp_const (N n) = N n |
asimp_const (V x) = V x |
asimp_const (Plus a₁ a₂) =
  (case (asimp_const a₁, asimp_const a₂) of
    (N n₁, N n₂) ⇒ N(n₁+n₂) |
    (b₁,b₂) ⇒ Plus b₁ b₂)

theorem aval_asimp_const:
  aval (asimp_const a) s = aval a s
apply(induction a)
apply (auto split: aexp.split)
done
```

Now we also eliminate all occurrences 0 in additions. The standard method: optimized versions of the constructors:

```
fun plus :: aexp ⇒ aexp ⇒ aexp where
plus (N i₁) (N i₂) = N(i₁+i₂) |
plus (N i) a = (if i=0 then a else Plus (N i) a) |
plus a (N i) = (if i=0 then a else Plus a (N i)) |
plus a₁ a₂ = Plus a₁ a₂

lemma aval_plus[simp]:
  aval (plus a1 a2) s = aval a1 s + aval a2 s
apply(induction a1 a2 rule: plus.induct)
apply simp_all
done

fun asimp :: aexp ⇒ aexp where
asimp (N n) = N n |
asimp (V x) = V x |
asimp (Plus a₁ a₂) = plus (asimp a₁) (asimp a₂)
```

Note that in *asimp_const* the optimized constructor was inlined. Making it a separate function *AExp.plus* improves modularity of the code and the proofs.

```
value asimp (Plus (Plus (N 0) (N 0)) (Plus (V ''x'') (N 0)))

theorem aval_asimp[simp]:
  aval (asimp a) s = aval a s
apply(induction a)
```

apply *simp_all*
done

end

1.3 Boolean Expressions

theory *BExp* **imports** *AExp* **begin**

datatype *bexp* = *Bc bool* | *Not bexp* | *And bexp bexp* | *Less aexp aexp*
fun *bval* :: *bexp* ⇒ *state* ⇒ *bool* **where**
bval (*Bc v*) *s* = *v* |
bval (*Not b*) *s* = (¬ *bval b s*) |
bval (*And* b_1 b_2) *s* = (*bval* b_1 *s* ∧ *bval* b_2 *s*) |
bval (*Less* a_1 a_2) *s* = (*aval* a_1 *s* < *aval* a_2 *s*)

value *bval* (*Less* (*V ''x''*) (*Plus* (*N 3*) (*V ''y''*)))
<''x'' := 3, ''y'' := 1>

1.4 Constant Folding

Optimizing constructors:

fun *less* :: *aexp* ⇒ *aexp* ⇒ *bexp* **where**
less (*N* n_1) (*N* n_2) = *Bc*(n_1 < n_2) |
less a_1 a_2 = *Less* a_1 a_2

lemma [*simp*]: *bval* (*less a1 a2*) *s* = (*aval a1 s* < *aval a2 s*)
apply(*induction a1 a2 rule*: *less.induct*)
apply *simp_all*
done

fun *and* :: *bexp* ⇒ *bexp* ⇒ *bexp* **where**
and (*Bc True*) *b* = *b* |
and b (*Bc True*) = *b* |
and (*Bc False*) *b* = *Bc False* |
and b (*Bc False*) = *Bc False* |
and b_1 b_2 = *And* b_1 b_2

lemma *bval_and*[*simp*]: *bval* (*and b1 b2*) *s* = (*bval b1 s* ∧ *bval b2 s*)
apply(*induction b1 b2 rule*: *and.induct*)
apply *simp_all*
done

fun *not* :: *bexp* ⇒ *bexp* **where**

not (*Bc True*) = *Bc False* |
not (*Bc False*) = *Bc True* |
not b = *Not b*

lemma *bval_not*[*simp*]: *bval* (*not b*) *s* = (¬ *bval b s*)
apply(*induction b rule*: *not.induct*)
apply *simp_all*
done

Now the overall optimizer:

fun *bsimp* :: *bexp* ⇒ *bexp* **where**
bsimp (*Bc v*) = *Bc v* |
bsimp (*Not b*) = *not*(*bsimp b*) |
bsimp (*And* b_1 b_2) = *and* (*bsimp* b_1) (*bsimp* b_2) |
bsimp (*Less* a_1 a_2) = *less* (*asimp* a_1) (*asimp* a_2)

value *bsimp* (*And* (*Less* (*N 0*) (*N 1*)) *b*)

value *bsimp* (*And* (*Less* (*N 1*) (*N 0*)) (*Bc True*))

theorem *bval* (*bsimp b*) *s* = *bval b s*
apply(*induction b*)
apply *simp_all*
done

end

2 Stack Machine and Compilation

theory *ASM* **imports** *AExp* **begin**

2.1 Stack Machine

datatype *instr* = *LOADI val* | *LOAD vname* | *ADD*

type_synonym *stack* = *val list*

Abbreviations are transparent: they are unfolded after parsing and folded back again before printing. Internally, they do not exist.

fun *exec1* :: *instr* ⇒ *state* ⇒ *stack* ⇒ *stack* **where**
exec1 (*LOADI n*) _ *stk* = *n* # *stk* |
exec1 (*LOAD x*) *s stk* = *s*(*x*) # *stk* |

exec1 ADD _ (j # i # stk) = (i + j) # stk

fun *exec* :: *instr list* ⇒ *state* ⇒ *stack* ⇒ *stack* **where**
exec [] _ *stk* = *stk* |
exec (*i#is*) *s stk* = *exec is s* (*exec1 i s stk*)

value *exec* [*LOADI 5*, *LOAD ''y''*, *ADD*] <''x'' := 42, ''y'' := 43> [50]

lemma *exec_append*[*simp*]:
  exec (*is1@is2*) *s stk* = *exec is2 s* (*exec is1 s stk*)
apply(*induction is1 arbitrary*: *stk*)
apply (*auto*)
done

2.2 Compilation

fun *comp* :: *aexp* ⇒ *instr list* **where**
comp (*N n*) = [*LOADI n*] |
comp (*V x*) = [*LOAD x*] |
comp (*Plus* e_1 e_2) = *comp* e_1 @ *comp* e_2 @ [*ADD*]

value *comp* (*Plus* (*Plus* (*V ''x''*) (*N 1*)) (*V ''z''*))

theorem *exec_comp*: *exec* (*comp a*) *s stk* = *aval a s* # *stk*
apply(*induction a arbitrary*: *stk*)
apply (*auto*)
done

end
theory *Star* **imports** *Main*
begin

inductive
  star :: (*'a* ⇒ *'a* ⇒ *bool*) ⇒ *'a* ⇒ *'a* ⇒ *bool*
for *r* **where**
refl: *star r x x* |
step: *r x y* ⟹ *star r y z* ⟹ *star r x z*

hide_fact (**open**) *refl step* — names too generic

lemma *star_trans*:
  star r x y ⟹ *star r y z* ⟹ *star r x z*
proof(*induction rule*: *star.induct*)
  case *refl* **thus** *?case* **.**

next
  case *step* **thus** *?case* **by** (*metis star.step*)
qed

lemmas *star_induct* =
  star.induct[*of r*:: *'a**'b* ⇒ *'a**'b* ⇒ *bool*, *split_format*(*complete*)]

declare *star.refl*[*simp*,*intro*]

lemma *star_step1*[*simp*, *intro*]: *r x y* ⟹ *star r x y*
by(*metis star.refl star.step*)

code_pred *star* **.**

end

3 IMP — A Simple Imperative Language

theory *Com* **imports** *BExp* **begin**

datatype
```
  com = SKIP
      | Assign vname aexp      (_ ::= _ [1000, 61] 61)
      | Seq    com  com        (_;;/ _  [60, 61] 60)
      | If     bexp com com    ((IF _/ THEN _/ ELSE _)  [0, 0, 61] 61)
      | While  bexp com        ((WHILE _/ DO _)  [0, 61] 61)
```

end

3.1 Big-Step Semantics of Commands

theory *Big_Step* **imports** *Com* **begin**

The big-step semantics is a straight-forward inductive definition with concrete syntax. Note that the first parameter is a tuple, so the syntax becomes $(c,s) \Rightarrow s'$.

inductive
  big_step :: *com* × *state* ⇒ *state* ⇒ *bool* (**infix** ⇒ *55*)
where
Skip: $(SKIP,s) \Rightarrow s$ |
Assign: $(x ::= a,s) \Rightarrow s(x := aval\ a\ s)$ |
Seq: $[\![(c_1,s_1) \Rightarrow s_2;\ (c_2,s_2) \Rightarrow s_3]\!] \Longrightarrow (c_1;;c_2,\ s_1) \Rightarrow s_3$ |
IfTrue: $[\![bval\ b\ s;\ (c_1,s) \Rightarrow t]\!] \Longrightarrow (IF\ b\ THEN\ c_1\ ELSE\ c_2,\ s) \Rightarrow t$ |
IfFalse: $[\![\neg bval\ b\ s;\ (c_2,s) \Rightarrow t]\!] \Longrightarrow (IF\ b\ THEN\ c_1\ ELSE\ c_2,\ s) \Rightarrow t$ |

WhileFalse: $\neg bval\ b\ s \Longrightarrow (WHILE\ b\ DO\ c,s) \Rightarrow s$ |
WhileTrue:
$[\![\ bval\ b\ s_1;\ \ (c,s_1) \Rightarrow s_2;\ \ (WHILE\ b\ DO\ c,\ s_2) \Rightarrow s_3\ ]\!]$
$\Longrightarrow (WHILE\ b\ DO\ c,\ s_1) \Rightarrow s_3$

schematic_goal *ex*: $(''x'' ::= N\ 5;;\ ''y'' ::= V\ ''x'',\ s) \Rightarrow ?t$
apply(*rule Seq*)
apply(*rule Assign*)
apply *simp*
apply(*rule Assign*)
done

thm *ex*[*simplified*]

We want to execute the big-step rules:

code_pred *big_step* **.**

For inductive definitions we need command `values` instead of `value`.

values $\{t.\ (SKIP,\ \lambda\_.\ 0) \Rightarrow t\}$

We need to translate the result state into a list to display it.

values $\{map\ t\ [''x'']\ |t.\ (SKIP,\ <''x'' := 42>) \Rightarrow t\}$

values $\{map\ t\ [''x'']\ |t.\ (''x'' ::= N\ 2,\ <''x'' := 42>) \Rightarrow t\}$

values $\{map\ t\ [''x'',''y'']\ |t.$
$(WHILE\ Less\ (V\ ''x'')\ (V\ ''y'')\ DO\ (''x'' ::= Plus\ (V\ ''x'')\ (N\ 5)),$
$<''x'' := 0,\ ''y'' := 13>) \Rightarrow t\}$

Proof automation:

The introduction rules are good for automatically construction small program executions. The recursive cases may require backtracking, so we declare the set as unsafe intro rules.

declare *big_step.intros* [*intro*]

The standard induction rule

$[\![x1 \Rightarrow x2;\ \bigwedge s.\ P\ (SKIP,\ s)\ s;\ \bigwedge x\ a\ s.\ P\ (x ::= a,\ s)\ (s(x := aval\ a\ s));$
$\bigwedge c_1\ s_1\ s_2\ c_2\ s_3.$
$[\![(c_1,\ s_1) \Rightarrow s_2;\ P\ (c_1,\ s_1)\ s_2;\ (c_2,\ s_2) \Rightarrow s_3;\ P\ (c_2,\ s_2)\ s_3]\!]$
$\Longrightarrow P\ (c_1;;\ c_2,\ s_1)\ s_3;$
$\bigwedge b\ s\ c_1\ t\ c_2.$
$[\![bval\ b\ s;\ (c_1,\ s) \Rightarrow t;\ P\ (c_1,\ s)\ t]\!] \Longrightarrow P\ (IF\ b\ THEN\ c_1\ ELSE\ c_2,\ s)\ t;$
$\bigwedge b\ s\ c_2\ t\ c_1.$

⟦¬ *bval b s*; (*c*$_2$, *s*) ⇒ *t*; *P* (*c*$_2$, *s*) *t*⟧ ⟹ *P* (*IF b THEN c*$_1$ *ELSE c*$_2$, *s*) *t*;
⋀*b s c*. ¬ *bval b s* ⟹ *P* (*WHILE b DO c*, *s*) *s*;
⋀*b s*$_1$ *c s*$_2$ *s*$_3$.
⟦*bval b s*$_1$; (*c*, *s*$_1$) ⇒ *s*$_2$; *P* (*c*, *s*$_1$) *s*$_2$; (*WHILE b DO c*, *s*$_2$) ⇒ *s*$_3$;
P (*WHILE b DO c*, *s*$_2$) *s*$_3$⟧
⟹ *P* (*WHILE b DO c*, *s*$_1$) *s*$_3$⟧
⟹ *P x1 x2*

thm *big_step.induct*

This induction schema is almost perfect for our purposes, but our trick for reusing the tuple syntax means that the induction schema has two parameters instead of the c, s, and s' that we are likely to encounter. Splitting the tuple parameter fixes this:

lemmas *big_step_induct* = *big_step.induct*[*split_format*(*complete*)]
thm *big_step_induct*

⟦(*x1a*, *x1b*) ⇒ *x2a*; ⋀*s*. *P SKIP s s*; ⋀*x a s*. *P* (*x* ::= *a*) *s* (*s*(*x* := *aval a s*));
⋀*c*$_1$ *s*$_1$ *s*$_2$ *c*$_2$ *s*$_3$.
⟦(*c*$_1$, *s*$_1$) ⇒ *s*$_2$; *P c*$_1$ *s*$_1$ *s*$_2$; (*c*$_2$, *s*$_2$) ⇒ *s*$_3$; *P c*$_2$ *s*$_2$ *s*$_3$⟧
⟹ *P* (*c*$_1$;; *c*$_2$) *s*$_1$ *s*$_3$;
⋀*b s c*$_1$ *t c*$_2$.
⟦*bval b s*; (*c*$_1$, *s*) ⇒ *t*; *P c*$_1$ *s t*⟧ ⟹ *P* (*IF b THEN c*$_1$ *ELSE c*$_2$) *s t*;
⋀*b s c*$_2$ *t c*$_1$.
⟦¬ *bval b s*; (*c*$_2$, *s*) ⇒ *t*; *P c*$_2$ *s t*⟧ ⟹ *P* (*IF b THEN c*$_1$ *ELSE c*$_2$) *s t*;
⋀*b s c*. ¬ *bval b s* ⟹ *P* (*WHILE b DO c*) *s s*;
⋀*b s*$_1$ *c s*$_2$ *s*$_3$.
⟦*bval b s*$_1$; (*c*, *s*$_1$) ⇒ *s*$_2$; *P c s*$_1$ *s*$_2$; (*WHILE b DO c*, *s*$_2$) ⇒ *s*$_3$;
P (*WHILE b DO c*) *s*$_2$ *s*$_3$⟧
⟹ *P* (*WHILE b DO c*) *s*$_1$ *s*$_3$⟧
⟹ *P x1a x1b x2a*

3.2 Rule inversion

What can we deduce from (*SKIP*, *s*) ⇒ *t* ? That $s = t$. This is how we can automatically prove it:

inductive_cases *SkipE*[*elim*!]: (*SKIP*,*s*) ⇒ *t*
thm *SkipE*

This is an *elimination rule*. The [elim] attribute tells auto, blast and friends (but not simp!) to use it automatically; [elim!] means that it is applied eagerly.

Similarly for the other commands:

inductive_cases *AssignE*[*elim*!]: $(x ::= a,s) \Rightarrow t$
thm *AssignE*
inductive_cases *SeqE*[*elim*!]: $(c1;;c2,s1) \Rightarrow s3$
thm *SeqE*
inductive_cases *IfE*[*elim*!]: $(IF\ b\ THEN\ c1\ ELSE\ c2,s) \Rightarrow t$
thm *IfE*

inductive_cases *WhileE*[*elim*]: $(WHILE\ b\ DO\ c,s) \Rightarrow t$
thm *WhileE*

Only [elim]: [elim!] would not terminate.

An automatic example:

lemma $(IF\ b\ THEN\ SKIP\ ELSE\ SKIP,\ s) \Rightarrow t \Longrightarrow t = s$
by *blast*

Rule inversion by hand via the "cases" method:

```
lemma assumes (IF b THEN SKIP ELSE SKIP, s) ⇒ t
shows t = s
proof−
  from assms show ?thesis
  proof cases — inverting assms
    case IfTrue thm IfTrue
    thus ?thesis by blast
  next
    case IfFalse thus ?thesis by blast
  qed
qed
```

lemma *assign_simp*:
$(x ::= a,s) \Rightarrow s' \longleftrightarrow (s' = s(x := aval\ a\ s))$
by *auto*

An example combining rule inversion and derivations

```
lemma Seq_assoc:
  (c1;; c2;; c3, s) ⇒ s' ⟷ (c1;; (c2;; c3), s) ⇒ s'
proof
  assume (c1;; c2;; c3, s) ⇒ s'
  then obtain s1 s2 where
    c1: (c1, s) ⇒ s1 and
    c2: (c2, s1) ⇒ s2 and
    c3: (c3, s2) ⇒ s' by auto
```

from *c2 c3*
have (*c2*;; *c3*, *s1*) ⇒ *s′* **by** (*rule Seq*)
with *c1*
show (*c1*;; (*c2*;; *c3*), *s*) ⇒ *s′* **by** (*rule Seq*)
next
— The other direction is analogous
assume (*c1*;; (*c2*;; *c3*), *s*) ⇒ *s′*
thus (*c1*;; *c2*;; *c3*, *s*) ⇒ *s′* **by** *auto*
qed

3.3 Command Equivalence

We call two statements c and c' equivalent wrt. the big-step semantics when *c started in s terminates in s′ iff c′ started in the same s also terminates in the same s′*. Formally:

abbreviation
equiv_c :: *com* ⇒ *com* ⇒ *bool* (**infix** ∼ *50*) **where**
c ∼ *c′* ≡ (∀ *s t*. (*c*,*s*) ⇒ *t* = (*c′*,*s*) ⇒ *t*)

Warning: ∼ is the symbol written `\ < s i m >` (without spaces).

As an example, we show that loop unfolding is an equivalence transformation on programs:

lemma *unfold_while*:
(*WHILE b DO c*) ∼ (*IF b THEN c*;; *WHILE b DO c ELSE SKIP*) (**is** *?w* ∼ *?iw*)
proof −
— to show the equivalence, we look at the derivation tree for
— each side and from that construct a derivation tree for the other side
have (*?iw*, *s*) ⇒ *t* **if** *assm*: (*?w*, *s*) ⇒ *t* **for** *s t*
proof −
from *assm* **show** *?thesis*
proof *cases* — rule inversion on (*?w*, *s*) ⇒ *t*
case *WhileFalse*
thus *?thesis* **by** *blast*
next
case *WhileTrue*
from ‹*bval b s*› ‹(*?w*, *s*) ⇒ *t*› **obtain** *s′* **where**
(*c*, *s*) ⇒ *s′* **and** (*?w*, *s′*) ⇒ *t* **by** *auto*
— now we can build a derivation tree for the *IF*
— first, the body of the True-branch:
hence (*c*;; *?w*, *s*) ⇒ *t* **by** (*rule Seq*)
— then the whole *IF*
with ‹*bval b s*› **show** *?thesis* **by** (*rule IfTrue*)

```
    qed
  qed
  moreover
  — now the other direction:
  have (?w, s) ⇒ t if assm: (?iw, s) ⇒ t for s t
  proof −
    from assm show ?thesis
    proof cases — rule inversion on (?iw, s) ⇒ t
      case IfFalse
      hence s = t using ‹(?iw, s) ⇒ t› by blast
      thus ?thesis using ‹¬bval b s› by blast
    next
      case IfTrue
      — and for this, only the Seq-rule is applicable:
      from ‹(c;; ?w, s) ⇒ t› obtain s' where
        (c, s) ⇒ s' and (?w, s') ⇒ t by auto
      — with this information, we can build a derivation tree for WHILE
      with ‹bval b s› show ?thesis by (rule WhileTrue)
    qed
  qed
  ultimately
  show ?thesis by blast
qed
```

Luckily, such lengthy proofs are seldom necessary. Isabelle can prove many such facts automatically.

```
lemma while_unfold:
  (WHILE b DO c) ∼ (IF b THEN c;; WHILE b DO c ELSE SKIP)
by blast

lemma triv_if:
  (IF b THEN c ELSE c) ∼ c
by blast

lemma commute_if:
  (IF b1 THEN (IF b2 THEN c11 ELSE c12) ELSE c2)
   ∼
    (IF b2 THEN (IF b1 THEN c11 ELSE c2) ELSE (IF b1 THEN c12
ELSE c2))
by blast

lemma sim_while_cong_aux:
  (WHILE b DO c,s) ⇒ t  ⟹ c ∼ c' ⟹  (WHILE b DO c',s) ⇒ t
apply(induction WHILE b DO c s t arbitrary: b c rule: big_step_induct)
```

apply *blast*
apply *blast*
done

lemma *sim_while_cong*: $c \sim c' \Longrightarrow$ *WHILE b DO c* $\sim$ *WHILE b DO c'*
by (*metis sim_while_cong_aux*)

Command equivalence is an equivalence relation, i.e. it is reflexive, symmetric, and transitive. Because we used an abbreviation above, Isabelle derives this automatically.

lemma *sim_refl*: $c \sim c$ **by** *simp*
lemma *sim_sym*: $(c \sim c') = (c' \sim c)$ **by** *auto*
lemma *sim_trans*: $c \sim c' \Longrightarrow c' \sim c'' \Longrightarrow c \sim c''$ **by** *auto*

3.4 Execution is deterministic

This proof is automatic.

theorem *big_step_determ*: $[\![(c,s) \Rightarrow t; (c,s) \Rightarrow u]\!] \Longrightarrow u = t$
  by (*induction arbitrary*: *u rule*: *big_step.induct*) *blast+*

This is the proof as you might present it in a lecture. The remaining cases are simple enough to be proved automatically:

theorem
  $(c,s) \Rightarrow t \Longrightarrow (c,s) \Rightarrow t' \Longrightarrow t' = t$
proof (*induction arbitrary*: t' *rule*: *big_step.induct*)
  — the only interesting case, *WhileTrue*:
  fix $b\ c\ s\ s_1\ t\ t'$
  — The assumptions of the rule:
  assume *bval b s* **and** $(c,s) \Rightarrow s_1$ **and** (*WHILE b DO c*,s_1) $\Rightarrow t$
  — Ind.Hyp; note the $\bigwedge$ because of arbitrary:
  assume *IHc*: $\bigwedge t'.\ (c,s) \Rightarrow t' \Longrightarrow t' = s_1$
  assume *IHw*: $\bigwedge t'.$ (*WHILE b DO c*,s_1) $\Rightarrow t' \Longrightarrow t' = t$
  — Premise of implication:
  assume (*WHILE b DO c*,s) $\Rightarrow t'$
  with ‹*bval b s*› **obtain** s_1' **where**
    c: $(c,s) \Rightarrow s_1'$ **and**
    w: (*WHILE b DO c*,s_1') $\Rightarrow t'$
   by *auto*
  from *c IHc* **have** $s_1' = s_1$ **by** *blast*
  with *w IHw* **show** $t' = t$ **by** *blast*
qed *blast+* — prove the rest automatically

end

4 Small-Step Semantics of Commands

theory *Small_Step* **imports** *Star Big_Step* **begin**

4.1 The transition relation

inductive
  small_step :: *com* ∗ *state* ⇒ *com* ∗ *state* ⇒ *bool* (**infix** → *55*)
where
Assign: (*x* ::= *a*, *s*) → (*SKIP*, *s*(*x* := *aval a s*)) |

Seq1:   (*SKIP*;;c_2,*s*) → (c_2,*s*) |
Seq2:   (c_1,*s*) → (c_1',*s'*) ⟹ (c_1;;c_2,*s*) → (c_1';;c_2,*s'*) |

IfTrue: *bval b s* ⟹ (*IF b THEN* c_1 *ELSE* c_2,*s*) → (c_1,*s*) |
IfFalse: ¬*bval b s* ⟹ (*IF b THEN* c_1 *ELSE* c_2,*s*) → (c_2,*s*) |

While: (*WHILE b DO c*,*s*) →
            (*IF b THEN c*;; *WHILE b DO c ELSE SKIP*,*s*)

abbreviation
  small_steps :: *com* ∗ *state* ⇒ *com* ∗ *state* ⇒ *bool* (**infix** →∗ *55*)
where *x* →∗ *y* == *star small_step x y*

4.2 Executability

code_pred *small_step* **.**

values {(*c'*,*map t* [''*x*'',''*y*'',''*z*'']) |*c' t*.
   (''*x*'' ::= *V* ''*z*'';; ''*y*'' ::= *V* ''*x*'',
    <''*x*'' := *3*, ''*y*'' := *7*, ''*z*'' := *5*>) →∗ (*c'*,*t*)}

4.3 Proof infrastructure

4.3.1 Induction rules

The default induction rule *small_step.induct* only works for lemmas of the form *a* → *b* ⟹ ... where *a* and *b* are not already pairs (*DUMMY*,*DUMMY*). We can generate a suitable variant of *small_step.induct* for pairs by "splitting" the arguments → into pairs:

lemmas *small_step_induct* = *small_step.induct*[*split_format*(*complete*)]

4.3.2 Proof automation

declare *small_step.intros*[*simp*,*intro*]

Rule inversion:

inductive_cases *SkipE*[*elim*!]: (*SKIP*,*s*) → *ct*
thm *SkipE*
inductive_cases *AssignE*[*elim*!]: (*x*::=*a*,*s*) → *ct*
thm *AssignE*
inductive_cases *SeqE*[*elim*]: (*c1*;;*c2*,*s*) → *ct*
thm *SeqE*
inductive_cases *IfE*[*elim*!]: (*IF b THEN c1 ELSE c2*,*s*) → *ct*
inductive_cases *WhileE*[*elim*]: (*WHILE b DO c*, *s*) → *ct*

A simple property:

lemma *deterministic*:
  cs → *cs′* ⟹ *cs* → *cs″* ⟹ *cs″* = *cs′*
apply(*induction arbitrary*: *cs″ rule*: *small_step.induct*)
apply *blast*+
done

4.4 Equivalence with big-step semantics

lemma *star_seq2*: (*c1*,*s*) →∗ (*c1′*,*s′*) ⟹ (*c1*;;*c2*,*s*) →∗ (*c1′*;;*c2*,*s′*)
proof(*induction rule*: *star_induct*)
  case *refl* **thus** *?case* **by** *simp*
next
  case *step*
  thus *?case* **by** (*metis Seq2 star.step*)
qed

lemma *seq_comp*:
  ⟦ (*c1*,*s1*) →∗ (*SKIP*,*s2*); (*c2*,*s2*) →∗ (*SKIP*,*s3*) ⟧
   ⟹ (*c1*;;*c2*, *s1*) →∗ (*SKIP*,*s3*)
by(*blast intro*: *star.step star_seq2 star_trans*)

The following proof corresponds to one on the board where one would show chains of → and →∗ steps.

lemma *big_to_small*:
  cs ⇒ *t* ⟹ *cs* →∗ (*SKIP*,*t*)
proof (*induction rule*: *big_step.induct*)
  fix *s* **show** (*SKIP*,*s*) →∗ (*SKIP*,*s*) **by** *simp*
next
  fix *x a s* **show** (*x* ::= *a*,*s*) →∗ (*SKIP*, *s*(*x* := *aval a s*)) **by** *auto*
next

```
  fix c1 c2 s1 s2 s3
  assume (c1,s1) →* (SKIP,s2) and (c2,s2) →* (SKIP,s3)
  thus (c1;;c2, s1) →* (SKIP,s3) by (rule seq_comp)
next
  fix s::state and b c0 c1 t
  assume bval b s
  hence (IF b THEN c0 ELSE c1,s) → (c0,s) by simp
  moreover assume (c0,s) →* (SKIP,t)
  ultimately
  show (IF b THEN c0 ELSE c1,s) →* (SKIP,t) by (metis star.simps)
next
  fix s::state and b c0 c1 t
  assume ¬bval b s
  hence (IF b THEN c0 ELSE c1,s) → (c1,s) by simp
  moreover assume (c1,s) →* (SKIP,t)
  ultimately
  show (IF b THEN c0 ELSE c1,s) →* (SKIP,t) by (metis star.simps)
next
  fix b c and s::state
  assume b: ¬bval b s
  let ?if = IF b THEN c;; WHILE b DO c ELSE SKIP
  have (WHILE b DO c,s) → (?if, s) by blast
  moreover have (?if,s) → (SKIP, s) by (simp add: b)
   ultimately show (WHILE b DO c,s) →* (SKIP,s) by(metis star.refl
star.step)
next
  fix b c s s' t
  let ?w = WHILE b DO c
  let ?if = IF b THEN c;; ?w ELSE SKIP
  assume w: (?w,s') →* (SKIP,t)
  assume c: (c,s) →* (SKIP,s')
  assume b: bval b s
  have (?w,s) → (?if, s) by blast
  moreover have (?if, s) → (c;; ?w, s) by (simp add: b)
  moreover have (c;; ?w,s) →* (SKIP,t) by(rule seq_comp[OF c w])
 ultimately show (WHILE b DO c,s) →* (SKIP,t) by (metis star.simps)
qed
```

Each case of the induction can be proved automatically:

```
lemma cs ⇒ t ⟹ cs →* (SKIP,t)
proof (induction rule: big_step.induct)
  case Skip show ?case by blast
next
  case Assign show ?case by blast
```

next
  case *Seq* **thus** *?case* **by** (*blast intro*: *seq_comp*)
next
  case *IfTrue* **thus** *?case* **by** (*blast intro*: *star.step*)
next
  case *IfFalse* **thus** *?case* **by** (*blast intro*: *star.step*)
next
  case *WhileFalse* **thus** *?case*
    by (*metis star.step star_step1 small_step.IfFalse small_step.While*)
next
  case *WhileTrue*
  thus *?case*
    by(*metis While seq_comp small_step.IfTrue star.step*[*of small_step*])
qed

lemma *small1_big_continue*:
  $cs \to cs' \Longrightarrow cs' \Rightarrow t \Longrightarrow cs \Rightarrow t$
apply (*induction arbitrary*: *t rule*: *small_step.induct*)
apply *auto*
done

lemma *small_to_big*:
  $cs \to* (SKIP,t) \Longrightarrow cs \Rightarrow t$
apply (*induction cs* (*SKIP*,*t*) *rule*: *star.induct*)
apply (*auto intro*: *small1_big_continue*)
done

Finally, the equivalence theorem:

theorem *big_iff_small*:
  $cs \Rightarrow t = cs \to* (SKIP,t)$
by(*metis big_to_small small_to_big*)

4.5 Final configurations and infinite reductions

definition *final cs* $\longleftrightarrow \neg(\exists\, cs'.\ cs \to cs')$

lemma *finalD*: *final* $(c,s) \Longrightarrow c = SKIP$
apply(*simp add*: *final_def*)
apply(*induction c*)
apply *blast+*
done

lemma *final_iff_SKIP*: *final* $(c,s) = (c = SKIP)$
by (*metis SkipE finalD final_def*)

Now we can show that ⇒ yields a final state iff → terminates:

lemma *big_iff_small_termination*:
$(\exists\, t.\ cs \Rightarrow t) \longleftrightarrow (\exists\, cs'.\ cs \rightarrow* cs' \wedge final\ cs')$
by(*simp add*: *big_iff_small final_iff_SKIP*)

This is the same as saying that the absence of a big step result is equivalent with absence of a terminating small step sequence, i.e. with nontermination. Since → is determininistic, there is no difference between may and must terminate.

end
theory *Finite_Reachable*
imports *Small_Step*
begin

4.6 Finite number of reachable commands

This theory shows that in the small-step semantics one can only reach a finite number of commands from any given command. Hence one can see the command component of a small-step configuration as a combination of the program to be executed and a pc.

definition *reachable* :: *com* ⇒ *com set* **where**
$reachable\ c = \{c'.\ \exists\, s\ t.\ (c,s) \rightarrow* (c',t)\}$

Proofs need induction on the length of a small-step reduction sequence.

fun *small_stepsn* :: *com* ∗ *state* ⇒ *nat* ⇒ *com* ∗ *state* ⇒ *bool*
  (_ →′(_′) _ [55,0,55] 55) **where**
$(cs \rightarrow(0)\ cs') = (cs' = cs)$ |
$cs \rightarrow(Suc\ n)\ cs'' = (\exists\, cs'.\ cs \rightarrow cs' \wedge cs' \rightarrow(n)\ cs'')$

lemma *stepsn_if_star*: $cs \rightarrow* cs' \Longrightarrow \exists\, n.\ cs \rightarrow(n)\ cs'$
proof(*induction rule*: *star.induct*)
  case *refl* **show** *?case* **by** (*metis small_stepsn.simps*(*1*))
next
  case *step* **thus** *?case* **by** (*metis small_stepsn.simps*(*2*))
qed

lemma *star_if_stepsn*: $cs \rightarrow(n)\ cs' \Longrightarrow cs \rightarrow* cs'$
by(*induction n arbitrary*: *cs*) (*auto elim*: *star.step*)

lemma *SKIP_starD*: $(SKIP,\ s) \rightarrow* (c,t) \Longrightarrow c = SKIP$
by(*induction SKIP s c t rule*: *star_induct*) *auto*

lemma *reachable_SKIP*: $reachable\ SKIP = \{SKIP\}$
by(*auto simp*: *reachable_def dest*: *SKIP_starD*)

```
lemma Assign_starD: (x::=a, s) →* (c,t) ⟹ c ∈ {x::=a, SKIP}
by (induction x::=a s c t rule: star_induct) (auto dest: SKIP_starD)

lemma reachable_Assign: reachable (x::=a) = {x::=a, SKIP}
by(auto simp: reachable_def dest:Assign_starD)


lemma Seq_stepsnD: (c1;; c2, s) →(n) (c', t) ⟹
  (∃ c1' m. c' = c1';; c2 ∧ (c1, s) →(m) (c1', t) ∧ m ≤ n) ∨
   (∃ s2 m1 m2. (c1,s) →(m1) (SKIP,s2) ∧ (c2, s2) →(m2) (c', t) ∧
m1+m2 < n)
proof(induction n arbitrary: c1 c2 s)
  case 0 thus ?case by auto
next
  case (Suc n)
  from Suc.prems obtain s' c12' where (c1;;c2, s) → (c12', s')
    and n: (c12',s') →(n) (c',t) by auto
  from this(1) show ?case
  proof
    assume c1 = SKIP (c12', s') = (c2, s)
    hence (c1,s) →(0) (SKIP, s') ∧ (c2, s') →(n) (c', t) ∧ 0 + n < Suc n
      using n by auto
    thus ?case by blast
  next
    fix c1' s'' assume 1: (c12', s') = (c1';; c2, s'') (c1, s) → (c1', s'')
    hence n': (c1';;c2,s') →(n) (c',t) using n by auto
    from Suc.IH[OF n'] show ?case
    proof
      assume ∃ c1'' m. c' = c1'';; c2 ∧ (c1', s') →(m) (c1'', t) ∧ m ≤ n
        (is ∃ a b. ?P a b)
      then obtain c1'' m where 2: ?P c1'' m by blast
      hence c' = c1'';;c2 ∧ (c1, s) →(Suc m) (c1'',t) ∧ Suc m ≤ Suc n
        using 1 by auto
      thus ?case by blast
    next
      assume ∃ s2 m1 m2. (c1',s') →(m1) (SKIP,s2) ∧
        (c2,s2) →(m2) (c',t) ∧ m1+m2 < n (is ∃ a b c. ?P a b c)
      then obtain s2 m1 m2 where ?P s2 m1 m2 by blast
      hence (c1,s) →(Suc m1) (SKIP,s2) ∧ (c2,s2) →(m2) (c',t) ∧
        Suc m1 + m2 < Suc n  using 1 by auto
      thus ?case by blast
    qed
```

  qed
qed

corollary *Seq_starD*: $(c1;;\ c2,\ s) \to* (c',\ t) \Longrightarrow$
  $(\exists\, c1'.\ c' = c1';;\ c2 \wedge (c1,\ s) \to* (c1',\ t)) \vee$
  $(\exists\, s2.\ (c1,s) \to* (SKIP,s2) \wedge (c2,\ s2) \to* (c',\ t))$
by(*metis Seq_stepsnD star_if_stepsn stepsn_if_star*)

lemma *reachable_Seq*: *reachable* $(c1;;c2) \subseteq$
  $(\lambda c1'.\ c1';;c2)$ ‘ *reachable c1* ∪ *reachable c2*
by(*auto simp*: *reachable_def image_def dest*!: *Seq_starD*)

lemma *If_starD*: $(IF\ b\ THEN\ c1\ ELSE\ c2,\ s) \to* (c,t) \Longrightarrow$
  $c = IF\ b\ THEN\ c1\ ELSE\ c2 \vee (c1,s) \to* (c,t) \vee (c2,s) \to* (c,t)$
by(*induction IF b THEN c1 ELSE c2 s c t rule*: *star_induct*) *auto*

lemma *reachable_If*: *reachable* (*IF b THEN c1 ELSE c2*) ⊆
  {*IF b THEN c1 ELSE c2*} ∪ *reachable c1* ∪ *reachable c2*
by(*auto simp*: *reachable_def dest*!: *If_starD*)

lemma *While_stepsnD*: $(WHILE\ b\ DO\ c,\ s) \to(n)\ (c2,t) \Longrightarrow$
  $c2 \in \{WHILE\ b\ DO\ c,\ IF\ b\ THEN\ c\ ;;\ WHILE\ b\ DO\ c\ ELSE\ SKIP,$
$SKIP\}$
  $\vee (\exists\, c1.\ c2 = c1\ ;;\ WHILE\ b\ DO\ c \wedge (\exists\ s1\ s2.\ (c,s1) \to* (c1,s2)))$
proof(*induction n arbitrary*: *s rule*: *less_induct*)
  case (*less n1*)
  show *?case*
  proof(*cases n1*)
    case *0* **thus** *?thesis* **using** *less.prems* **by** (*simp*)
  next
    case (*Suc n2*)
    let *?w* = *WHILE b DO c*
    let *?iw* = *IF b THEN c ;; ?w ELSE SKIP*
     from *Suc less.prems* **have** *n2*: $(?iw,s) \to(n2)\ (c2,t)$ **by**(*auto elim*!:
WhileE)
    show *?thesis*
    proof(*cases n2*)
      case *0* **thus** *?thesis* **using** *n2* **by** *auto*
    next
      case (*Suc n3*)
      then obtain *iw′ s′* **where** $(?iw,s) \to (iw',s')$
        and *n3*: $(iw',s') \to(n3)\ (c2,t)$ **using** *n2* **by** *auto*

from *this*(*1*)
show *?thesis*
proof
assume (*iw*′, *s*′) = (*c*;; *WHILE b DO c*, *s*)
with *n3* **have** (*c*;;*?w*, *s*) →(*n3*) (*c2*,*t*) **by** *auto*
from *Seq_stepsnD*[*OF this*] **show** *?thesis*
proof
assume ∃ *c1*′ *m*. *c2* = *c1*′;; *?w* ∧ (*c*,*s*) →(*m*) (*c1*′, *t*) ∧ *m* ≤ *n3*
thus *?thesis* **by** (*metis star_if_stepsn*)
next
assume ∃ *s2 m1 m2*. (*c*, *s*) →(*m1*) (*SKIP*, *s2*) ∧
(*WHILE b DO c*, *s2*) →(*m2*) (*c2*, *t*) ∧ *m1* + *m2* < *n3* (**is** ∃ *x y z*. *?P x y z*)
then obtain *s2 m1 m2* **where** *?P s2 m1 m2* **by** *blast*
with ‹*n2* = *Suc n3*› ‹*n1* = *Suc n2*›**have** *m2* < *n1* **by** *arith*
from *less.IH*[*OF this*] ‹*?P s2 m1 m2*› **show** *?thesis* **by** *blast*
qed
next
assume (*iw*′, *s*′) = (*SKIP*, *s*)
thus *?thesis* **using** *star_if_stepsn*[*OF n3*] **by**(*auto dest*!: *SKIP_starD*)
qed
qed
qed
qed

lemma *reachable_While*: *reachable* (*WHILE b DO c*) ⊆
{*WHILE b DO c*, *IF b THEN c* ;; *WHILE b DO c ELSE SKIP*, *SKIP*} ∪
(λ*c*′. *c*′ ;; *WHILE b DO c*) ‘ *reachable c*
apply(*auto simp*: *reachable_def image_def*)
by (*metis While_stepsnD insertE singletonE stepsn_if_star*)

theorem *finite_reachable*: *finite*(*reachable c*)
apply(*induction c*)
apply(*auto simp*: *reachable_SKIP reachable_Assign*
finite_subset[*OF reachable_Seq*] *finite_subset*[*OF reachable_If*]
finite_subset[*OF reachable_While*])
done

end

5 Denotational Semantics of Commands

theory *Denotational* **imports** *Big_Step* **begin**

type_synonym *com_den* = (*state* × *state*) *set*

definition *W* :: (*state* ⇒ *bool*) ⇒ *com_den* ⇒ (*com_den* ⇒ *com_den*)
where
W db dc = (λ*dw*. {(*s*,*t*). *if db s then* (*s*,*t*) ∈ *dc O dw else s*=*t*})

fun *D* :: *com* ⇒ *com_den* **where**
D SKIP = *Id* |
D (*x* ::= *a*) = {(*s*,*t*). *t* = *s*(*x* := *aval a s*)} |
D (*c1*;;*c2*) = *D*(*c1*) *O D*(*c2*) |
D (*IF b THEN c1 ELSE c2*)
= {(*s*,*t*). *if bval b s then* (*s*,*t*) ∈ *D c1 else* (*s*,*t*) ∈ *D c2*} |
D (*WHILE b DO c*) = *lfp* (*W* (*bval b*) (*D c*))

lemma *W_mono*: *mono* (*W b r*)
by (*unfold W_def mono_def*) *auto*

lemma *D_While_If*:
D(*WHILE b DO c*) = *D*(*IF b THEN c*;;*WHILE b DO c ELSE SKIP*)
proof−
let *?w* = *WHILE b DO c* **let** *?f* = *W* (*bval b*) (*D c*)
have *D ?w* = *lfp ?f* **by** *simp*
also have ... = *?f* (*lfp ?f*) **by**(*rule lfp_unfold* [*OF W_mono*])
also have ... = *D*(*IF b THEN c*;;*?w ELSE SKIP*) **by** (*simp add*: *W_def*)
finally show *?thesis* **.**
qed

Equivalence of denotational and big-step semantics:

lemma *D_if_big_step*: (*c*,*s*) ⇒ *t* ⟹ (*s*,*t*) ∈ *D*(*c*)
proof (*induction rule*: *big_step_induct*)
case *WhileFalse*
with *D_While_If* **show** *?case* **by** *auto*
next
case *WhileTrue*
show *?case* **unfolding** *D_While_If* **using** *WhileTrue* **by** *auto*
qed *auto*

abbreviation *Big_step* :: *com* ⇒ *com_den* **where**
Big_step c ≡ {(*s*,*t*). (*c*,*s*) ⇒ *t*}

```
lemma Big_step_if_D:  (s,t) ∈ D(c) ⟹ (s,t) ∈ Big_step c
proof (induction c arbitrary: s t)
  case Seq thus ?case by fastforce
next
  case (While b c)
  let ?B = Big_step (WHILE b DO c) let ?f = W (bval b) (D c)
  have ?f ?B ⊆ ?B using While.IH by (auto simp: W_def)
  from lfp_lowerbound[where ?f = ?f, OF this] While.prems
  show ?case by auto
qed (auto split: if_splits)

theorem denotational_is_big_step:
  (s,t) ∈ D(c)  =  ((c,s) ⇒ t)
by (metis D_if_big_step Big_step_if_D[simplified])

corollary equiv_c_iff_equal_D: (c1 ∼ c2) ⟷ D c1 = D c2
by(simp add: denotational_is_big_step[symmetric] set_eq_iff)
```

5.1 Continuity

```
definition chain :: (nat ⇒ 'a set) ⇒ bool where
chain S = (∀ i. S i ⊆ S(Suc i))

lemma chain_total: chain S ⟹ S i ≤ S j ∨ S j ≤ S i
by (metis chain_def le_cases lift_Suc_mono_le)

definition cont :: ('a set ⇒ 'b set) ⇒ bool where
cont f = (∀ S. chain S ⟶ f(UN n. S n) = (UN n. f(S n)))

lemma mono_if_cont: fixes f :: 'a set ⇒ 'b set
  assumes cont f shows mono f
proof
  fix a b :: 'a set assume a ⊆ b
  let ?S = λn::nat. if n=0 then a else b
  have chain ?S using ‹a ⊆ b› by(auto simp: chain_def)
  hence f(UN n. ?S n) = (UN n. f(?S n))
    using assms by (simp add: cont_def del: if_image_distrib)
  moreover have (UN n. ?S n) = b using ‹a ⊆ b› by (auto split: if_splits)
  moreover have (UN n. f(?S n)) = f a ∪ f b by (auto split: if_splits)
  ultimately show f a ⊆ f b by (metis Un_upper1)
qed

lemma chain_iterates: fixes f :: 'a set ⇒ 'a set
  assumes mono f shows chain(λn. (f^^n) {})
```

```
proof−
  have (f ^^ n) {} ⊆ (f ^^ Suc n) {} for n
  proof (induction n)
    case 0 show ?case by simp
  next
    case (Suc n) thus ?case using assms by (auto simp: mono_def)
  qed
  thus ?thesis by(auto simp: chain_def assms)
qed

theorem lfp_if_cont:
  assumes cont f shows lfp f = (UN n. (f^^n) {}) (is _ = ?U)
proof
  from assms mono_if_cont
  have mono: (f ^^ n) {} ⊆ (f ^^ Suc n) {} for n
    using funpow_decreasing [of n Suc n] by auto
  show lfp f ⊆ ?U
  proof (rule lfp_lowerbound)
    have f ?U = (UN n. (f^^Suc n){})
      using chain_iterates[OF mono_if_cont[OF assms]] assms
      by(simp add: cont_def)
    also have . . . = (f^^0){} ∪ . . . by simp
    also have . . . = ?U
       using mono by auto (metis funpow_simps_right(2) funpow_swap1
o_apply)
    finally show f ?U ⊆ ?U by simp
  qed
next
  have (f^^n){} ⊆ p if f p ⊆ p for n p
  proof −
    show ?thesis
    proof(induction n)
      case 0 show ?case by simp
    next
      case Suc
      from monoD[OF mono_if_cont[OF assms] Suc] ‹f p ⊆ p›
      show ?case by simp
    qed
  qed
  thus ?U ⊆ lfp f by(auto simp: lfp_def)
qed

lemma cont_W: cont(W b r)
by(auto simp: cont_def W_def)
```

5.2 The denotational semantics is deterministic

```
lemma single_valued_UN_chain:
  assumes chain S (⋀n. single_valued (S n))
  shows single_valued(UN n. S n)
proof(auto simp: single_valued_def)
  fix m n x y z assume (x, y) ∈ S m (x, z) ∈ S n
  with chain_total[OF assms(1), of m n] assms(2)
  show y = z by (auto simp: single_valued_def)
qed

lemma single_valued_lfp: fixes f :: com_den ⇒ com_den
assumes cont f ⋀r. single_valued r ⟹ single_valued (f r)
shows single_valued(lfp f)
unfolding lfp_if_cont[OF assms(1)]
proof(rule single_valued_UN_chain[OF chain_iterates[OF mono_if_cont[OF
assms(1)]]])
  fix n show single_valued ((f ^^ n) {})
  by(induction n)(auto simp: assms(2))
qed

lemma single_valued_D: single_valued (D c)
proof(induction c)
  case Seq thus ?case by(simp add: single_valued_relcomp)
next
  case (While b c)
  let ?f = W (bval b) (D c)
  have single_valued (lfp ?f)
  proof(rule single_valued_lfp[OF cont_W])
    show ⋀r. single_valued r ⟹ single_valued (?f r)
      using While.IH by(force simp: single_valued_def W_def)
  qed
  thus ?case by simp
qed (auto simp add: single_valued_def)

end
```

6 Compiler for IMP

```
theory Compiler imports Big_Step Star
begin
```

6.1 List setup

In the following, we use the length of lists as integers instead of natural numbers. Instead of converting *nat* to *int* explicitly, we tell Isabelle to coerce *nat* automatically when necessary.

```
declare [[coercion_enabled]]
declare [[coercion int :: nat ⇒ int]]
```

Similarly, we will want to access the ith element of a list, where *i* is an *int*.

```
fun inth :: 'a list ⇒ int ⇒ 'a (infixl !! 100) where
(x # xs) !! i = (if i = 0 then x else xs !! (i − 1))
```

The only additional lemma we need about this function is indexing over append:

```
lemma inth_append [simp]:
  0 ≤ i ⟹
  (xs @ ys) !! i = (if i < size xs then xs !! i else ys !! (i − size xs))
by (induction xs arbitrary: i) (auto simp: algebra_simps)
```

We hide coercion *int* applied to *length*:

```
abbreviation (output)
  isize xs == int (length xs)

notation isize (size)
```

6.2 Instructions and Stack Machine

```
datatype instr =
  LOADI int | LOAD vname | ADD | STORE vname |
  JMP int | JMPLESS int | JMPGE int
type_synonym stack = val list
type_synonym config = int × state × stack

abbreviation hd2 xs == hd(tl xs)
abbreviation tl2 xs == tl(tl xs)

fun iexec :: instr ⇒ config ⇒ config where
iexec instr (i,s,stk) = (case instr of
  LOADI n ⇒ (i+1,s, n#stk) |
  LOAD x ⇒ (i+1,s, s x # stk) |
  ADD ⇒ (i+1,s, (hd2 stk + hd stk) # tl2 stk) |
  STORE x ⇒ (i+1,s(x := hd stk),tl stk) |
  JMP n ⇒  (i+1+n,s,stk) |
  JMPLESS n ⇒ (if hd2 stk < hd stk then i+1+n else i+1,s,tl2 stk) |
```

JMPGE n ⇒ (*if hd2 stk* >= *hd stk then i*+*1*+*n else i*+*1*,*s*,*tl2 stk*))

definition
exec1 :: *instr list* ⇒ *config* ⇒ *config* ⇒ *bool*
((_/ ⊢ (_ →/ _)) [*59*,*0*,*59*] *60*)
where
P ⊢ *c* → *c′* =
(∃ *i s stk*. *c* = (*i*,*s*,*stk*) ∧ *c′* = *iexec*(*P*!!*i*) (*i*,*s*,*stk*) ∧ *0* ≤ *i* ∧ *i* < *size P*)

lemma *exec1I* [*intro*, *code_pred_intro*]:
c′ = *iexec* (*P*!!*i*) (*i*,*s*,*stk*) ⟹ *0* ≤ *i* ⟹ *i* < *size P*
⟹ *P* ⊢ (*i*,*s*,*stk*) → *c′*
by (*simp add*: *exec1_def*)

abbreviation
exec :: *instr list* ⇒ *config* ⇒ *config* ⇒ *bool* ((_/ ⊢ (_ →∗/ _)) *50*)
where
exec P ≡ *star* (*exec1 P*)

lemmas *exec_induct* = *star.induct* [*of exec1 P*, *split_format*(*complete*)]

code_pred *exec1* **by** (*metis exec1_def*)

values
{(*i*,*map t* [′′*x*′′,′′*y*′′],*stk*) | *i t stk*.
[*LOAD* ′′*y*′′, *STORE* ′′*x*′′] ⊢
(*0*, <′′*x*′′ := *3*, ′′*y*′′ := *4*>, []) →∗ (*i*,*t*,*stk*)}

6.3 Verification infrastructure

Below we need to argue about the execution of code that is embedded in larger programs. For this purpose we show that execution is preserved by appending code to the left or right of a program.

lemma *iexec_shift* [*simp*]:
((*n*+*i′*,*s′*,*stk′*) = *iexec x* (*n*+*i*,*s*,*stk*)) = ((*i′*,*s′*,*stk′*) = *iexec x* (*i*,*s*,*stk*))
by(*auto split*:*instr.split*)

lemma *exec1_appendR*: *P* ⊢ *c* → *c′* ⟹ *P*@*P′* ⊢ *c* → *c′*
by (*auto simp*: *exec1_def*)

lemma *exec_appendR*: *P* ⊢ *c* →∗ *c′* ⟹ *P*@*P′* ⊢ *c* →∗ *c′*
by (*induction rule*: *star.induct*) (*fastforce intro*: *star.step exec1_appendR*)+

lemma *exec1_appendL*:

fixes *i i′* :: *int*
shows
$P \vdash (i,s,stk) \rightarrow (i',s',stk') \Longrightarrow$
$P' \; @ \; P \vdash (size(P')+i,s,stk) \rightarrow (size(P')+i',s',stk')$
unfolding *exec1_def*
by (*auto simp del*: *iexec.simps*)

lemma *exec_appendL*:
fixes *i i′* :: *int*
shows
$P \vdash (i,s,stk) \rightarrow* (i',s',stk') \Longrightarrow$
$P' \; @ \; P \vdash (size(P')+i,s,stk) \rightarrow* (size(P')+i',s',stk')$
by (*induction rule*: *exec_induct*) (*blast intro*: *star.step exec1_appendL*)+

Now we specialise the above lemmas to enable automatic proofs of $P \vdash c \rightarrow* c'$ where P is a mixture of concrete instructions and pieces of code that we already know how they execute (by induction), combined by @ and #. Backward jumps are not supported. The details should be skipped on a first reading.

If we have just executed the first instruction of the program, drop it:

lemma *exec_Cons_1* [*intro*]:
$P \vdash (0,s,stk) \rightarrow* (j,t,stk') \Longrightarrow$
$instr\#P \vdash (1,s,stk) \rightarrow* (1+j,t,stk')$
by (*drule exec_appendL*[**where** *P′*=[*instr*]]) *simp*

lemma *exec_appendL_if*[*intro*]:
fixes *i i′ j* :: *int*
shows
$size\ P' <= i$
$\Longrightarrow P \vdash (i - size\ P',s,stk) \rightarrow* (j,s',stk')$
$\Longrightarrow i' = size\ P' + j$
$\Longrightarrow P' \; @ \; P \vdash (i,s,stk) \rightarrow* (i',s',stk')$
by (*drule exec_appendL*[**where** *P′*=*P′*]) *simp*

Split the execution of a compound program up into the execution of its parts:

lemma *exec_append_trans*[*intro*]:
fixes *i′ i″ j″* :: *int*
shows
$P \vdash (0,s,stk) \rightarrow* (i',s',stk') \Longrightarrow$
$size\ P \le i' \Longrightarrow$
$P' \vdash (i' - size\ P,s',stk') \rightarrow* (i'',s'',stk'') \Longrightarrow$
$j'' = size\ P + i''$
$\Longrightarrow$

```
 P @ P' ⊢ (0,s,stk) →* (j'',s'',stk'')
by(metis star_trans[OF exec_appendR exec_appendL_if])


declare Let_def[simp]
```

6.4 Compilation

```
fun acomp :: aexp ⇒ instr list where
acomp (N n) = [LOADI n] |
acomp (V x) = [LOAD x] |
acomp (Plus a1 a2) = acomp a1 @ acomp a2 @ [ADD]

lemma acomp_correct[intro]:
  acomp a ⊢ (0,s,stk) →* (size(acomp a),s,aval a s#stk)
by (induction a arbitrary: stk) fastforce+

fun bcomp :: bexp ⇒ bool ⇒ int ⇒ instr list where
bcomp (Bc v) f n = (if v=f then [JMP n] else []) |
bcomp (Not b) f n = bcomp b (¬f) n |
bcomp (And b1 b2) f n =
 (let cb2 = bcomp b2 f n;
        m = if f then size cb2 else (size cb2)+n;
      cb1 = bcomp b1 False m
  in cb1 @ cb2) |
bcomp (Less a1 a2) f n =
 acomp a1 @ acomp a2 @ (if f then [JMPLESS n] else [JMPGE n])

value
  bcomp (And (Less (V ''x'') (V ''y'')) (Not(Less (V ''u'') (V ''v''))))
     False 3

lemma bcomp_correct[intro]:
  fixes n :: int
  shows
  0 ≤ n ⟹
  bcomp b f n ⊢
 (0,s,stk)  →*  (size(bcomp b f n) + (if f = bval b s then n else 0),s,stk)
proof(induction b arbitrary: f n)
  case Not
  from Not(1)[where f=~f] Not(2) show ?case by fastforce
next
  case (And b1 b2)
  from And(1)[of if f then size(bcomp b2 f n) else size(bcomp b2 f n) + n
```

```
              False]
       And(2)[of n f] And(3)
  show ?case by fastforce
qed fastforce+

fun ccomp :: com ⇒ instr list where
ccomp SKIP = [] |
ccomp (x ::= a) = acomp a @ [STORE x] |
ccomp (c₁;;c₂) = ccomp c₁ @ ccomp c₂ |
ccomp (IF b THEN c₁ ELSE c₂) =
  (let cc₁ = ccomp c₁; cc₂ = ccomp c₂; cb = bcomp b False (size cc₁ + 1)
   in cb @ cc₁ @ JMP (size cc₂) # cc₂) |
ccomp (WHILE b DO c) =
 (let cc = ccomp c; cb = bcomp b False (size cc + 1)
  in cb @ cc @ [JMP (−(size cb + size cc + 1))])


value ccomp
 (IF Less (V ''u'') (N 1) THEN ''u'' ::= Plus (V ''u'') (N 1)
  ELSE ''v'' ::= V ''u'')

value ccomp (WHILE Less (V ''u'') (N 1) DO (''u'' ::= Plus (V ''u'') (N
1)))
```

6.5 Preservation of semantics

```
lemma ccomp_bigstep:
  (c,s) ⇒ t ⟹ ccomp c ⊢ (0,s,stk) →* (size(ccomp c),t,stk)
proof(induction arbitrary: stk rule: big_step_induct)
  case (Assign x a s)
  show ?case by (fastforce simp:fun_upd_def cong: if_cong)
next
  case (Seq c1 s1 s2 c2 s3)
  let ?cc1 = ccomp c1  let ?cc2 = ccomp c2
  have ?cc1 @ ?cc2 ⊢ (0,s1,stk) →* (size ?cc1,s2,stk)
    using Seq.IH(1) by fastforce
  moreover
  have ?cc1 @ ?cc2 ⊢ (size ?cc1,s2,stk) →* (size(?cc1 @ ?cc2),s3,stk)
    using Seq.IH(2) by fastforce
  ultimately show ?case by simp (blast intro: star_trans)
next
  case (WhileTrue b s1 c s2 s3)
  let ?cc = ccomp c
  let ?cb = bcomp b False (size ?cc + 1)
```

```
  let ?cw = ccomp(WHILE b DO c)
  have ?cw ⊢ (0,s1,stk) →* (size ?cb,s1,stk)
    using ‹bval b s1› by fastforce
  moreover
  have ?cw ⊢ (size ?cb,s1,stk) →* (size ?cb + size ?cc,s2,stk)
    using WhileTrue.IH(1) by fastforce
  moreover
  have ?cw ⊢ (size ?cb + size ?cc,s2,stk) →* (0,s2,stk)
    by fastforce
  moreover
  have ?cw ⊢ (0,s2,stk) →* (size ?cw,s3,stk) by(rule WhileTrue.IH(2))
  ultimately show ?case by(blast intro: star_trans)
qed fastforce+

end
```

7 Compiler Correctness, Reverse Direction

```
theory Compiler2
imports Compiler
begin
```

The preservation of the source code semantics is already shown in the parent theory *Compiler*. This here shows the second direction.

7.1 Definitions

Execution in *n* steps for simpler induction

```
primrec
  exec_n :: instr list ⇒ config ⇒ nat ⇒ config ⇒ bool
  (_/ ⊢ (_ →^_/ _) [65,0,1000,55] 55)
where
  P ⊢ c →^0 c' = (c'=c) |
  P ⊢ c →^(Suc n) c'' = (∃ c'. (P ⊢ c → c') ∧ P ⊢ c' →^n c'')
```

The possible successor PCs of an instruction at position *n*

```
definition isuccs :: instr ⇒ int ⇒ int set where
isuccs i n = (case i of
  JMP j ⇒ {n + 1 + j} |
  JMPLESS j ⇒ {n + 1 + j, n + 1} |
  JMPGE j ⇒ {n + 1 + j, n + 1} |
  _ ⇒ {n +1})
```

The possible successors PCs of an instruction list

definition *succs* :: *instr list* ⇒ *int* ⇒ *int set* **where**
$succs\ P\ n = \{s.\ \exists\, i{::}int.\ 0 \le i \wedge i < size\ P \wedge s \in isuccs\ (P!!i)\ (n+i)\}$

Possible exit PCs of a program

definition *exits* :: *instr list* ⇒ *int set* **where**
$exits\ P = succs\ P\ 0 - \{0..< size\ P\}$

7.2 Basic properties of *exec_n*

lemma *exec_n_exec*:
$P \vdash c \rightarrow\hat{}n\ c' \Longrightarrow P \vdash c \rightarrow\!*\ c'$
by (*induct n arbitrary*: *c*) (*auto intro*: *star.step*)

lemma *exec_0* [*intro*!]: $P \vdash c \rightarrow\hat{}0\ c$ **by** *simp*

lemma *exec_Suc*:
$[\![\ P \vdash c \rightarrow c';\ P \vdash c' \rightarrow\hat{}n\ c''\ ]\!] \Longrightarrow P \vdash c \rightarrow\hat{}(Suc\ n)\ c''$
by (*fastforce simp del*: *split_paired_Ex*)

lemma *exec_exec_n*:
$P \vdash c \rightarrow\!*\ c' \Longrightarrow \exists\, n.\ P \vdash c \rightarrow\hat{}n\ c'$
by (*induct rule*: *star.induct*) (*auto intro*: *exec_Suc*)

lemma *exec_eq_exec_n*:
$(P \vdash c \rightarrow\!*\ c') = (\exists\, n.\ P \vdash c \rightarrow\hat{}n\ c')$
by (*blast intro*: *exec_exec_n exec_n_exec*)

lemma *exec_n_Nil* [*simp*]:
$[] \vdash c \rightarrow\hat{}k\ c' = (c' = c \wedge k = 0)$
by (*induct k*) (*auto simp*: *exec1_def*)

lemma *exec1_exec_n* [*intro*!]:
$P \vdash c \rightarrow c' \Longrightarrow P \vdash c \rightarrow\hat{}1\ c'$
by (*cases c'*) *simp*

7.3 Concrete symbolic execution steps

lemma *exec_n_step*:
$n \neq n' \Longrightarrow$
$P \vdash (n,stk,s) \rightarrow\hat{}k\ (n',stk',s') =$
$(\exists\, c.\ P \vdash (n,stk,s) \rightarrow c \wedge P \vdash c \rightarrow\hat{}(k - 1)\ (n',stk',s') \wedge 0 < k)$
by (*cases k*) *auto*

```
lemma exec1_end:
  size P <= fst c ⟹ ¬ P ⊢ c → c′
  by (auto simp: exec1_def)

lemma exec_n_end:
  size P <= (n::int) ⟹
  P ⊢ (n,s,stk) →^k (n′,s′,stk′) = (n′ = n ∧ stk′=stk ∧ s′=s ∧ k =0)
  by (cases k) (auto simp: exec1_end)

lemmas exec_n_simps = exec_n_step exec_n_end
```

7.4 Basic properties of *succs*

```
lemma succs_simps [simp]:
  succs [ADD] n = {n + 1}
  succs [LOADI v] n = {n + 1}
  succs [LOAD x] n = {n + 1}
  succs [STORE x] n = {n + 1}
  succs [JMP i] n = {n + 1 + i}
  succs [JMPGE i] n = {n + 1 + i, n + 1}
  succs [JMPLESS i] n = {n + 1 + i, n + 1}
  by (auto simp: succs_def isuccs_def)

lemma succs_empty [iff]: succs [] n = {}
  by (simp add: succs_def)

lemma succs_Cons:
  succs (x#xs) n = isuccs x n ∪ succs xs (1+n) (is _ = ?x ∪ ?xs)
proof
  let ?isuccs = λp P n i::int. 0 ≤ i ∧ i < size P ∧ p ∈ isuccs (P!!i) (n+i)
  have p ∈ ?x ∪ ?xs if assm: p ∈ succs (x#xs) n for p
  proof −
    from assm obtain i::int where isuccs: ?isuccs p (x#xs) n i
      unfolding succs_def by auto
    show ?thesis
    proof cases
      assume i = 0 with isuccs show ?thesis by simp
    next
      assume i ≠ 0
      with isuccs
      have ?isuccs p xs (1+n) (i − 1) by auto
      hence p ∈ ?xs unfolding succs_def by blast
      thus ?thesis ..
    qed
```

```
  qed
  thus succs (x#xs) n ⊆ ?x ∪ ?xs ..

  have p ∈ succs (x#xs) n if assm: p ∈ ?x ∨ p ∈ ?xs for p
  proof –
    from assm show ?thesis
    proof
      assume p ∈ ?x thus ?thesis by (fastforce simp: succs_def)
    next
      assume p ∈ ?xs
      then obtain i where ?isuccs p xs (1+n) i
        unfolding succs_def by auto
      hence ?isuccs p (x#xs) n (1+i)
        by (simp add: algebra_simps)
      thus ?thesis unfolding succs_def by blast
    qed
  qed
  thus ?x ∪ ?xs ⊆ succs (x#xs) n by blast
qed

lemma succs_iexec1:
  assumes c′ = iexec (P!!i) (i,s,stk) 0 ≤ i i < size P
  shows fst c′ ∈ succs P 0
  using assms by (auto simp: succs_def isuccs_def split: instr.split)

lemma succs_shift:
  (p − n ∈ succs P 0) = (p ∈ succs P n)
  by (fastforce simp: succs_def isuccs_def split: instr.split)

lemma inj_op_plus [simp]:
  inj ((+) (i::int))
  by (metis add_minus_cancel inj_on_inverseI)

lemma succs_set_shift [simp]:
  (+) i ‘ succs xs 0 = succs xs i
  by (force simp: succs_shift [where n=i, symmetric] intro: set_eqI)

lemma succs_append [simp]:
  succs (xs @ ys) n = succs xs n ∪ succs ys (n + size xs)
  by (induct xs arbitrary: n) (auto simp: succs_Cons algebra_simps)


lemma exits_append [simp]:
  exits (xs @ ys) = exits xs ∪ ((+) (size xs)) ‘ exits ys −
```

{0..<size xs + size ys}
by (*auto simp*: *exits_def image_set_diff*)

lemma *exits_single*:
exits [*x*] = *isuccs x 0* − {*0*}
by (*auto simp*: *exits_def succs_def*)

lemma *exits_Cons*:
exits (*x* # *xs*) = (*isuccs x 0* − {*0*}) ∪ ((+) *1*) ‘ *exits xs* −
{*0*..<*1* + *size xs*}
using *exits_append* [*of* [*x*] *xs*]
by (*simp add*: *exits_single*)

lemma *exits_empty* [*iff*]: *exits* [] = {} **by** (*simp add*: *exits_def*)

lemma *exits_simps* [*simp*]:
exits [*ADD*] = {*1*}
exits [*LOADI v*] = {*1*}
exits [*LOAD x*] = {*1*}
exits [*STORE x*] = {*1*}
$i \neq -1 \Longrightarrow$ *exits* [*JMP i*] = {*1* + *i*}
$i \neq -1 \Longrightarrow$ *exits* [*JMPGE i*] = {*1* + *i*, *1*}
$i \neq -1 \Longrightarrow$ *exits* [*JMPLESS i*] = {*1* + *i*, *1*}
by (*auto simp*: *exits_def*)

lemma *acomp_succs* [*simp*]:
succs (*acomp a*) *n* = {*n* + *1* .. *n* + *size* (*acomp a*)}
by (*induct a arbitrary*: *n*) *auto*

lemma *acomp_size*:
(*1*::*int*) ≤ *size* (*acomp a*)
by (*induct a*) *auto*

lemma *acomp_exits* [*simp*]:
exits (*acomp a*) = {*size* (*acomp a*)}
by (*auto simp*: *exits_def acomp_size*)

lemma *bcomp_succs*:
0 ≤ *i* ⟹
succs (*bcomp b f i*) *n* ⊆ {*n* .. *n* + *size* (*bcomp b f i*)}
∪ {*n* + *i* + *size* (*bcomp b f i*)}
proof (*induction b arbitrary*: *f i n*)
case (*And b1 b2*)
from *And.prems*

```
  show ?case
    by (cases f)
       (auto dest: And.IH(1) [THEN subsetD, rotated]
                   And.IH(2) [THEN subsetD, rotated])
qed auto

lemmas bcomp_succsD [dest!] = bcomp_succs [THEN subsetD, rotated]

lemma bcomp_exits:
  fixes i :: int
  shows
  0 ≤ i ⟹
  exits (bcomp b f i) ⊆ {size (bcomp b f i), i + size (bcomp b f i)}
  by (auto simp: exits_def)

lemma bcomp_exitsD [dest!]:
  p ∈ exits (bcomp b f i) ⟹ 0 ≤ i ⟹
  p = size (bcomp b f i) ∨ p = i + size (bcomp b f i)
  using bcomp_exits by auto

lemma ccomp_succs:
  succs (ccomp c) n ⊆ {n..n + size (ccomp c)}
proof (induction c arbitrary: n)
  case SKIP thus ?case by simp
next
  case Assign thus ?case by simp
next
  case (Seq c1 c2)
  from Seq.prems
  show ?case
    by (fastforce dest: Seq.IH [THEN subsetD])
next
  case (If b c1 c2)
  from If.prems
  show ?case
    by (auto dest!: If.IH [THEN subsetD] simp: isuccs_def succs_Cons)
next
  case (While b c)
  from While.prems
  show ?case by (auto dest!: While.IH [THEN subsetD])
qed

lemma ccomp_exits:
  exits (ccomp c) ⊆ {size (ccomp c)}
```

using *ccomp_succs* [*of c 0*] **by** (*auto simp*: *exits_def*)

lemma *ccomp_exitsD* [*dest*!]:
$p \in exits\ (ccomp\ c) \Longrightarrow p = size\ (ccomp\ c)$
using *ccomp_exits* **by** *auto*

7.5 Splitting up machine executions

lemma *exec1_split*:
fixes *i j* :: *int*
shows
$P\ @\ c\ @\ P' \vdash (size\ P + i,\ s) \to (j,s') \Longrightarrow 0 \le i \Longrightarrow i < size\ c \Longrightarrow$
$c \vdash (i,s) \to (j - size\ P,\ s')$
by (*auto split*: *instr.splits simp*: *exec1_def*)

lemma *exec_n_split*:
fixes *i j* :: *int*
assumes $P\ @\ c\ @\ P' \vdash (size\ P + i,\ s) \to\hat{}n\ (j,\ s')$
$0 \le i\ i < size\ c$
$j \notin \{size\ P\ ..< size\ P + size\ c\}$
shows $\exists\, s''\ (i'::int)\ k\ m.$
$c \vdash (i,\ s) \to\hat{}k\ (i',\ s'') \wedge$
$i' \in exits\ c \wedge$
$P\ @\ c\ @\ P' \vdash (size\ P + i',\ s'') \to\hat{}m\ (j,\ s') \wedge$
$n = k + m$
using *assms* **proof** (*induction n arbitrary*: *i j s*)
case *0*
thus *?case* **by** *simp*
next
case (*Suc n*)
have *i*: $0 \le i\ i < size\ c$ **by** *fact+*
from *Suc.prems*
have *j*: $\neg\ (size\ P \le j \wedge j < size\ P + size\ c)$ **by** *simp*
from *Suc.prems*
obtain *i0 s0* **where**
step: $P\ @\ c\ @\ P' \vdash (size\ P + i,\ s) \to (i0,s0)$ **and**
rest: $P\ @\ c\ @\ P' \vdash (i0,s0) \to\hat{}n\ (j,\ s')$
by *clarsimp*

from *step i*
have *c*: $c \vdash (i,s) \to (i0 - size\ P,\ s0)$ **by** (*rule exec1_split*)

have $i0 = size\ P + (i0 - size\ P)$ **by** *simp*
then obtain *j0*::*int* **where** *j0*: $i0 = size\ P + j0$ **..**

note *split_paired_Ex* [*simp del*]

have *?case* **if** *assm*: $j0 \in \{0 ..< size\ c\}$
proof −
from *assm j0 j rest c* **show** *?case*
by (*fastforce dest*!: *Suc.IH intro*!: *exec_Suc*)
qed
moreover
have *?case* **if** *assm*: $j0 \notin \{0 ..< size\ c\}$
proof −
from *c j0* **have** $j0 \in succs\ c\ 0$
by (*auto dest*: *succs_iexec1 simp*: *exec1_def simp del*: *iexec.simps*)
with *assm* **have** $j0 \in exits\ c$ **by** (*simp add*: *exits_def*)
with *c j0 rest* **show** *?case* **by** *fastforce*
qed
ultimately
show *?case* **by** *cases*
qed

lemma *exec_n_drop_right*:
fixes *j* :: *int*
assumes $c\ @\ P' \vdash (0, s) \rightarrow\hat{}n\ (j, s')$ $j \notin \{0..<size\ c\}$
shows $\exists s''\ i'\ k\ m.$
$(if\ c = []\ then\ s'' = s \wedge i' = 0 \wedge k = 0$
$else\ c \vdash (0, s) \rightarrow\hat{}k\ (i', s'') \wedge$
$i' \in exits\ c) \wedge$
$c\ @\ P' \vdash (i', s'') \rightarrow\hat{}m\ (j, s') \wedge$
$n = k + m$
using *assms*
by (*cases* $c = []$)
(*auto dest*: *exec_n_split* [**where** *P*=[], *simplified*])

Dropping the left context of a potentially incomplete execution of *c*.

lemma *exec1_drop_left*:
fixes *i n* :: *int*
assumes $P1\ @\ P2 \vdash (i, s, stk) \rightarrow (n, s', stk')$ **and** $size\ P1 \le i$
shows $P2 \vdash (i - size\ P1, s, stk) \rightarrow (n - size\ P1, s', stk')$
proof −
have $i = size\ P1 + (i - size\ P1)$ **by** *simp*
then obtain *i'* :: *int* **where** $i = size\ P1 + i'$ **..**
moreover
have $n = size\ P1 + (n - size\ P1)$ **by** *simp*
then obtain *n'* :: *int* **where** $n = size\ P1 + n'$ **..**

```
  ultimately
  show ?thesis using assms
    by (clarsimp simp: exec1_def simp del: iexec.simps)
qed

lemma exec_n_drop_left:
  fixes i n :: int
  assumes P @ P' ⊢ (i, s, stk) →^k (n, s', stk')
          size P ≤ i exits P' ⊆ {0..}
  shows P' ⊢ (i − size P, s, stk) →^k (n − size P, s', stk')
using assms proof (induction k arbitrary: i s stk)
  case 0 thus ?case by simp
next
  case (Suc k)
  from Suc.prems
  obtain i' s'' stk'' where
    step: P @ P' ⊢ (i, s, stk) → (i', s'', stk'') and
    rest: P @ P' ⊢ (i', s'', stk'') →^k (n, s', stk')
    by auto
  from step ‹size P ≤ i›
  have *: P' ⊢ (i − size P, s, stk) → (i' − size P, s'', stk'')
    by (rule exec1_drop_left)
  then have i' − size P ∈ succs P' 0
    by (fastforce dest!: succs_iexec1 simp: exec1_def simp del: iexec.simps)
  with ‹exits P' ⊆ {0..}›
  have size P ≤ i' by (auto simp: exits_def)
  from rest this ‹exits P' ⊆ {0..}›
  have P' ⊢ (i' − size P, s'', stk'') →^k (n − size P, s', stk')
    by (rule Suc.IH)
  with * show ?case by auto
qed

lemmas exec_n_drop_Cons =
  exec_n_drop_left [where P=[instr], simplified] for instr

definition
  closed P ⟷ exits P ⊆ {size P}

lemma ccomp_closed [simp, intro!]: closed (ccomp c)
  using ccomp_exits by (auto simp: closed_def)

lemma acomp_closed [simp, intro!]: closed (acomp c)
  by (simp add: closed_def)
```

```
lemma exec_n_split_full:
  fixes j :: int
  assumes exec: P @ P' ⊢ (0,s,stk) →^k (j, s', stk')
  assumes P: size P ≤ j
  assumes closed: closed P
  assumes exits: exits P' ⊆ {0..}
  shows ∃ k1 k2 s'' stk''. P ⊢ (0,s,stk) →^k1 (size P, s'', stk'') ∧
                           P' ⊢ (0,s'',stk'') →^k2 (j − size P, s', stk')
proof (cases P)
  case Nil with exec
  show ?thesis by fastforce
next
  case Cons
  hence 0 < size P by simp
  with exec P closed
  obtain k1 k2 s'' stk'' where
    1: P ⊢ (0,s,stk) →^k1 (size P, s'', stk'') and
    2: P @ P' ⊢ (size P,s'',stk'') →^k2 (j, s', stk')
    by (auto dest!: exec_n_split [where P=[] and i=0, simplified]
             simp: closed_def)
  moreover
  have j = size P + (j − size P) by simp
  then obtain j0 :: int where j = size P + j0 ..
  ultimately
  show ?thesis using exits
    by (fastforce dest: exec_n_drop_left)
qed
```

7.6 Correctness theorem

```
lemma acomp_neq_Nil [simp]:
  acomp a ≠ []
  by (induct a) auto

lemma acomp_exec_n [dest!]:
  acomp a ⊢ (0,s,stk) →^n (size (acomp a),s',stk') ⟹
  s' = s ∧ stk' = aval a s#stk
proof (induction a arbitrary: n s' stk stk')
  case (Plus a1 a2)
  let ?sz = size (acomp a1) + (size (acomp a2) + 1)
  from Plus.prems
  have acomp a1 @ acomp a2 @ [ADD] ⊢ (0,s,stk) →^n (?sz, s', stk')
    by (simp add: algebra_simps)
```

```
  then obtain n1 s1 stk1 n2 s2 stk2 n3 where
    acomp a1 ⊢ (0,s,stk) →^n1 (size (acomp a1), s1, stk1)
    acomp a2 ⊢ (0,s1,stk1) →^n2 (size (acomp a2), s2, stk2)
      [ADD] ⊢ (0,s2,stk2) →^n3 (1, s', stk')
    by (auto dest!: exec_n_split_full)

  thus ?case by (fastforce dest: Plus.IH simp: exec_n_simps exec1_def)
qed (auto simp: exec_n_simps exec1_def)

lemma bcomp_split:
  fixes i j :: int
  assumes bcomp b f i @ P' ⊢ (0, s, stk) →^n (j, s', stk')
          j ∉ {0..<size (bcomp b f i)} 0 ≤ i
  shows ∃ s'' stk'' (i'::int) k m.
           bcomp b f i ⊢ (0, s, stk) →^k (i', s'', stk'') ∧
           (i' = size (bcomp b f i) ∨ i' = i + size (bcomp b f i)) ∧
           bcomp b f i @ P' ⊢ (i', s'', stk'') →^m (j, s', stk') ∧
           n = k + m
 using assms by (cases bcomp b f i = []) (fastforce dest!: exec_n_drop_right)+

lemma bcomp_exec_n [dest]:
  fixes i j :: int
  assumes bcomp b f j ⊢ (0, s, stk) →^n (i, s', stk')
          size (bcomp b f j) ≤ i 0 ≤ j
  shows i = size(bcomp b f j) + (if f = bval b s then j else 0) ∧
         s' = s ∧ stk' = stk
using assms proof (induction b arbitrary: f j i n s' stk')
  case Bc thus ?case
    by (simp split: if_split_asm add: exec_n_simps exec1_def)
next
  case (Not b)
  from Not.prems show ?case
    by (fastforce dest!: Not.IH)
next
  case (And b1 b2)

  let ?b2 = bcomp b2 f j
  let ?m  = if f then size ?b2 else size ?b2 + j
  let ?b1 = bcomp b1 False ?m

  have j: size (bcomp (And b1 b2) f j) ≤ i 0 ≤ j by fact+

  from And.prems
  obtain s'' stk'' and i'::int and k m where
```

b1: *?b1* ⊢ (*0*, *s*, *stk*) →^*k* (*i'*, *s''*, *stk''*)
i' = *size ?b1* ∨ *i'* = *?m* + *size ?b1* **and**
b2: *?b2* ⊢ (*i'* − *size ?b1*, *s''*, *stk''*) →^*m* (*i* − *size ?b1*, *s'*, *stk'*)
by (*auto dest*!: *bcomp_split dest*: *exec_n_drop_left*)
from *b1 j*
have *i'* = *size ?b1* + (*if* ¬*bval b1 s then ?m else 0*) ∧ *s''* = *s* ∧ *stk''* = *stk*
by (*auto dest*!: *And.IH*)
with *b2 j*
show *?case*
by (*fastforce dest*!: *And.IH simp*: *exec_n_end split*: *if_split_asm*)
next
case *Less*
thus *?case* **by** (*auto dest*!: *exec_n_split_full simp*: *exec_n_simps exec1_def*)

qed

lemma *ccomp_empty* [*elim*!]:
ccomp c = [] ⟹ (*c*,*s*) ⇒ *s*
by (*induct c*) *auto*

declare *assign_simp* [*simp*]

lemma *ccomp_exec_n*:
ccomp c ⊢ (*0*,*s*,*stk*) →^*n* (*size*(*ccomp c*),*t*,*stk'*)
⟹ (*c*,*s*) ⇒ *t* ∧ *stk'*=*stk*
proof (*induction c arbitrary*: *s t stk stk' n*)
case *SKIP*
thus *?case* **by** *auto*
next
case (*Assign x a*)
thus *?case*
by *simp* (*fastforce dest*!: *exec_n_split_full simp*: *exec_n_simps exec1_def*)
next
case (*Seq c1 c2*)
thus *?case* **by** (*fastforce dest*!: *exec_n_split_full*)
next
case (*If b c1 c2*)
note *If.IH* [*dest*!]

let *?if* = *IF b THEN c1 ELSE c2*
let *?cs* = *ccomp ?if*
let *?bcomp* = *bcomp b False* (*size* (*ccomp c1*) + *1*)

```
  from ‹?cs ⊢ (0,s,stk) →^n (size ?cs,t,stk')›
  obtain i' :: int and k m s'' stk'' where
    cs: ?cs ⊢ (i',s'',stk'') →^m (size ?cs,t,stk') and
        ?bcomp ⊢ (0,s,stk) →^k (i', s'', stk'')
        i' = size ?bcomp ∨ i' = size ?bcomp + size (ccomp c1) + 1
    by (auto dest!: bcomp_split)

  hence i':
    s''=s stk'' = stk
    i' = (if bval b s then size ?bcomp else size ?bcomp+size(ccomp c1)+1)
    by auto

  with cs have cs':
    ccomp c1@JMP (size (ccomp c2))#ccomp c2 ⊢
       (if bval b s then 0 else size (ccomp c1)+1, s, stk) →^m
       (1 + size (ccomp c1) + size (ccomp c2), t, stk')
    by (fastforce dest: exec_n_drop_left simp: exits_Cons isuccs_def alge-
bra_simps)

  show ?case
  proof (cases bval b s)
    case True with cs'
    show ?thesis
      by simp
         (fastforce dest: exec_n_drop_right
                    split: if_split_asm
                    simp: exec_n_simps exec1_def)
  next
    case False with cs'
    show ?thesis
      by (auto dest!: exec_n_drop_Cons exec_n_drop_left
               simp: exits_Cons isuccs_def)
  qed
next
  case (While b c)

  from While.prems
  show ?case
  proof (induction n arbitrary: s rule: nat_less_induct)
    case (1 n)

    have ?case if assm: ¬ bval b s
    proof −
      from assm 1.prems
```

```
  show ?case
    by simp (fastforce dest!: bcomp_split simp: exec_n_simps)
qed
moreover
have ?case if b: bval b s
proof -
  let ?c0 = WHILE b DO c
  let ?cs = ccomp ?c0
  let ?bs = bcomp b False (size (ccomp c) + 1)
  let ?jmp = [JMP (-((size ?bs + size (ccomp c) + 1)))]

  from 1.prems b
  obtain k where
    cs: ?cs ⊢ (size ?bs, s, stk) →^k (size ?cs, t, stk') and
    k:  k ≤ n
    by (fastforce dest!: bcomp_split)

  show ?case
  proof cases
    assume ccomp c = []
    with cs k
    obtain m where
      ?cs ⊢ (0,s,stk) →^m (size (ccomp ?c0), t, stk')
      m < n
      by (auto simp: exec_n_step [where k=k] exec1_def)
    with 1.IH
    show ?case by blast
  next
    assume ccomp c ≠ []
    with cs
    obtain m m' s'' stk'' where
      c: ccomp c ⊢ (0, s, stk) →^m' (size (ccomp c), s'', stk'') and
      rest: ?cs ⊢ (size ?bs + size (ccomp c), s'', stk'') →^m
                 (size ?cs, t, stk') and
      m: k = m + m'
      by (auto dest: exec_n_split [where i=0, simplified])
    from c
    have (c,s) ⇒ s'' and stk: stk'' = stk
      by (auto dest!: While.IH)
    moreover
    from rest m k stk
    obtain k' where
      ?cs ⊢ (0, s'', stk) →^k' (size ?cs, t, stk')
      k' < n
```

```
      by (auto simp: exec_n_step [where k=m] exec1_def)
    with 1.IH
    have (?c0, s'') ⇒ t ∧ stk' = stk by blast
    ultimately
    show ?case using b by blast
  qed
 qed
 ultimately show ?case by cases
 qed
qed

theorem ccomp_exec:
  ccomp c ⊢ (0,s,stk) →* (size(ccomp c),t,stk') ⟹ (c,s) ⇒ t
  by (auto dest: exec_exec_n ccomp_exec_n)

corollary ccomp_sound:
  ccomp c ⊢ (0,s,stk) →* (size(ccomp c),t,stk)  ⟷  (c,s) ⇒ t
  by (blast intro!: ccomp_exec ccomp_bigstep)

end
```

8 A Typed Language

theory *Types* **imports** *Star Complex_Main* **begin**

We build on *Complex_Main* instead of *Main* to access the real numbers.

8.1 Arithmetic Expressions

```
datatype val = Iv int | Rv real

type_synonym vname = string
type_synonym state = vname ⇒ valdatatype aexp =  Ic int | Rc real |
V vname | Plus aexp aexp

inductive taval :: aexp ⇒ state ⇒ val ⇒ bool where
taval (Ic i) s (Iv i) |
taval (Rc r) s (Rv r) |
taval (V x) s (s x) |
taval a1 s (Iv i1) ⟹ taval a2 s (Iv i2)
 ⟹ taval (Plus a1 a2) s (Iv(i1+i2)) |
taval a1 s (Rv r1) ⟹ taval a2 s (Rv r2)
 ⟹ taval (Plus a1 a2) s (Rv(r1+r2))
```

inductive_cases [*elim*!]:
taval (*Ic i*) *s v* *taval* (*Rc i*) *s v*
taval (*V x*) *s v*
taval (*Plus a1 a2*) *s v*

8.2 Boolean Expressions

datatype *bexp* = *Bc bool* | *Not bexp* | *And bexp bexp* | *Less aexp aexp*

inductive *tbval* :: *bexp* ⇒ *state* ⇒ *bool* ⇒ *bool* **where**
tbval (*Bc v*) *s v* |
tbval b s bv ⟹ *tbval* (*Not b*) *s* (¬ *bv*) |
tbval b1 s bv1 ⟹ *tbval b2 s bv2* ⟹ *tbval* (*And b1 b2*) *s* (*bv1* & *bv2*) |
taval a1 s (*Iv i1*) ⟹ *taval a2 s* (*Iv i2*) ⟹ *tbval* (*Less a1 a2*) *s* (*i1* < *i2*)
|
taval a1 s (*Rv r1*) ⟹ *taval a2 s* (*Rv r2*) ⟹ *tbval* (*Less a1 a2*) *s* (*r1* <
r2)

8.3 Syntax of Commands

datatype
com = *SKIP*
| *Assign vname aexp* (_ ::= _ [*1000*, *61*] *61*)
| *Seq com com* (_;; _ [*60*, *61*] *60*)
| *If bexp com com* (*IF* _ *THEN* _ *ELSE* _ [*0*, *0*, *61*] *61*)
| *While bexp com* (*WHILE* _ *DO* _ [*0*, *61*] *61*)

8.4 Small-Step Semantics of Commands

inductive
small_step :: (*com* × *state*) ⇒ (*com* × *state*) ⇒ *bool* (**infix** → *55*)
where
Assign: *taval a s v* ⟹ (*x* ::= *a*, *s*) → (*SKIP*, *s*(*x* := *v*)) |

Seq1: (*SKIP*;;*c*,*s*) → (*c*,*s*) |
Seq2: (*c1*,*s*) → (*c1'*,*s'*) ⟹ (*c1*;;*c2*,*s*) → (*c1'*;;*c2*,*s'*) |

IfTrue: *tbval b s True* ⟹ (*IF b THEN c1 ELSE c2*,*s*) → (*c1*,*s*) |
IfFalse: *tbval b s False* ⟹ (*IF b THEN c1 ELSE c2*,*s*) → (*c2*,*s*) |

While: (*WHILE b DO c*,*s*) → (*IF b THEN c*;; *WHILE b DO c ELSE*
SKIP,*s*)

lemmas *small_step_induct* = *small_step.induct*[*split_format*(*complete*)]

8.5 The Type System

datatype *ty* = *Ity* | *Rty*

type_synonym *tyenv* = *vname* ⇒ *ty*

inductive *atyping* :: *tyenv* ⇒ *aexp* ⇒ *ty* ⇒ *bool*
((*1*_/ ⊢/ (_ :/ _)) [*50*,*0*,*50*] *50*)
where
Ic_ty: Γ ⊢ *Ic i* : *Ity* |
Rc_ty: Γ ⊢ *Rc r* : *Rty* |
V_ty: Γ ⊢ *V x* : Γ *x* |
Plus_ty: Γ ⊢ *a1* : τ ⟹ Γ ⊢ *a2* : τ ⟹ Γ ⊢ *Plus a1 a2* : τ

declare *atyping.intros* [*intro*!]
inductive_cases [*elim*!]:
Γ ⊢ *V x* : τ Γ ⊢ *Ic i* : τ Γ ⊢ *Rc r* : τ Γ ⊢ *Plus a1 a2* : τ

Warning: the ":" notation leads to syntactic ambiguities, i.e. multiple parse trees, because ":" also stands for set membership. In most situations Isabelle's type system will reject all but one parse tree, but will still inform you of the potential ambiguity.

inductive *btyping* :: *tyenv* ⇒ *bexp* ⇒ *bool* (**infix** ⊢ *50*)
where
B_ty: Γ ⊢ *Bc v* |
Not_ty: Γ ⊢ *b* ⟹ Γ ⊢ *Not b* |
And_ty: Γ ⊢ *b1* ⟹ Γ ⊢ *b2* ⟹ Γ ⊢ *And b1 b2* |
Less_ty: Γ ⊢ *a1* : τ ⟹ Γ ⊢ *a2* : τ ⟹ Γ ⊢ *Less a1 a2*

declare *btyping.intros* [*intro*!]
inductive_cases [*elim*!]: Γ ⊢ *Not b* Γ ⊢ *And b1 b2* Γ ⊢ *Less a1 a2*

inductive *ctyping* :: *tyenv* ⇒ *com* ⇒ *bool* (**infix** ⊢ *50*) **where**
Skip_ty: Γ ⊢ *SKIP* |
Assign_ty: Γ ⊢ *a* : Γ(*x*) ⟹ Γ ⊢ *x* ::= *a* |
Seq_ty: Γ ⊢ *c1* ⟹ Γ ⊢ *c2* ⟹ Γ ⊢ *c1*;;*c2* |
If_ty: Γ ⊢ *b* ⟹ Γ ⊢ *c1* ⟹ Γ ⊢ *c2* ⟹ Γ ⊢ *IF b THEN c1 ELSE c2* |
While_ty: Γ ⊢ *b* ⟹ Γ ⊢ *c* ⟹ Γ ⊢ *WHILE b DO c*

declare *ctyping.intros* [*intro*!]
inductive_cases [*elim*!]:
Γ ⊢ *x* ::= *a* Γ ⊢ *c1*;;*c2*
Γ ⊢ *IF b THEN c1 ELSE c2*
Γ ⊢ *WHILE b DO c*

8.6 Well-typed Programs Do Not Get Stuck

```
fun type :: val ⇒ ty where
type (Iv i) = Ity |
type (Rv r) = Rty

lemma type_eq_Ity[simp]: type v = Ity ⟷ (∃ i. v = Iv i)
by (cases v) simp_all

lemma type_eq_Rty[simp]: type v = Rty ⟷ (∃ r. v = Rv r)
by (cases v) simp_all

definition styping :: tyenv ⇒ state ⇒ bool (infix ⊢ 50)
where Γ ⊢ s  ⟷  (∀ x. type (s x) = Γ x)

lemma apreservation:
  Γ ⊢ a : τ ⟹ taval a s v ⟹ Γ ⊢ s ⟹ type v = τ
apply(induction arbitrary: v rule: atyping.induct)
apply (fastforce simp: styping_def)+
done

lemma aprogress: Γ ⊢ a : τ ⟹ Γ ⊢ s ⟹ ∃ v. taval a s v
proof(induction rule: atyping.induct)
  case (Plus_ty Γ a1 t a2)
  then obtain v1 v2 where v: taval a1 s v1 taval a2 s v2 by blast
  show ?case
  proof (cases v1)
    case Iv
    with Plus_ty v show ?thesis
      by(fastforce intro: taval.intros(4) dest!: apreservation)
  next
    case Rv
    with Plus_ty v show ?thesis
      by(fastforce intro: taval.intros(5) dest!: apreservation)
  qed
qed (auto intro: taval.intros)

lemma bprogress: Γ ⊢ b ⟹ Γ ⊢ s ⟹ ∃ v. tbval b s v
proof(induction rule: btyping.induct)
  case (Less_ty Γ a1 t a2)
  then obtain v1 v2 where v: taval a1 s v1 taval a2 s v2
    by (metis aprogress)
  show ?case
  proof (cases v1)
```

```
    case Iv
    with Less_ty v show ?thesis
      by (fastforce intro!: tbval.intros(4) dest!:apreservation)
  next
    case Rv
    with Less_ty v show ?thesis
      by (fastforce intro!: tbval.intros(5) dest!:apreservation)
  qed
qed (auto intro: tbval.intros)

theorem progress:
  Γ ⊢ c ⟹ Γ ⊢ s ⟹ c ≠ SKIP ⟹ ∃ cs'. (c,s) → cs'
proof(induction rule: ctyping.induct)
  case Skip_ty thus ?case by simp
next
  case Assign_ty
  thus ?case by (metis Assign aprogress)
next
  case Seq_ty thus ?case by simp (metis Seq1 Seq2)
next
  case (If_ty Γ b c1 c2)
  then obtain bv where tbval b s bv by (metis bprogress)
  show ?case
  proof(cases bv)
    assume bv
    with ‹tbval b s bv› show ?case by simp (metis IfTrue)
  next
    assume ¬bv
    with ‹tbval b s bv› show ?case by simp (metis IfFalse)
  qed
next
  case While_ty show ?case by (metis While)
qed

theorem styping_preservation:
  (c,s) → (c',s') ⟹ Γ ⊢ c ⟹ Γ ⊢ s ⟹ Γ ⊢ s'
proof(induction rule: small_step_induct)
  case Assign thus ?case
    by (auto simp: styping_def) (metis Assign(1,3) apreservation)
qed auto

theorem ctyping_preservation:
  (c,s) → (c',s') ⟹ Γ ⊢ c ⟹ Γ ⊢ c'
by (induct rule: small_step_induct) (auto simp: ctyping.intros)
```

abbreviation *small_steps* :: *com* ∗ *state* ⇒ *com* ∗ *state* ⇒ *bool* (**infix** →∗ *55*)
where *x* →∗ *y* == *star small_step x y*

theorem *type_sound*:
$(c,s) \to* (c',s') \Longrightarrow \Gamma \vdash c \Longrightarrow \Gamma \vdash s \Longrightarrow c' \neq SKIP$
$\Longrightarrow \exists cs''. (c',s') \to cs''$
apply(*induction rule*:*star_induct*)
apply (*metis progress*)
by (*metis styping_preservation ctyping_preservation*)

end

8.7 Type Variables

theory *Poly_Types* **imports** *Types* **begin**

datatype *ty* = *Ity* | *Rty* | *TV nat*

Everything else remains the same.

type_synonym *tyenv* = *vname* ⇒ *ty*

inductive *atyping* :: *tyenv* ⇒ *aexp* ⇒ *ty* ⇒ *bool*
((*1_*/ ⊢*p*/ (_ :/ _)) [*50*,*0*,*50*] *50*)
where
Γ ⊢*p* *Ic i* : *Ity* |
Γ ⊢*p* *Rc r* : *Rty* |
Γ ⊢*p* *V x* : Γ *x* |
$\Gamma \vdash p\ a1 : \tau \Longrightarrow \Gamma \vdash p\ a2 : \tau \Longrightarrow \Gamma \vdash p\ Plus\ a1\ a2 : \tau$

inductive *btyping* :: *tyenv* ⇒ *bexp* ⇒ *bool* (**infix** ⊢*p* *50*)
where
Γ ⊢*p* *Bc v* |
Γ ⊢*p* *b* ⟹ Γ ⊢*p* *Not b* |
Γ ⊢*p* *b1* ⟹ Γ ⊢*p* *b2* ⟹ Γ ⊢*p* *And b1 b2* |
$\Gamma \vdash p\ a1 : \tau \Longrightarrow \Gamma \vdash p\ a2 : \tau \Longrightarrow \Gamma \vdash p\ Less\ a1\ a2$

inductive *ctyping* :: *tyenv* ⇒ *com* ⇒ *bool* (**infix** ⊢*p* *50*) **where**
Γ ⊢*p* *SKIP* |
Γ ⊢*p* *a* : Γ(*x*) ⟹ Γ ⊢*p* *x* ::= *a* |
Γ ⊢*p* *c1* ⟹ Γ ⊢*p* *c2* ⟹ Γ ⊢*p* *c1*;;*c2* |
Γ ⊢*p* *b* ⟹ Γ ⊢*p* *c1* ⟹ Γ ⊢*p* *c2* ⟹ Γ ⊢*p* *IF b THEN c1 ELSE c2* |
Γ ⊢*p* *b* ⟹ Γ ⊢*p* *c* ⟹ Γ ⊢*p* *WHILE b DO c*

```
fun type :: val ⇒ ty where
type (Iv i) = Ity |
type (Rv r) = Rty

definition styping :: tyenv ⇒ state ⇒ bool (infix ⊢p 50)
where Γ ⊢p s ⟷ (∀x. type (s x) = Γ x)

fun tsubst :: (nat ⇒ ty) ⇒ ty ⇒ ty where
tsubst S (TV n) = S n |
tsubst S t = t
```

8.8 Typing is Preserved by Substitution

```
lemma subst_atyping: E ⊢p a : t ⟹ tsubst S ∘ E ⊢p a : tsubst S t
apply(induction rule: atyping.induct)
apply(auto intro: atyping.intros)
done

lemma subst_btyping: E ⊢p (b::bexp) ⟹ tsubst S ∘ E ⊢p b
apply(induction rule: btyping.induct)
apply(auto intro: btyping.intros)
apply(drule subst_atyping[where S=S])
apply(drule subst_atyping[where S=S])
apply(simp add: o_def btyping.intros)
done

lemma subst_ctyping: E ⊢p (c::com) ⟹ tsubst S ∘ E ⊢p c
apply(induction rule: ctyping.induct)
apply(auto intro: ctyping.intros)
apply(drule subst_atyping[where S=S])
apply(simp add: o_def ctyping.intros)
apply(drule subst_btyping[where S=S])
apply(simp add: o_def ctyping.intros)
apply(drule subst_btyping[where S=S])
apply(simp add: o_def ctyping.intros)
done

end
```

9 Security Type Systems

9.1 Security Levels and Expressions

theory *Sec_Type_Expr* **imports** *Big_Step*
begin

type_synonym *level* = *nat*

class *sec* =
fixes *sec* :: *'a* ⇒ *nat*

The security/confidentiality level of each variable is globally fixed for simplicity. For the sake of examples — the general theory does not rely on it! — a variable of length n has security level n:

instantiation *list* :: (*type*)*sec*
begin

definition *sec*(*x* :: *'a list*) = *length x*

instance ..

end

instantiation *aexp* :: *sec*
begin

fun *sec_aexp* :: *aexp* ⇒ *level* **where**
sec (*N n*) = *0* |
sec (*V x*) = *sec x* |
sec (*Plus* a_1 a_2) = *max* (*sec* a_1) (*sec* a_2)

instance ..

end

instantiation *bexp* :: *sec*
begin

fun *sec_bexp* :: *bexp* ⇒ *level* **where**
sec (*Bc v*) = *0* |
sec (*Not b*) = *sec b* |
sec (*And* b_1 b_2) = *max* (*sec* b_1) (*sec* b_2) |
sec (*Less* a_1 a_2) = *max* (*sec* a_1) (*sec* a_2)

instance ..

end

abbreviation *eq_le* :: *state* ⇒ *state* ⇒ *level* ⇒ *bool*
  ((_ = _ ′(≤ _′)) [*51*,*51*,*0*] *50*) **where**
$s = s' \ (\leq l) == (\forall \ x. \ sec \ x \leq l \longrightarrow s \ x = s' \ x)$

abbreviation *eq_less* :: *state* ⇒ *state* ⇒ *level* ⇒ *bool*
  ((_ = _ ′(< _′)) [*51*,*51*,*0*] *50*) **where**
$s = s' \ (< l) == (\forall \ x. \ sec \ x < l \longrightarrow s \ x = s' \ x)$

lemma *aval_eq_if_eq_le*:
  $[\![\ s_1 = s_2 \ (\leq l); \ \ sec \ a \leq l \ ]\!] \Longrightarrow aval \ a \ s_1 = aval \ a \ s_2$
by (*induct a*) *auto*

lemma *bval_eq_if_eq_le*:
  $[\![\ s_1 = s_2 \ (\leq l); \ \ sec \ b \leq l \ ]\!] \Longrightarrow bval \ b \ s_1 = bval \ b \ s_2$
by (*induct b*) (*auto simp add*: *aval_eq_if_eq_le*)

end

9.2 Security Typing of Commands

theory *Sec_Typing* **imports** *Sec_Type_Expr*
begin

9.2.1 Syntax Directed Typing

inductive *sec_type* :: *nat* ⇒ *com* ⇒ *bool* ((_/ ⊢ _) [*0*,*0*] *50*) **where**
Skip:
  $l \vdash SKIP$ |
Assign:
  $[\![\ sec \ x \geq sec \ a; \ \ sec \ x \geq l \ ]\!] \Longrightarrow l \vdash x ::= a$ |
Seq:
  $[\![\ l \vdash c_1; \ \ l \vdash c_2 \ ]\!] \Longrightarrow l \vdash c_1;;c_2$ |
If:
  $[\![\ max \ (sec \ b) \ l \vdash c_1; \ \ max \ (sec \ b) \ l \vdash c_2 \ ]\!] \Longrightarrow l \vdash IF \ b \ THEN \ c_1 \ ELSE \ c_2$ |
While:
  $max \ (sec \ b) \ l \vdash c \Longrightarrow l \vdash WHILE \ b \ DO \ c$

code_pred (*expected_modes*: *i* => *i* => *bool*) *sec_type* **.**

```
value 0 ⊢ IF Less (V ''x1'') (V ''x'') THEN ''x1'' ::= N 0 ELSE SKIP
value 1 ⊢ IF Less (V ''x1'') (V ''x'') THEN ''x'' ::= N 0 ELSE SKIP
value 2 ⊢ IF Less (V ''x1'') (V ''x'') THEN ''x1'' ::= N 0 ELSE SKIP

inductive_cases [elim!]:
  l ⊢ x ::= a  l ⊢ c1;;c2  l ⊢ IF b THEN c1 ELSE c2  l ⊢ WHILE b DO c
```

An important property: anti-monotonicity.

```
lemma anti_mono: ⟦ l ⊢ c;  l' ≤ l ⟧ ⟹ l' ⊢ c
apply(induction arbitrary: l' rule: sec_type.induct)
apply (metis sec_type.intros(1))
apply (metis le_trans sec_type.intros(2))
apply (metis sec_type.intros(3))
apply (metis If le_refl sup_mono sup_nat_def)
apply (metis While le_refl sup_mono sup_nat_def)
done

lemma confinement: ⟦ (c,s) ⇒ t;  l ⊢ c ⟧ ⟹ s = t (< l)
proof(induction rule: big_step_induct)
  case Skip thus ?case by simp
next
  case Assign thus ?case by auto
next
  case Seq thus ?case by auto
next
  case (IfTrue b s c1)
  hence max (sec b) l ⊢ c1 by auto
  hence l ⊢ c1 by (metis max.cobounded2 anti_mono)
  thus ?case using IfTrue.IH by metis
next
  case (IfFalse b s c2)
  hence max (sec b) l ⊢ c2 by auto
  hence l ⊢ c2 by (metis max.cobounded2 anti_mono)
  thus ?case using IfFalse.IH by metis
next
  case WhileFalse thus ?case by auto
next
  case (WhileTrue b s1 c)
  hence max (sec b) l ⊢ c by auto
  hence l ⊢ c by (metis max.cobounded2 anti_mono)
  thus ?case using WhileTrue by metis
qed
```

```
theorem noninterference:
  ⟦ (c,s) ⇒ s'; (c,t) ⇒ t';  0 ⊢ c;  s = t (≤ l) ⟧
   ⟹ s' = t' (≤ l)
proof(induction arbitrary: t t' rule: big_step_induct)
  case Skip thus ?case by auto
next
  case (Assign x a s)
  have [simp]: t' = t(x := aval a t) using Assign by auto
  have sec x >= sec a using ‹0 ⊢ x ::= a› by auto
  show ?case
  proof auto
    assume sec x ≤ l
    with ‹sec x >= sec a› have sec a ≤ l by arith
    thus aval a s = aval a t
      by (rule aval_eq_if_eq_le[OF ‹s = t (≤ l)›])
  next
    fix y assume y ≠ x sec y ≤ l
    thus s y = t y using ‹s = t (≤ l)› by simp
  qed
next
  case Seq thus ?case by blast
next
  case (IfTrue b s c1 s' c2)
  have sec b ⊢ c1 sec b ⊢ c2 using ‹0 ⊢ IF b THEN c1 ELSE c2› by auto
  show ?case
  proof cases
    assume sec b ≤ l
    hence s = t (≤ sec b) using ‹s = t (≤ l)› by auto
    hence bval b t using ‹bval b s› by(simp add: bval_eq_if_eq_le)
    with IfTrue.IH IfTrue.prems(1,3) ‹sec b ⊢ c1›  anti_mono
    show ?thesis by auto
  next
    assume ¬ sec b ≤ l
    have 1: sec b ⊢ IF b THEN c1 ELSE c2
      by(rule sec_type.intros)(simp_all add: ‹sec b ⊢ c1› ‹sec b ⊢ c2›)
    from confinement[OF ‹(c1, s) ⇒ s'› ‹sec b ⊢ c1›] ‹¬ sec b ≤ l›
    have s = s' (≤ l) by auto
    moreover
    from confinement[OF ‹(IF b THEN c1 ELSE c2, t) ⇒ t'› 1] ‹¬ sec b
≤ l›
    have t = t' (≤ l) by auto
    ultimately show s' = t' (≤ l) using ‹s = t (≤ l)› by auto
  qed
```

```
next
  case (IfFalse b s c2 s' c1)
  have sec b ⊢ c1 sec b ⊢ c2 using ‹0 ⊢ IF b THEN c1 ELSE c2› by auto
  show ?case
  proof cases
    assume sec b ≤ l
    hence s = t (≤ sec b) using ‹s = t (≤ l)› by auto
    hence ¬ bval b t using ‹¬ bval b s› by(simp add: bval_eq_if_eq_le)
    with IfFalse.IH IfFalse.prems(1,3) ‹sec b ⊢ c2› anti_mono
    show ?thesis by auto
  next
    assume ¬ sec b ≤ l
    have 1: sec b ⊢ IF b THEN c1 ELSE c2
      by(rule sec_type.intros)(simp_all add: ‹sec b ⊢ c1› ‹sec b ⊢ c2›)
     from confinement[OF big_step.IfFalse[OF IfFalse(1,2)] 1] ‹¬ sec b ≤
l›
    have s = s' (≤ l) by auto
    moreover
     from confinement[OF ‹(IF b THEN c1 ELSE c2, t) ⇒ t'› 1] ‹¬ sec b
≤ l›
    have t = t' (≤ l) by auto
    ultimately show s' = t' (≤ l) using ‹s = t (≤ l)› by auto
  qed
next
  case (WhileFalse b s c)
  have sec b ⊢ c using WhileFalse.prems(2) by auto
  show ?case
  proof cases
    assume sec b ≤ l
    hence s = t (≤ sec b) using ‹s = t (≤ l)› by auto
    hence ¬ bval b t using ‹¬ bval b s› by(simp add: bval_eq_if_eq_le)
    with WhileFalse.prems(1,3) show ?thesis by auto
  next
    assume ¬ sec b ≤ l
    have 1: sec b ⊢ WHILE b DO c
      by(rule sec_type.intros)(simp_all add: ‹sec b ⊢ c›)
    from confinement[OF ‹(WHILE b DO c, t) ⇒ t'› 1] ‹¬ sec b ≤ l›
    have t = t' (≤ l) by auto
    thus s = t' (≤ l) using ‹s = t (≤ l)› by auto
  qed
next
  case (WhileTrue b s1 c s2 s3 t1 t3)
  let ?w = WHILE b DO c
  have sec b ⊢ c using ‹0 ⊢ WHILE b DO c› by auto
```

```
  show ?case
  proof cases
    assume sec b ≤ l
    hence s1 = t1 (≤ sec b) using ‹s1 = t1 (≤ l)› by auto
    hence bval b t1
      using ‹bval b s1› by(simp add: bval_eq_if_eq_le)
    then obtain t2 where (c,t1) ⇒ t2 (?w,t2) ⇒ t3
      using ‹(?w,t1) ⇒ t3› by auto
    from WhileTrue.IH(2)[OF ‹(?w,t2) ⇒ t3› ‹0 ⊢ ?w›
      WhileTrue.IH(1)[OF ‹(c,t1) ⇒ t2› anti_mono[OF ‹sec b ⊢ c›]
        ‹s1 = t1 (≤ l)›]]
    show ?thesis by simp
  next
    assume ¬ sec b ≤ l
    have 1: sec b ⊢ ?w by(rule sec_type.intros)(simp_all add: ‹sec b ⊢ c›)
     from confinement[OF big_step.WhileTrue[OF WhileTrue.hyps] 1] ‹¬
sec b ≤ l›
    have s1 = s3 (≤ l) by auto
    moreover
    from confinement[OF ‹(WHILE b DO c, t1) ⇒ t3› 1] ‹¬ sec b ≤ l›
    have t1 = t3 (≤ l) by auto
    ultimately show s3 = t3 (≤ l) using ‹s1 = t1 (≤ l)› by auto
  qed
qed
```

9.2.2 The Standard Typing System

The predicate $l \vdash c$ is nicely intuitive and executable. The standard formulation, however, is slightly different, replacing the maximum computation by an antimonotonicity rule. We introduce the standard system now and show the equivalence with our formulation.

```
inductive sec_type' :: nat ⇒ com ⇒ bool ((_/ ⊢'' _) [0,0] 50) where
Skip':
  l ⊢' SKIP |
Assign':
  ⟦ sec x ≥ sec a; sec x ≥ l ⟧ ⟹ l ⊢' x ::= a |
Seq':
  ⟦ l ⊢' c₁;  l ⊢' c₂ ⟧ ⟹ l ⊢' c₁;;c₂ |
If':
  ⟦ sec b ≤ l;  l ⊢' c₁;  l ⊢' c₂ ⟧ ⟹ l ⊢' IF b THEN c₁ ELSE c₂ |
While':
  ⟦ sec b ≤ l;  l ⊢' c ⟧ ⟹ l ⊢' WHILE b DO c |
anti_mono':
  ⟦ l ⊢' c;  l' ≤ l ⟧ ⟹ l' ⊢' c
```

lemma *sec_type_sec_type'*: $l \vdash c \Longrightarrow l \vdash' c$
apply(*induction rule*: *sec_type.induct*)
apply (*metis Skip'*)
apply (*metis Assign'*)
apply (*metis Seq'*)
apply (*metis max.commute max.absorb_iff2 nat_le_linear If' anti_mono'*)
by (*metis less_or_eq_imp_le max.absorb1 max.absorb2 nat_le_linear While' anti_mono'*)

lemma *sec_type'_sec_type*: $l \vdash' c \Longrightarrow l \vdash c$
apply(*induction rule*: *sec_type'.induct*)
apply (*metis Skip*)
apply (*metis Assign*)
apply (*metis Seq*)
apply (*metis max.absorb2 If*)
apply (*metis max.absorb2 While*)
by (*metis anti_mono*)

9.2.3 A Bottom-Up Typing System

inductive *sec_type2* :: *com* ⇒ *level* ⇒ *bool* ((⊢ _ : _) [*0*,*0*] *50*) **where**
Skip2:
  $\vdash$ *SKIP* : l |
Assign2:
  $sec\ x \geq sec\ a \Longrightarrow \vdash x ::= a : sec\ x$ |
Seq2:
  $[\![\vdash c_1 : l_1;\ \vdash c_2 : l_2]\!] \Longrightarrow \vdash c_1;;c_2 : min\ l_1\ l_2$ |
If2:
  $[\![sec\ b \leq min\ l_1\ l_2;\ \vdash c_1 : l_1;\ \vdash c_2 : l_2]\!]$
  $\Longrightarrow \vdash$ *IF* b *THEN* c_1 *ELSE* c_2 : $min\ l_1\ l_2$ |
While2:
  $[\![sec\ b \leq l;\ \vdash c : l]\!] \Longrightarrow \vdash$ *WHILE* b *DO* c : l

lemma *sec_type2_sec_type'*: $\vdash c : l \Longrightarrow l \vdash' c$
apply(*induction rule*: *sec_type2.induct*)
apply (*metis Skip'*)
apply (*metis Assign' eq_imp_le*)
apply (*metis Seq' anti_mono' min.cobounded1 min.cobounded2*)
apply (*metis If' anti_mono' min.absorb2 min.absorb_iff1 nat_le_linear*)
by (*metis While'*)

lemma *sec_type'_sec_type2*: $l \vdash' c \Longrightarrow \exists\ l' \geq l.\ \vdash c : l'$
apply(*induction rule*: *sec_type'.induct*)
apply (*metis Skip2 le_refl*)
apply (*metis Assign2*)
apply (*metis Seq2 min.boundedI*)
apply (*metis If2 inf_greatest inf_nat_def le_trans*)
apply (*metis While2 le_trans*)
by (*metis le_trans*)

end

9.3 Termination-Sensitive Systems

theory *Sec_TypingT* **imports** *Sec_Type_Expr*
begin

9.3.1 A Syntax Directed System

inductive *sec_type* :: *nat* ⇒ *com* ⇒ *bool* ((_/ ⊢ _) [*0*,*0*] *50*) **where**
Skip:
$l \vdash SKIP$ |
Assign:
$[\![\ sec\ x \geq sec\ a;\ \ sec\ x \geq l\ ]\!] \Longrightarrow l \vdash x ::= a$ |
Seq:
$l \vdash c_1 \Longrightarrow l \vdash c_2 \Longrightarrow l \vdash c_1;;c_2$ |
If:
$[\![\ max\ (sec\ b)\ l \vdash c_1;\ \ max\ (sec\ b)\ l \vdash c_2\ ]\!]$
$\Longrightarrow l \vdash IF\ b\ THEN\ c_1\ ELSE\ c_2$ |
While:
$sec\ b = 0 \Longrightarrow 0 \vdash c \Longrightarrow 0 \vdash WHILE\ b\ DO\ c$

code_pred (*expected_modes*: *i* => *i* => *bool*) *sec_type* **.**

inductive_cases [*elim*!]:
$l \vdash x ::= a$ $l \vdash c_1;;c_2$ $l \vdash IF\ b\ THEN\ c_1\ ELSE\ c_2$ $l \vdash WHILE\ b\ DO\ c$

lemma *anti_mono*: $l \vdash c \Longrightarrow l' \leq l \Longrightarrow l' \vdash c$
apply(*induction arbitrary*: *l' rule*: *sec_type.induct*)
apply (*metis sec_type.intros*(*1*))
apply (*metis le_trans sec_type.intros*(*2*))
apply (*metis sec_type.intros*(*3*))
apply (*metis If le_refl sup_mono sup_nat_def*)
by (*metis While le_0_eq*)

```
lemma confinement: (c,s) ⇒ t ⟹ l ⊢ c ⟹ s = t (< l)
proof(induction rule: big_step_induct)
  case Skip thus ?case by simp
next
  case Assign thus ?case by auto
next
  case Seq thus ?case by auto
next
  case (IfTrue b s c1)
  hence max (sec b) l ⊢ c1 by auto
  hence l ⊢ c1 by (metis max.cobounded2 anti_mono)
  thus ?case using IfTrue.IH by metis
next
  case (IfFalse b s c2)
  hence max (sec b) l ⊢ c2 by auto
  hence l ⊢ c2 by (metis max.cobounded2 anti_mono)
  thus ?case using IfFalse.IH by metis
next
  case WhileFalse thus ?case by auto
next
  case (WhileTrue b s1 c)
  hence l ⊢ c by auto
  thus ?case using WhileTrue by metis
qed

lemma termi_if_non0: l ⊢ c ⟹ l ≠ 0 ⟹ ∃ t. (c,s) ⇒ t
apply(induction arbitrary: s rule: sec_type.induct)
apply (metis big_step.Skip)
apply (metis big_step.Assign)
apply (metis big_step.Seq)
apply (metis IfFalse IfTrue le0 le_antisym max.cobounded2)
apply simp
done

theorem noninterference: (c,s) ⇒ s' ⟹ 0 ⊢ c ⟹  s = t (≤ l)
  ⟹ ∃ t'. (c,t) ⇒ t' ∧ s' = t' (≤ l)
proof(induction arbitrary: t rule: big_step_induct)
  case Skip thus ?case by auto
next
  case (Assign x a s)
  have sec x >= sec a using ‹0 ⊢ x ::= a› by auto
  have (x ::= a,t) ⇒ t(x := aval a t) by auto
```

```
  moreover
  have s(x := aval a s) = t(x := aval a t) (≤ l)
  proof auto
    assume sec x ≤ l
    with ‹sec x ≥ sec a› have sec a ≤ l by arith
    thus aval a s = aval a t
      by (rule aval_eq_if_eq_le[OF ‹s = t (≤ l)›])
  next
    fix y assume y ≠ x sec y ≤ l
    thus s y = t y using ‹s = t (≤ l)› by simp
  qed
  ultimately show ?case by blast
next
  case Seq thus ?case by blast
next
  case (IfTrue b s c1 s' c2)
  have sec b ⊢ c1 sec b ⊢ c2 using ‹0 ⊢ IF b THEN c1 ELSE c2› by auto
  obtain t' where t': (c1, t) ⇒ t' s' = t' (≤ l)
   using IfTrue.IH[OF anti_mono[OF ‹sec b ⊢ c1›] ‹s = t (≤ l)›] by blast
  show ?case
  proof cases
    assume sec b ≤ l
    hence s = t (≤ sec b) using ‹s = t (≤ l)› by auto
    hence bval b t using ‹bval b s› by(simp add: bval_eq_if_eq_le)
    thus ?thesis by (metis t' big_step.IfTrue)
  next
    assume ¬ sec b ≤ l
    hence 0: sec b ≠ 0 by arith
    have 1: sec b ⊢ IF b THEN c1 ELSE c2
      by(rule sec_type.intros)(simp_all add: ‹sec b ⊢ c1› ‹sec b ⊢ c2›)
    from confinement[OF big_step.IfTrue[OF IfTrue(1,2)] 1] ‹¬ sec b ≤ l›
    have s = s' (≤ l) by auto
    moreover
    from termi_if_non0[OF 1 0, of t] obtain t' where
      t': (IF b THEN c1 ELSE c2,t) ⇒ t' ..
    moreover
    from confinement[OF t' 1] ‹¬ sec b ≤ l›
    have t = t' (≤ l) by auto
    ultimately
    show ?case using ‹s = t (≤ l)› by auto
  qed
next
  case (IfFalse b s c2 s' c1)
  have sec b ⊢ c1 sec b ⊢ c2 using ‹0 ⊢ IF b THEN c1 ELSE c2› by auto
```

```
  obtain t' where t': (c2, t) ⇒ t' s' = t' (≤ l)
    using IfFalse.IH[OF anti_mono[OF ‹sec b ⊢ c2›] ‹s = t (≤ l)›] by blast
  show ?case
  proof cases
    assume sec b ≤ l
    hence s = t (≤ sec b) using ‹s = t (≤ l)› by auto
    hence ¬ bval b t using ‹¬ bval b s› by(simp add: bval_eq_if_eq_le)
    thus ?thesis by (metis t' big_step.IfFalse)
  next
    assume ¬ sec b ≤ l
    hence 0: sec b ≠ 0 by arith
    have 1: sec b ⊢ IF b THEN c1 ELSE c2
      by(rule sec_type.intros)(simp_all add: ‹sec b ⊢ c1› ‹sec b ⊢ c2›)
    from confinement[OF big_step.IfFalse[OF IfFalse(1,2)] 1] ‹¬ sec b ≤ l›
    have s = s' (≤ l) by auto
    moreover
    from termi_if_non0[OF 1 0, of t] obtain t' where
      t': (IF b THEN c1 ELSE c2,t) ⇒ t' ..
    moreover
    from confinement[OF t' 1] ‹¬ sec b ≤ l›
    have t = t' (≤ l) by auto
    ultimately
    show ?case using ‹s = t (≤ l)› by auto
  qed
next
  case (WhileFalse b s c)
  hence [simp]: sec b = 0 by auto
  have s = t (≤ sec b) using ‹s = t (≤ l)› by auto
  hence ¬ bval b t using ‹¬ bval b s› by (metis bval_eq_if_eq_le le_refl)
  with WhileFalse.prems(2) show ?case by auto
next
  case (WhileTrue b s c s'' s')
  let ?w = WHILE b DO c
  from ‹0 ⊢ ?w› have [simp]: sec b = 0 by auto
  have 0 ⊢ c using ‹0 ⊢ WHILE b DO c› by auto
  from WhileTrue.IH(1)[OF this ‹s = t (≤ l)›]
  obtain t'' where (c,t) ⇒ t'' and s'' = t'' (≤l) by blast
  from WhileTrue.IH(2)[OF ‹0 ⊢ ?w› this(2)]
  obtain t' where (?w,t'') ⇒ t' and s' = t' (≤l) by blast
  from ‹bval b s› have bval b t
    using bval_eq_if_eq_le[OF ‹s = t (≤l)›] by auto
  show ?case
```

using *big_step.WhileTrue*[*OF* ‹*bval b t*› ‹(*c*,*t*) ⇒ *t''*› ‹(*?w*,*t''*) ⇒ *t'*›]
by (*metis* ‹*s'* = *t'* (≤ *l*)›)
qed

9.3.2 The Standard System

The predicate $l \vdash c$ is nicely intuitive and executable. The standard formulation, however, is slightly different, replacing the maximum computation by an antimonotonicity rule. We introduce the standard system now and show the equivalence with our formulation.

inductive *sec_type'* :: *nat* ⇒ *com* ⇒ *bool* ((_/ ⊢'' _) [*0*,*0*] *50*) **where**
Skip':
$l \vdash' SKIP$ |
Assign':
$[\![sec\ x \geq sec\ a;\ \ sec\ x \geq l]\!] \Longrightarrow l \vdash' x ::= a$ |
Seq':
$l \vdash' c_1 \Longrightarrow l \vdash' c_2 \Longrightarrow l \vdash' c_1;;c_2$ |
If':
$[\![sec\ b \leq l;\ \ l \vdash' c_1;\ \ l \vdash' c_2]\!] \Longrightarrow l \vdash' IF\ b\ THEN\ c_1\ ELSE\ c_2$ |
While':
$[\![sec\ b = 0;\ \ 0 \vdash' c]\!] \Longrightarrow 0 \vdash' WHILE\ b\ DO\ c$ |
anti_mono':
$[\![l \vdash' c;\ \ l' \leq l]\!] \Longrightarrow l' \vdash' c$

lemma *sec_type_sec_type'*:
$l \vdash c \Longrightarrow l \vdash' c$
apply(*induction rule*: *sec_type.induct*)
apply (*metis Skip'*)
apply (*metis Assign'*)
apply (*metis Seq'*)
apply (*metis max.commute max.absorb_iff2 nat_le_linear If' anti_mono'*)
by (*metis While'*)

lemma *sec_type'_sec_type*:
$l \vdash' c \Longrightarrow l \vdash c$
apply(*induction rule*: *sec_type'.induct*)
apply (*metis Skip*)
apply (*metis Assign*)
apply (*metis Seq*)
apply (*metis max.absorb2 If*)
apply (*metis While*)
by (*metis anti_mono*)

corollary *sec_type_eq*: $l \vdash c \longleftrightarrow l \vdash' c$
by (*metis sec_type'_sec_type sec_type_sec_type'*)

end

10 Definite Initialization Analysis

theory *Vars* **imports** *Com*
begin

10.1 The Variables in an Expression

We need to collect the variables in both arithmetic and boolean expressions. For a change we do not introduce two functions, e.g. *avars* and *bvars*, but we overload the name *vars* via a *type class*, a device that originated with Haskell:

class *vars* =
fixes *vars* :: *'a* ⇒ *vname set*

This defines a type class "vars" with a single function of (coincidentally) the same name. Then we define two separated instances of the class, one for *aexp* and one for *bexp*:

instantiation *aexp* :: *vars*
begin

fun *vars_aexp* :: *aexp* ⇒ *vname set* **where**
vars (*N n*) = {} |
vars (*V x*) = {*x*} |
vars (*Plus* a_1 a_2) = *vars* a_1 ∪ *vars* a_2

instance ..

end

value *vars* (*Plus* (*V ''x''*) (*V ''y''*))

instantiation *bexp* :: *vars*
begin

fun *vars_bexp* :: *bexp* ⇒ *vname set* **where**
vars (*Bc v*) = {} |
vars (*Not b*) = *vars b* |
vars (*And* b_1 b_2) = *vars* b_1 ∪ *vars* b_2 |

$vars\ (Less\ a_1\ \ a_2) = vars\ a_1 \cup vars\ a_2$

instance ..

end

value *vars* (*Less* (*Plus* (*V* ''*z*'') (*V* ''*y*'')) (*V* ''*x*''))

abbreviation
eq_on :: ('*a* ⇒ '*b*) ⇒ ('*a* ⇒ '*b*) ⇒ '*a set* ⇒ *bool*
((_ =/ _/ *on* _) [*50*,*0*,*50*] *50*) **where**
$f = g\ on\ X == \forall\ x \in X.\ f\ x = g\ x$

lemma *aval_eq_if_eq_on_vars*[*simp*]:
$s_1 = s_2\ on\ vars\ a \Longrightarrow aval\ a\ s_1 = aval\ a\ s_2$
apply(*induction a*)
apply *simp_all*
done

lemma *bval_eq_if_eq_on_vars*:
$s_1 = s_2\ on\ vars\ b \Longrightarrow bval\ b\ s_1 = bval\ b\ s_2$
proof(*induction b*)
case (*Less a1 a2*)
hence $aval\ a1\ s_1 = aval\ a1\ s_2$ **and** $aval\ a2\ s_1 = aval\ a2\ s_2$ **by** *simp_all*
thus *?case* **by** *simp*
qed *simp_all*

fun *lvars* :: *com* ⇒ *vname set* **where**
lvars SKIP = {} |
lvars (*x*::=*e*) = {*x*} |
lvars (*c1*;;*c2*) = *lvars c1* ∪ *lvars c2* |
lvars (*IF b THEN c1 ELSE c2*) = *lvars c1* ∪ *lvars c2* |
lvars (*WHILE b DO c*) = *lvars c*

fun *rvars* :: *com* ⇒ *vname set* **where**
rvars SKIP = {} |
rvars (*x*::=*e*) = *vars e* |
rvars (*c1*;;*c2*) = *rvars c1* ∪ *rvars c2* |
rvars (*IF b THEN c1 ELSE c2*) = *vars b* ∪ *rvars c1* ∪ *rvars c2* |
rvars (*WHILE b DO c*) = *vars b* ∪ *rvars c*

instantiation *com* :: *vars*
begin

definition *vars_com c* = *lvars c* ∪ *rvars c*

instance ..

end

lemma *vars_com_simps*[*simp*]:
  vars SKIP = {}
  vars (*x*::=*e*) = {*x*} ∪ *vars e*
  vars (*c1*;;*c2*) = *vars c1* ∪ *vars c2*
  vars (*IF b THEN c1 ELSE c2*) = *vars b* ∪ *vars c1* ∪ *vars c2*
  vars (*WHILE b DO c*) = *vars b* ∪ *vars c*
by(*auto simp*: *vars_com_def*)

lemma *finite_avars*[*simp*]: *finite*(*vars*(*a*::*aexp*))
by(*induction a*) *simp_all*

lemma *finite_bvars*[*simp*]: *finite*(*vars*(*b*::*bexp*))
by(*induction b*) *simp_all*

lemma *finite_lvars*[*simp*]: *finite*(*lvars*(*c*))
by(*induction c*) *simp_all*

lemma *finite_rvars*[*simp*]: *finite*(*rvars*(*c*))
by(*induction c*) *simp_all*

lemma *finite_cvars*[*simp*]: *finite*(*vars*(*c*::*com*))
by(*simp add*: *vars_com_def*)

end

theory *Def_Init_Exp*
imports *Vars*
begin

10.2 Initialization-Sensitive Expressions Evaluation

type_synonym *state* = *vname* ⇒ *val option*

fun *aval* :: *aexp* ⇒ *state* ⇒ *val option* **where**
aval (*N i*) *s* = *Some i* |
aval (*V x*) *s* = *s x* |

$aval\ (Plus\ a_1\ a_2)\ s =$
  $(case\ (aval\ a_1\ s,\ aval\ a_2\ s)\ of$
    $(Some\ i_1, Some\ i_2) \Rightarrow Some(i_1+i_2) \mid \_ \Rightarrow None)$

fun *bval* :: *bexp* ⇒ *state* ⇒ *bool option* **where**
$bval\ (Bc\ v)\ s = Some\ v \mid$
$bval\ (Not\ b)\ s = (case\ bval\ b\ s\ of\ None \Rightarrow None \mid Some\ bv \Rightarrow Some(\neg\ bv))$
|
$bval\ (And\ b_1\ b_2)\ s = (case\ (bval\ b_1\ s,\ bval\ b_2\ s)\ of$
  $(Some\ bv_1,\ Some\ bv_2) \Rightarrow Some(bv_1\ \&\ bv_2) \mid \_ \Rightarrow None) \mid$
$bval\ (Less\ a_1\ a_2)\ s = (case\ (aval\ a_1\ s,\ aval\ a_2\ s)\ of$
 $(Some\ i_1,\ Some\ i_2) \Rightarrow Some(i_1 < i_2) \mid \_ \Rightarrow None)$

lemma *aval_Some*: *vars a* ⊆ *dom s* ⟹ ∃ *i. aval a s* = *Some i*
by (*induct a*) *auto*

lemma *bval_Some*: *vars b* ⊆ *dom s* ⟹ ∃ *bv. bval b s* = *Some bv*
by (*induct b*) (*auto dest*!: *aval_Some*)

end
theory *Def_Init*
imports *Vars Com*
begin

10.3 Definite Initialization Analysis

inductive *D* :: *vname set* ⇒ *com* ⇒ *vname set* ⇒ *bool* **where**
Skip: *D A SKIP A* |
Assign: *vars a* ⊆ *A* ⟹ *D A* (*x* ::= *a*) (*insert x A*) |
Seq: $[\![\ D\ A_1\ c_1\ A_2;\ \ D\ A_2\ c_2\ A_3\ ]\!] \Longrightarrow D\ A_1\ (c_1;;\ c_2)\ A_3 \mid$
If: $[\![\ vars\ b \subseteq A;\ \ D\ A\ c_1\ A_1;\ \ D\ A\ c_2\ A_2\ ]\!] \Longrightarrow$
  $D\ A\ (IF\ b\ THEN\ c_1\ ELSE\ c_2)\ (A_1\ Int\ A_2) \mid$
While: $[\![\ vars\ b \subseteq A;\ \ D\ A\ c\ A'\ ]\!] \Longrightarrow D\ A\ (WHILE\ b\ DO\ c)\ A$

inductive_cases [*elim*!]:
$D\ A\ SKIP\ A'$
$D\ A\ (x ::= a)\ A'$
$D\ A\ (c1;;c2)\ A'$
$D\ A\ (IF\ b\ THEN\ c1\ ELSE\ c2)\ A'$
$D\ A\ (WHILE\ b\ DO\ c)\ A'$

lemma *D_incr*:

$D\ A\ c\ A' \Longrightarrow A \subseteq A'$
by (*induct rule*: *D.induct*) *auto*

end

theory *Def_Init_Big*
imports *Def_Init_Exp Def_Init*
begin

10.4 Initialization-Sensitive Big Step Semantics

inductive
big_step :: (*com* × *state option*) ⇒ *state option* ⇒ *bool* (**infix** ⇒ *55*)
where
None: (*c*,*None*) ⇒ *None* |
Skip: (*SKIP*,*s*) ⇒ *s* |
AssignNone: *aval a s* = *None* ⟹ (*x* ::= *a*, *Some s*) ⇒ *None* |
Assign: *aval a s* = *Some i* ⟹ (*x* ::= *a*, *Some s*) ⇒ *Some*(*s*(*x* := *Some i*))
|
Seq: $(c_1,s_1) \Rightarrow s_2 \Longrightarrow (c_2,s_2) \Rightarrow s_3 \Longrightarrow (c_1;;c_2,s_1) \Rightarrow s_3$ |

IfNone: *bval b s* = *None* ⟹ (*IF b THEN* c_1 *ELSE* c_2,*Some s*) ⇒ *None* |
IfTrue: ⟦ *bval b s* = *Some True*; (c_1,*Some s*) ⇒ *s'* ⟧ ⟹
(*IF b THEN* c_1 *ELSE* c_2,*Some s*) ⇒ *s'* |
IfFalse: ⟦ *bval b s* = *Some False*; (c_2,*Some s*) ⇒ *s'* ⟧ ⟹
(*IF b THEN* c_1 *ELSE* c_2,*Some s*) ⇒ *s'* |

WhileNone: *bval b s* = *None* ⟹ (*WHILE b DO c*,*Some s*) ⇒ *None* |
WhileFalse: *bval b s* = *Some False* ⟹ (*WHILE b DO c*,*Some s*) ⇒ *Some s* |
WhileTrue:
⟦ *bval b s* = *Some True*; (*c*,*Some s*) ⇒ *s'*; (*WHILE b DO c*,*s'*) ⇒ *s''* ⟧
⟹
(*WHILE b DO c*,*Some s*) ⇒ *s''*

lemmas *big_step_induct* = *big_step.induct*[*split_format*(*complete*)]

10.5 Soundness wrt Big Steps

Note the special form of the induction because one of the arguments of the inductive predicate is not a variable but the term *Some s*:

theorem *Sound*:
⟦ (*c*,*Some s*) ⇒ *s'*; *D A c A'*; *A* ⊆ *dom s* ⟧

$\Longrightarrow \exists\ t.\ s' = Some\ t \wedge A' \subseteq dom\ t$
proof (*induction c Some s s' arbitrary*: *s A A' rule*:*big_step_induct*)
  case *AssignNone* **thus** *?case*
    by *auto* (*metis aval_Some option.simps*(*3*) *subset_trans*)
next
  case *Seq* **thus** *?case* **by** *auto metis*
next
  case *IfTrue* **thus** *?case* **by** *auto blast*
next
  case *IfFalse* **thus** *?case* **by** *auto blast*
next
  case *IfNone* **thus** *?case*
    by *auto* (*metis bval_Some option.simps*(*3*) *order_trans*)
next
  case *WhileNone* **thus** *?case*
    by *auto* (*metis bval_Some option.simps*(*3*) *order_trans*)
next
  case (*WhileTrue b s c s' s''*)
  from ‹*D A* (*WHILE b DO c*) *A'*› **obtain** *A'* **where** *D A c A'* **by** *blast*
  then obtain *t'* **where** $s' = Some\ t'\ A \subseteq dom\ t'$
    by (*metis D_incr WhileTrue*(*3*,*7*) *subset_trans*)
  from *WhileTrue*(*5*)[*OF this*(*1*) *WhileTrue*(*6*) *this*(*2*)] **show** *?case* **.**
qed *auto*

corollary *sound*: ⟦ *D* (*dom s*) *c A'*; (*c*,*Some s*) ⇒ *s'* ⟧ $\Longrightarrow s' \neq None$
by (*metis Sound not_Some_eq subset_refl*)

end

theory *Def_Init_Small*
imports *Star Def_Init_Exp Def_Init*
begin

10.6 Initialization-Sensitive Small Step Semantics

inductive
  small_step :: (*com* × *state*) ⇒ (*com* × *state*) ⇒ *bool* (**infix** → *55*)
where
Assign: *aval a s* = *Some i* ⟹ (*x* ::= *a*, *s*) → (*SKIP*, *s*(*x* := *Some i*)) |

Seq1: (*SKIP*;;*c*,*s*) → (*c*,*s*) |
Seq2: $(c_1,s) \to (c_1',s') \Longrightarrow (c_1;;c_2,s) \to (c_1';;c_2,s')$ |

IfTrue: *bval b s = Some True* $\Longrightarrow$ (*IF b THEN* c_1 *ELSE* c_2,*s*) $\to$ (c_1,*s*) |
IfFalse: *bval b s = Some False* $\Longrightarrow$ (*IF b THEN* c_1 *ELSE* c_2,*s*) $\to$ (c_2,*s*) |

While: (*WHILE b DO c*,*s*) $\to$ (*IF b THEN c*;; *WHILE b DO c ELSE SKIP*,*s*)

lemmas *small_step_induct = small_step.induct*[*split_format*(*complete*)]

abbreviation *small_steps* :: *com* ∗ *state* ⇒ *com* ∗ *state* ⇒ *bool* (**infix** →∗ *55*)
where *x* →∗ *y* == *star small_step x y*

10.7 Soundness wrt Small Steps

theorem *progress*:
D (*dom s*) *c A'* ⟹ *c* ≠ *SKIP* ⟹ ∃ *cs'*. (*c*,*s*) → *cs'*
proof (*induction c arbitrary*: *s A'*)
case *Assign* **thus** *?case* **by** *auto* (*metis aval_Some small_step.Assign*)
next
case (*If b c1 c2*)
then obtain *bv* **where** *bval b s = Some bv* **by** (*auto dest*!:*bval_Some*)
then show *?case*
by(*cases bv*)(*auto intro*: *small_step.IfTrue small_step.IfFalse*)
qed (*fastforce intro*: *small_step.intros*)+

lemma *D_mono*: *D A c M* ⟹ *A* ⊆ *A'* ⟹ ∃ *M'*. *D A' c M'* & *M* <= *M'*
proof (*induction c arbitrary*: *A A' M*)
case *Seq* **thus** *?case* **by** *auto* (*metis D.intros*(*3*))
next
case (*If b c1 c2*)
then obtain *M1 M2* **where** *vars b* ⊆ *A D A c1 M1 D A c2 M2 M* = *M1* ∩ *M2*
by *auto*
with *If.IH* ‹*A* ⊆ *A'*› **obtain** *M1' M2'*
where *D A' c1 M1' D A' c2 M2'* **and** *M1* ⊆ *M1' M2* ⊆ *M2'* **by** *metis*
hence *D A'* (*IF b THEN c1 ELSE c2*) (*M1'* ∩ *M2'*) **and** *M* ⊆ *M1'* ∩ *M2'*
using ‹*vars b* ⊆ *A*› ‹*A* ⊆ *A'*› ‹*M* = *M1* ∩ *M2*› **by**(*fastforce intro*: *D.intros*)+
thus *?case* **by** *metis*
next
case *While* **thus** *?case* **by** *auto* (*metis D.intros*(*5*) *subset_trans*)
qed (*auto intro*: *D.intros*)

theorem *D_preservation*:
(*c*,*s*) → (*c*′,*s*′) ⟹ *D* (*dom s*) *c A* ⟹ ∃ *A*′. *D* (*dom s*′) *c*′ *A*′ & *A* <= *A*′
proof (*induction arbitrary*: *A rule*: *small_step_induct*)
case (*While b c s*)
then obtain *A*′ **where** *A*′: *vars b* ⊆ *dom s A* = *dom s D* (*dom s*) *c A*′ **by** *blast*
then obtain *A*″ **where** *D A*′ *c A*″ **by** (*metis D_incr D_mono*)
with *A*′ **have** *D* (*dom s*) (*IF b THEN c*;; *WHILE b DO c ELSE SKIP*) (*dom s*)
by (*metis D.If*[*OF* ‹*vars b* ⊆ *dom s*› *D.Seq*[*OF* ‹*D* (*dom s*) *c A*′› *D.While*[*OF* _ ‹*D A*′ *c A*″›]] *D.Skip*] *D_incr Int_absorb1 subset_trans*)
thus *?case* **by** (*metis D_incr* ‹*A* = *dom s*›)
next
case *Seq2* **thus** *?case* **by** *auto* (*metis D_mono D.intros*(*3*))
qed (*auto intro*: *D.intros*)

theorem *D_sound*:
(*c*,*s*) →∗ (*c*′,*s*′) ⟹ *D* (*dom s*) *c A*′
⟹ (∃ *cs*″. (*c*′,*s*′) → *cs*″) ∨ *c*′ = *SKIP*
apply(*induction arbitrary*: *A*′ *rule*:*star_induct*)
apply (*metis progress*)
by (*metis D_preservation*)

end

11 Constant Folding

theory *Sem_Equiv*
imports *Big_Step*
begin

11.1 Semantic Equivalence up to a Condition

type_synonym *assn* = *state* ⇒ *bool*

definition
equiv_up_to :: *assn* ⇒ *com* ⇒ *com* ⇒ *bool* (_ ⊨ _ ∼ _ [*50*,*0*,*10*] *50*)
where
(*P* ⊨ *c* ∼ *c*′) = (∀ *s s*′. *P s* ⟶ (*c*,*s*) ⇒ *s*′ ⟷ (*c*′,*s*) ⇒ *s*′)

definition
bequiv_up_to :: *assn* ⇒ *bexp* ⇒ *bexp* ⇒ *bool* (_ ⊨ _ <∼> _ [*50*,*0*,*10*] *50*)

where
  $(P \models b <\sim> b') = (\forall s.\ P\ s \longrightarrow bval\ b\ s = bval\ b'\ s)$

lemma *equiv_up_to_True*:
  $((\lambda\_.\ True) \models c \sim c') = (c \sim c')$
  by (*simp add*: *equiv_def equiv_up_to_def*)

lemma *equiv_up_to_weaken*:
  $P \models c \sim c' \Longrightarrow (\bigwedge s.\ P'\ s \Longrightarrow P\ s) \Longrightarrow P' \models c \sim c'$
  by (*simp add*: *equiv_up_to_def*)

lemma *equiv_up_toI*:
  $(\bigwedge s\ s'.\ P\ s \Longrightarrow (c,\ s) \Rightarrow s' = (c',\ s) \Rightarrow s') \Longrightarrow P \models c \sim c'$
  by (*unfold equiv_up_to_def*) *blast*

lemma *equiv_up_toD1*:
  $P \models c \sim c' \Longrightarrow (c,\ s) \Rightarrow s' \Longrightarrow P\ s \Longrightarrow (c',\ s) \Rightarrow s'$
  by (*unfold equiv_up_to_def*) *blast*

lemma *equiv_up_toD2*:
  $P \models c \sim c' \Longrightarrow (c',\ s) \Rightarrow s' \Longrightarrow P\ s \Longrightarrow (c,\ s) \Rightarrow s'$
  by (*unfold equiv_up_to_def*) *blast*

lemma *equiv_up_to_refl* [*simp*, *intro*!]:
  $P \models c \sim c$
  by (*auto simp*: *equiv_up_to_def*)

lemma *equiv_up_to_sym*:
  $(P \models c \sim c') = (P \models c' \sim c)$
  by (*auto simp*: *equiv_up_to_def*)

lemma *equiv_up_to_trans*:
  $P \models c \sim c' \Longrightarrow P \models c' \sim c'' \Longrightarrow P \models c \sim c''$
  by (*auto simp*: *equiv_up_to_def*)

lemma *bequiv_up_to_refl* [*simp*, *intro*!]:
  $P \models b <\sim> b$
  by (*auto simp*: *bequiv_up_to_def*)

lemma *bequiv_up_to_sym*:
  $(P \models b <\sim> b') = (P \models b' <\sim> b)$
  by (*auto simp*: *bequiv_up_to_def*)

```
lemma bequiv_up_to_trans:
  P ⊨ b <∼> b′ ⟹ P ⊨ b′ <∼> b″ ⟹ P ⊨ b <∼> b″
  by (auto simp: bequiv_up_to_def)

lemma bequiv_up_to_subst:
  P ⊨ b <∼> b′ ⟹ P s ⟹ bval b s = bval b′ s
  by (simp add: bequiv_up_to_def)


lemma equiv_up_to_seq:
  P ⊨ c ∼ c′ ⟹ Q ⊨ d ∼ d′ ⟹
  (⋀s s′. (c,s) ⇒ s′ ⟹ P s ⟹ Q s′) ⟹
  P ⊨ (c;; d) ∼ (c′;; d′)
  by (clarsimp simp: equiv_up_to_def) blast

lemma equiv_up_to_while_lemma_weak:
  shows (d,s) ⇒ s′ ⟹
        P ⊨ b <∼> b′ ⟹
        P ⊨ c ∼ c′ ⟹
        (⋀s s′. (c, s) ⇒ s′ ⟹ P s ⟹ bval b s ⟹ P s′) ⟹
        P s ⟹
        d = WHILE b DO c ⟹
        (WHILE b′ DO c′, s) ⇒ s′
proof (induction rule: big_step_induct)
  case (WhileTrue b s1 c s2 s3)
  hence IH: P s2 ⟹ (WHILE b′ DO c′, s2) ⇒ s3 by auto
  from WhileTrue.prems
  have P ⊨ b <∼> b′ by simp
  with ‹bval b s1› ‹P s1›
  have bval b′ s1 by (simp add: bequiv_up_to_def)
  moreover
  from WhileTrue.prems
  have P ⊨ c ∼ c′ by simp
  with ‹bval b s1› ‹P s1› ‹(c, s1) ⇒ s2›
  have (c′, s1) ⇒ s2 by (simp add: equiv_up_to_def)
  moreover
  from WhileTrue.prems
  have ⋀s s′. (c,s) ⇒ s′ ⟹ P s ⟹ bval b s ⟹ P s′ by simp
  with ‹P s1› ‹bval b s1› ‹(c, s1) ⇒ s2›
  have P s2 by simp
  hence (WHILE b′ DO c′, s2) ⇒ s3 by (rule IH)
  ultimately
  show ?case by blast
```

next
  case *WhileFalse*
  thus *?case* **by** (*auto simp*: *bequiv_up_to_def*)
qed (*fastforce simp*: *equiv_up_to_def bequiv_up_to_def*)+

lemma *equiv_up_to_while_weak*:
  assumes *b*: *P* ⊨ *b* <∼> *b′*
  assumes *c*: *P* ⊨ *c* ∼ *c′*
  assumes *I*: ⋀*s s′*. (*c*, *s*) ⇒ *s′* ⟹ *P s* ⟹ *bval b s* ⟹ *P s′*
  shows *P* ⊨ *WHILE b DO c* ∼ *WHILE b′ DO c′*
proof −
  from *b* **have** *b′*: *P* ⊨ *b′* <∼> *b* **by** (*simp add*: *bequiv_up_to_sym*)

  from *c b* **have** *c′*: *P* ⊨ *c′* ∼ *c* **by** (*simp add*: *equiv_up_to_sym*)

  from *I*
  have *I′*: ⋀*s s′*. (*c′*, *s*) ⇒ *s′* ⟹ *P s* ⟹ *bval b′ s* ⟹ *P s′*
    by (*auto dest*!: *equiv_up_toD1* [*OF c′*] *simp*: *bequiv_up_to_subst* [*OF b′*])

  note *equiv_up_to_while_lemma_weak* [*OF _ b c*]
       equiv_up_to_while_lemma_weak [*OF _ b′ c′*]
  thus *?thesis* **using** *I I′* **by** (*auto intro*!: *equiv_up_toI*)
qed

lemma *equiv_up_to_if_weak*:
  P ⊨ *b* <∼> *b′* ⟹ *P* ⊨ *c* ∼ *c′* ⟹ *P* ⊨ *d* ∼ *d′* ⟹
   P ⊨ *IF b THEN c ELSE d* ∼ *IF b′ THEN c′ ELSE d′*
  by (*auto simp*: *bequiv_up_to_def equiv_up_to_def*)

lemma *equiv_up_to_if_True* [*intro*!]:
  (⋀*s*. *P s* ⟹ *bval b s*) ⟹ *P* ⊨ *IF b THEN c1 ELSE c2* ∼ *c1*
  by (*auto simp*: *equiv_up_to_def*)

lemma *equiv_up_to_if_False* [*intro*!]:
  (⋀*s*. *P s* ⟹ ¬ *bval b s*) ⟹ *P* ⊨ *IF b THEN c1 ELSE c2* ∼ *c2*
  by (*auto simp*: *equiv_up_to_def*)

lemma *equiv_up_to_while_False* [*intro*!]:
  (⋀*s*. *P s* ⟹ ¬ *bval b s*) ⟹ *P* ⊨ *WHILE b DO c* ∼ *SKIP*
  by (*auto simp*: *equiv_up_to_def*)

lemma *while_never*: (*c*, *s*) ⇒ *u* ⟹ *c* ≠ *WHILE* (*Bc True*) *DO c′*
 by (*induct rule*: *big_step_induct*) *auto*

lemma *equiv_up_to_while_True* [*intro*!,*simp*]:
  $P \models$ *WHILE Bc True DO c* $\sim$ *WHILE Bc True DO SKIP*
  unfolding *equiv_up_to_def*
  by (*blast dest*: *while_never*)

end
theory *Fold* **imports** *Sem_Equiv Vars* **begin**

11.2 Simple folding of arithmetic expressions

type_synonym
  tab = *vname* ⇒ *val option*

fun *afold* :: *aexp* ⇒ *tab* ⇒ *aexp* **where**
afold (*N n*) _ = *N n* |
afold (*V x*) *t* = (*case t x of None* ⇒ *V x* | *Some k* ⇒ *N k*) |
afold (*Plus e1 e2*) *t* = (*case* (*afold e1 t*, *afold e2 t*) *of*
  (*N n1*, *N n2*) ⇒ *N*(*n1*+*n2*) | (*e1'*,*e2'*) ⇒ *Plus e1' e2'*)

definition *approx t s* ⟷ (∀ *x k*. *t x* = *Some k* ⟶ *s x* = *k*)

theorem *aval_afold*[*simp*]:
assumes *approx t s*
shows *aval* (*afold a t*) *s* = *aval a s*
  using *assms*
  by (*induct a*) (*auto simp*: *approx_def split*: *aexp.split option.split*)

theorem *aval_afold_N*:
assumes *approx t s*
shows *afold a t* = *N n* ⟹ *aval a s* = *n*
  by (*metis assms aval.simps*(*1*) *aval_afold*)

definition
  merge t1 t2 = (λ*m*. *if t1 m* = *t2 m then t1 m else None*)

primrec *defs* :: *com* ⇒ *tab* ⇒ *tab* **where**
defs SKIP t = *t* |
defs (*x* ::= *a*) *t* =
  (*case afold a t of N k* ⇒ *t*(*x* ↦ *k*) | _ ⇒ *t*(*x*:=*None*)) |
defs (*c1*;;*c2*) *t* = (*defs c2 o defs c1*) *t* |
defs (*IF b THEN c1 ELSE c2*) *t* = *merge* (*defs c1 t*) (*defs c2 t*) |
defs (*WHILE b DO c*) *t* = *t* |‘ (−*lvars c*)

```
primrec fold where
fold SKIP _ = SKIP |
fold (x ::= a) t = (x ::= (afold a t)) |
fold (c1;;c2) t = (fold c1 t;; fold c2 (defs c1 t)) |
fold (IF b THEN c1 ELSE c2) t = IF b THEN fold c1 t ELSE fold c2 t |
fold (WHILE b DO c) t = WHILE b DO fold c (t |` (−lvars c))

lemma approx_merge:
  approx t1 s ∨ approx t2 s ⟹ approx (merge t1 t2) s
  by (fastforce simp: merge_def approx_def)

lemma approx_map_le:
  approx t2 s ⟹ t1 ⊆m t2 ⟹ approx t1 s
  by (clarsimp simp: approx_def map_le_def dom_def)

lemma restrict_map_le [intro!, simp]: t |` S ⊆m t
  by (clarsimp simp: restrict_map_def map_le_def)

lemma merge_restrict:
  assumes t1 |` S = t |` S
  assumes t2 |` S = t |` S
  shows merge t1 t2 |` S = t |` S
proof −
  from assms
  have ∀x. (t1 |` S) x = (t |` S) x
   and ∀x. (t2 |` S) x = (t |` S) x by auto
  thus ?thesis
    by (auto simp: merge_def restrict_map_def
             split: if_splits)
qed


lemma defs_restrict:
  defs c t |` (− lvars c) = t |` (− lvars c)
proof (induction c arbitrary: t)
  case (Seq c1 c2)
  hence defs c1 t |` (− lvars c1) = t |` (− lvars c1)
    by simp
  hence defs c1 t |` (− lvars c1) |` (−lvars c2) =
         t |` (− lvars c1) |` (−lvars c2) by simp
  moreover
  from Seq
  have defs c2 (defs c1 t) |` (− lvars c2) =
```

```
      defs c1 t |` (- lvars c2)
    by simp
  hence defs c2 (defs c1 t) |` (- lvars c2) |` (- lvars c1) =
       defs c1 t |` (- lvars c2) |` (- lvars c1)
    by simp
  ultimately
  show ?case by (clarsimp simp: Int_commute)
next
  case (If b c1 c2)
  hence defs c1 t |` (- lvars c1) = t |` (- lvars c1) by simp
  hence defs c1 t |` (- lvars c1) |` (-lvars c2) =
       t |` (- lvars c1) |` (-lvars c2) by simp
  moreover
  from If
  have defs c2 t |` (- lvars c2) = t |` (- lvars c2) by simp
  hence defs c2 t |` (- lvars c2) |` (-lvars c1) =
       t |` (- lvars c2) |` (-lvars c1) by simp
  ultimately
  show ?case by (auto simp: Int_commute intro: merge_restrict)
qed (auto split: aexp.split)


lemma big_step_pres_approx:
  (c,s) ⇒ s' ⟹ approx t s ⟹ approx (defs c t) s'
proof (induction arbitrary: t rule: big_step_induct)
  case Skip thus ?case by simp
next
  case Assign
  thus ?case
    by (clarsimp simp: aval_afold_N approx_def split: aexp.split)
next
  case (Seq c1 s1 s2 c2 s3)
  have approx (defs c1 t) s2 by (rule Seq.IH(1)[OF Seq.prems])
  hence approx (defs c2 (defs c1 t)) s3 by (rule Seq.IH(2))
  thus ?case by simp
next
  case (IfTrue b s c1 s')
  hence approx (defs c1 t) s' by simp
  thus ?case by (simp add: approx_merge)
next
  case (IfFalse b s c2 s')
  hence approx (defs c2 t) s' by simp
  thus ?case by (simp add: approx_merge)
next
```

```
  case WhileFalse
  thus ?case by (simp add: approx_def restrict_map_def)
next
  case (WhileTrue b s1 c s2 s3)
  hence approx (defs c t) s2 by simp
  with WhileTrue
  have approx (defs c t |` (−lvars c)) s3 by simp
  thus ?case by (simp add: defs_restrict)
qed


lemma big_step_pres_approx_restrict:
  (c,s) ⇒ s′ ⟹ approx (t |` (−lvars c)) s ⟹ approx (t |` (−lvars c)) s′
proof (induction arbitrary: t rule: big_step_induct)
  case Assign
  thus ?case by (clarsimp simp: approx_def)
next
  case (Seq c1 s1 s2 c2 s3)
  hence approx (t |` (−lvars c2) |` (−lvars c1)) s1
    by (simp add: Int_commute)
  hence approx (t |` (−lvars c2) |` (−lvars c1)) s2
    by (rule Seq)
  hence approx (t |` (−lvars c1) |` (−lvars c2)) s2
    by (simp add: Int_commute)
  hence approx (t |` (−lvars c1) |` (−lvars c2)) s3
    by (rule Seq)
  thus ?case by simp
next
  case (IfTrue b s c1 s′ c2)
  hence approx (t |` (−lvars c2) |` (−lvars c1)) s
    by (simp add: Int_commute)
  hence approx (t |` (−lvars c2) |` (−lvars c1)) s′
    by (rule IfTrue)
  thus ?case by (simp add: Int_commute)
next
  case (IfFalse b s c2 s′ c1)
  hence approx (t |` (−lvars c1) |` (−lvars c2)) s
    by simp
  hence approx (t |` (−lvars c1) |` (−lvars c2)) s′
    by (rule IfFalse)
  thus ?case by simp
qed auto
```

```
declare assign_simp [simp]

lemma approx_eq:
  approx t ⊨ c ∼ fold c t
proof (induction c arbitrary: t)
  case SKIP show ?case by simp
next
  case Assign
  show ?case by (simp add: equiv_up_to_def)
next
  case Seq
  thus ?case by (auto intro!: equiv_up_to_seq big_step_pres_approx)
next
  case If
  thus ?case by (auto intro!: equiv_up_to_if_weak)
next
  case (While b c)
  hence approx (t |` (− lvars c)) ⊨
        WHILE b DO c ∼ WHILE b DO fold c (t |` (− lvars c))
   by (auto intro: equiv_up_to_while_weak big_step_pres_approx_restrict)
  thus ?case
    by (auto intro: equiv_up_to_weaken approx_map_le)
qed


lemma approx_empty [simp]:
  approx Map.empty = (λ_. True)
  by (auto simp: approx_def)


theorem constant_folding_equiv:
  fold c Map.empty ∼ c
  using approx_eq [of Map.empty c]
  by (simp add: equiv_up_to_True sim_sym)


end
```

12 Live Variable Analysis

```
theory Live imports Vars Big_Step
begin
```

12.1 Liveness Analysis

fun $L :: com \Rightarrow vname\ set \Rightarrow vname\ set$ **where**
$L\ SKIP\ X = X\ |$
$L\ (x ::= a)\ X = vars\ a \cup (X - \{x\})\ |$
$L\ (c_1;;\ c_2)\ X = L\ c_1\ (L\ c_2\ X)\ |$
$L\ (IF\ b\ THEN\ c_1\ ELSE\ c_2)\ X = vars\ b \cup L\ c_1\ X \cup L\ c_2\ X\ |$
$L\ (WHILE\ b\ DO\ c)\ X = vars\ b \cup X \cup L\ c\ X$

value $show\ (L\ (''y'' ::= V\ ''z'';;\ ''x'' ::= Plus\ (V\ ''y'')\ (V\ ''z''))\ \{''x''\})$

value $show\ (L\ (WHILE\ Less\ (V\ ''x'')\ (V\ ''x'')\ DO\ ''y'' ::= V\ ''z'')\ \{''x''\})$

fun $kill :: com \Rightarrow vname\ set$ **where**
$kill\ SKIP = \{\}\ |$
$kill\ (x ::= a) = \{x\}\ |$
$kill\ (c_1;;\ c_2) = kill\ c_1 \cup kill\ c_2\ |$
$kill\ (IF\ b\ THEN\ c_1\ ELSE\ c_2) = kill\ c_1 \cap kill\ c_2\ |$
$kill\ (WHILE\ b\ DO\ c) = \{\}$

fun $gen :: com \Rightarrow vname\ set$ **where**
$gen\ SKIP = \{\}\ |$
$gen\ (x ::= a) = vars\ a\ |$
$gen\ (c_1;;\ c_2) = gen\ c_1 \cup (gen\ c_2 - kill\ c_1)\ |$
$gen\ (IF\ b\ THEN\ c_1\ ELSE\ c_2) = vars\ b \cup gen\ c_1 \cup gen\ c_2\ |$
$gen\ (WHILE\ b\ DO\ c) = vars\ b \cup gen\ c$

lemma *L_gen_kill*: $L\ c\ X = gen\ c \cup (X - kill\ c)$
by(*induct c arbitrary*:*X*) *auto*

lemma *L_While_pfp*: $L\ c\ (L\ (WHILE\ b\ DO\ c)\ X) \subseteq L\ (WHILE\ b\ DO\ c)$
X
by(*auto simp add*:*L_gen_kill*)

lemma *L_While_lpfp*:
$vars\ b \cup X \cup L\ c\ P \subseteq P \Longrightarrow L\ (WHILE\ b\ DO\ c)\ X \subseteq P$
by(*simp add*: *L_gen_kill*)

lemma *L_While_vars*: $vars\ b \subseteq L\ (WHILE\ b\ DO\ c)\ X$
by *auto*

lemma *L_While_X*: $X \subseteq L\ (WHILE\ b\ DO\ c)\ X$
by *auto*

Disable L WHILE equation and reason only with L WHILE constraints

declare *L.simps*(*5*)[*simp del*]

12.2 Correctness

theorem *L_correct*:
  $(c,s) \Rightarrow s' \Longrightarrow s = t$ *on* $L\ c\ X \Longrightarrow$
  $\exists\ t'.\ (c,t) \Rightarrow t'\ \&\ s' = t'$ *on* X
proof (*induction arbitrary*: *X t rule*: *big_step_induct*)
  case *Skip* **then show** *?case* **by** *auto*
next
  case *Assign* **then show** *?case*
    by (*auto simp*: *ball_Un*)
next
  case (*Seq c1 s1 s2 c2 s3 X t1*)
  from *Seq.IH*(*1*) *Seq.prems* **obtain** *t2* **where**
    t12: (*c1*, *t1*) ⇒ *t2* **and** *s2t2*: *s2* = *t2 on L c2 X*
    by *simp blast*
  from *Seq.IH*(*2*)[*OF s2t2*] **obtain** *t3* **where**
    t23: (*c2*, *t2*) ⇒ *t3* **and** *s3t3*: *s3* = *t3 on X*
    by *auto*
  show *?case* **using** *t12 t23 s3t3* **by** *auto*
next
  case (*IfTrue b s c1 s′ c2*)
  hence *s* = *t on vars b s* = *t on L c1 X* **by** *auto*
  from *bval_eq_if_eq_on_vars*[*OF this*(*1*)] *IfTrue*(*1*) **have** *bval b t* **by** *simp*
  from *IfTrue.IH*[*OF* ‹*s* = *t on L c1 X*›] **obtain** *t′* **where**
    (*c1*, *t*) ⇒ *t′ s′* = *t′ on X* **by** *auto*
  thus *?case* **using** ‹*bval b t*› **by** *auto*
next
  case (*IfFalse b s c2 s′ c1*)
  hence *s* = *t on vars b s* = *t on L c2 X* **by** *auto*
  from *bval_eq_if_eq_on_vars*[*OF this*(*1*)] *IfFalse*(*1*) **have** ~*bval b t* **by** *simp*
  from *IfFalse.IH*[*OF* ‹*s* = *t on L c2 X*›] **obtain** *t′* **where**
    (*c2*, *t*) ⇒ *t′ s′* = *t′ on X* **by** *auto*
  thus *?case* **using** ‹~*bval b t*› **by** *auto*
next
  case (*WhileFalse b s c*)
  hence ~ *bval b t*
    by (*metis L_While_vars bval_eq_if_eq_on_vars subsetD*)
  thus *?case* **by**(*metis WhileFalse.prems L_While_X big_step.WhileFalse subsetD*)
next

case (*WhileTrue b s1 c s2 s3 X t1*)
let *?w* = *WHILE b DO c*
from ‹*bval b s1*› *WhileTrue.prems* **have** *bval b t1*
by (*metis L_While_vars bval_eq_if_eq_on_vars subsetD*)
have *s1* = *t1 on L c* (*L ?w X*) **using** *L_While_pfp WhileTrue.prems*
by (*blast*)
from *WhileTrue.IH*(*1*)[*OF this*] **obtain** *t2* **where**
(*c*, *t1*) ⇒ *t2 s2* = *t2 on L ?w X* **by** *auto*
from *WhileTrue.IH*(*2*)[*OF this*(*2*)] **obtain** *t3* **where** (*?w*,*t2*) ⇒ *t3 s3* = *t3 on X*
by *auto*
with ‹*bval b t1*› ‹(*c*, *t1*) ⇒ *t2*› **show** *?case* **by** *auto*
qed

12.3 Program Optimization

Burying assignments to dead variables:

fun *bury* :: *com* ⇒ *vname set* ⇒ *com* **where**
bury SKIP X = *SKIP* |
bury (*x* ::= *a*) *X* = (*if x* ∈ *X then x* ::= *a else SKIP*) |
bury (c_1;; c_2) *X* = (*bury* c_1 (*L* c_2 *X*);; *bury* c_2 *X*) |
bury (*IF b THEN* c_1 *ELSE* c_2) *X* = *IF b THEN bury* c_1 *X ELSE bury* c_2 *X* |
bury (*WHILE b DO c*) *X* = *WHILE b DO bury c* (*L* (*WHILE b DO c*) *X*)

We could prove the analogous lemma to *L_correct*, and the proof would be very similar. However, we phrase it as a semantics preservation property:

theorem *bury_correct*:
(*c*,*s*) ⇒ *s'* ⟹ *s* = *t on L c X* ⟹
∃ *t'*. (*bury c X*,*t*) ⇒ *t'* & *s'* = *t' on X*
proof (*induction arbitrary*: *X t rule*: *big_step_induct*)
case *Skip* **then show** *?case* **by** *auto*
next
case *Assign* **then show** *?case*
by (*auto simp*: *ball_Un*)
next
case (*Seq c1 s1 s2 c2 s3 X t1*)
from *Seq.IH*(*1*) *Seq.prems* **obtain** *t2* **where**
t12: (*bury c1* (*L c2 X*), *t1*) ⇒ *t2* **and** *s2t2*: *s2* = *t2 on L c2 X*
by *simp blast*
from *Seq.IH*(*2*)[*OF s2t2*] **obtain** *t3* **where**
t23: (*bury c2 X*, *t2*) ⇒ *t3* **and** *s3t3*: *s3* = *t3 on X*
by *auto*
show *?case* **using** *t12 t23 s3t3* **by** *auto*

next
  case (*IfTrue b s c1 s' c2*)
  hence *s = t on vars b s = t on L c1 X* **by** *auto*
  from *bval_eq_if_eq_on_vars*[*OF this*(*1*)] *IfTrue*(*1*) **have** *bval b t* **by** *simp*
  from *IfTrue.IH*[*OF* ‹*s = t on L c1 X*›] **obtain** *t'* **where**
    (*bury c1 X*, *t*) ⇒ *t' s' =t' on X* **by** *auto*
  thus *?case* **using** ‹*bval b t*› **by** *auto*
next
  case (*IfFalse b s c2 s' c1*)
  hence *s = t on vars b s = t on L c2 X* **by** *auto*
  from *bval_eq_if_eq_on_vars*[*OF this*(*1*)] *IfFalse*(*1*) **have** ~*bval b t* **by** *simp*
  from *IfFalse.IH*[*OF* ‹*s = t on L c2 X*›] **obtain** *t'* **where**
    (*bury c2 X*, *t*) ⇒ *t' s' = t' on X* **by** *auto*
  thus *?case* **using** ‹~*bval b t*› **by** *auto*
next
  case (*WhileFalse b s c*)
  hence ~ *bval b t* **by** (*metis L_While_vars bval_eq_if_eq_on_vars subsetD*)
  thus *?case*
    by *simp* (*metis L_While_X WhileFalse.prems big_step.WhileFalse subsetD*)
next
  case (*WhileTrue b s1 c s2 s3 X t1*)
  let *?w = WHILE b DO c*
  from ‹*bval b s1*› *WhileTrue.prems* **have** *bval b t1*
    by (*metis L_While_vars bval_eq_if_eq_on_vars subsetD*)
  have *s1 = t1 on L c* (*L ?w X*)
    using *L_While_pfp WhileTrue.prems* **by** *blast*
  from *WhileTrue.IH*(*1*)[*OF this*] **obtain** *t2* **where**
    (*bury c* (*L ?w X*), *t1*) ⇒ *t2 s2 = t2 on L ?w X* **by** *auto*
  from *WhileTrue.IH*(*2*)[*OF this*(*2*)] **obtain** *t3*
    where (*bury ?w X*,*t2*) ⇒ *t3 s3 = t3 on X*
    by *auto*
  with ‹*bval b t1*› ‹(*bury c* (*L ?w X*), *t1*) ⇒ *t2*› **show** *?case* **by** *auto*
qed

corollary *final_bury_correct*: (*c*,*s*) ⇒ *s'* ⟹ (*bury c UNIV*,*s*) ⇒ *s'*
using *bury_correct*[*of c s s' UNIV*]
by (*auto simp*: *fun_eq_iff*[*symmetric*])

Now the opposite direction.

lemma *SKIP_bury*[*simp*]:

$SKIP = bury\ c\ X \longleftrightarrow c = SKIP \mid (\exists x\ a.\ c = x{::=}a\ \&\ x \notin X)$
by (*cases c*) *auto*

lemma *Assign_bury*[*simp*]: $x{::=}a = bury\ c\ X \longleftrightarrow c = x{::=}a \land x \in X$
by (*cases c*) *auto*

lemma *Seq_bury*[*simp*]: $bc_1;;bc_2 = bury\ c\ X \longleftrightarrow$
$(\exists c_1\ c_2.\ c = c_1;;c_2\ \&\ bc_2 = bury\ c_2\ X\ \&\ bc_1 = bury\ c_1\ (L\ c_2\ X))$
by (*cases c*) *auto*

lemma *If_bury*[*simp*]: *IF b THEN bc1 ELSE bc2* = *bury c X* $\longleftrightarrow$
($\exists$ *c1 c2*. *c* = *IF b THEN c1 ELSE c2* &
bc1 = *bury c1 X* & *bc2* = *bury c2 X*)
by (*cases c*) *auto*

lemma *While_bury*[*simp*]: *WHILE b DO bc'* = *bury c X* $\longleftrightarrow$
($\exists$ *c'*. *c* = *WHILE b DO c'* & *bc'* = *bury c'* (*L* (*WHILE b DO c'*) *X*))
by (*cases c*) *auto*

theorem *bury_correct2*:
(*bury c X*,*s*) $\Rightarrow$ *s'* $\Longrightarrow$ *s* = *t on L c X* $\Longrightarrow$
$\exists$ *t'*. (*c*,*t*) $\Rightarrow$ *t'* & *s'* = *t' on X*
proof (*induction bury c X s s' arbitrary*: *c X t rule*: *big_step_induct*)
case *Skip* **then show** *?case* **by** *auto*
next
case *Assign* **then show** *?case*
by (*auto simp*: *ball_Un*)
next
case (*Seq bc1 s1 s2 bc2 s3 c X t1*)
then obtain *c1 c2* **where** *c*: *c* = *c1*;;*c2*
and *bc2*: *bc2* = *bury c2 X* **and** *bc1*: *bc1* = *bury c1* (*L c2 X*) **by** *auto*
note *IH* = *Seq.hyps*(*2*,*4*)
from *IH*(*1*)[*OF bc1*, *of t1*] *Seq.prems c* **obtain** *t2* **where**
t12: (*c1*, *t1*) $\Rightarrow$ *t2* **and** *s2t2*: *s2* = *t2 on L c2 X* **by** *auto*
from *IH*(*2*)[*OF bc2 s2t2*] **obtain** *t3* **where**
t23: (*c2*, *t2*) $\Rightarrow$ *t3* **and** *s3t3*: *s3* = *t3 on X*
by *auto*
show *?case* **using** *c t12 t23 s3t3* **by** *auto*
next
case (*IfTrue b s bc1 s' bc2*)
then obtain *c1 c2* **where** *c*: *c* = *IF b THEN c1 ELSE c2*
and *bc1*: *bc1* = *bury c1 X* **and** *bc2*: *bc2* = *bury c2 X* **by** *auto*
have *s* = *t on vars b s* = *t on L c1 X* **using** *IfTrue.prems c* **by** *auto*
from *bval_eq_if_eq_on_vars*[*OF this*(*1*)] *IfTrue*(*1*) **have** *bval b t* **by**

```
simp
  note IH = IfTrue.hyps(3)
  from IH[OF bc1 ‹s = t on L c1 X›] obtain t' where
    (c1, t) ⇒ t' s' =t' on X by auto
  thus ?case using c ‹bval b t› by auto
next
  case (IfFalse b s bc2 s' bc1)
  then obtain c1 c2 where c: c = IF b THEN c1 ELSE c2
    and bc1: bc1 = bury c1 X and bc2: bc2 = bury c2 X by auto
  have s = t on vars b s = t on L c2 X using IfFalse.prems c by auto
  from bval_eq_if_eq_on_vars[OF this(1)] IfFalse(1) have ~bval b t by
simp
  note IH = IfFalse.hyps(3)
  from IH[OF bc2 ‹s = t on L c2 X›] obtain t' where
    (c2, t) ⇒ t' s' =t' on X by auto
  thus ?case using c ‹~bval b t› by auto
next
  case (WhileFalse b s c)
  hence ~ bval b t
    by auto (metis L_While_vars bval_eq_if_eq_on_vars rev_subsetD)
  thus ?case using WhileFalse
    by auto (metis L_While_X big_step.WhileFalse subsetD)
next
  case (WhileTrue b s1 bc' s2 s3 w X t1)
  then obtain c' where w: w = WHILE b DO c'
    and bc': bc' = bury c' (L (WHILE b DO c') X) by auto
  from ‹bval b s1› WhileTrue.prems w have bval b t1
    by auto (metis L_While_vars bval_eq_if_eq_on_vars subsetD)
  note IH = WhileTrue.hyps(3,5)
  have s1 = t1 on L c' (L w X)
    using L_While_pfp WhileTrue.prems w by blast
  with IH(1)[OF bc', of t1] w obtain t2 where
    (c', t1) ⇒ t2 s2 = t2 on L w X by auto
  from IH(2)[OF WhileTrue.hyps(6), of t2] w this(2) obtain t3
    where (w,t2) ⇒ t3 s3 = t3 on X
    by auto
  with ‹bval b t1› ‹(c', t1) ⇒ t2› w show ?case by auto
qed

corollary final_bury_correct2: (bury c UNIV,s) ⇒ s' ⟹ (c,s) ⇒ s'
using bury_correct2[of c UNIV]
by (auto simp: fun_eq_iff[symmetric])

corollary bury_sim: bury c UNIV ∼ c
```

by(*metis final_bury_correct final_bury_correct2*)

end

12.4 True Liveness Analysis

theory *Live_True*
imports *HOL−Library.While_Combinator Vars Big_Step*
begin

12.4.1 Analysis

```
fun L :: com ⇒ vname set ⇒ vname set where
L SKIP X = X |
L (x ::= a) X = (if x ∈ X then vars a ∪ (X − {x}) else X) |
L (c1;; c2) X = L c1 (L c2 X) |
L (IF b THEN c1 ELSE c2) X = vars b ∪ L c1 X ∪ L c2 X |
L (WHILE b DO c) X = lfp(λY. vars b ∪ X ∪ L c Y)

lemma L_mono: mono (L c)
proof−
  have X ⊆ Y ⟹ L c X ⊆ L c Y for X Y
  proof(induction c arbitrary: X Y)
    case (While b c)
    show ?case
    proof(simp, rule lfp_mono)
      fix Z show vars b ∪ X ∪ L c Z ⊆ vars b ∪ Y ∪ L c Z
        using While by auto
    qed
  next
    case If thus ?case by(auto simp: subset_iff)
  qed auto
  thus ?thesis by(rule monoI)
qed

lemma mono_union_L:
  mono (λY. X ∪ L c Y)
by (metis (no_types) L_mono mono_def order_eq_iff set_eq_subset sup_mono)

lemma L_While_unfold:
  L (WHILE b DO c) X = vars b ∪ X ∪ L c (L (WHILE b DO c) X)
by(metis lfp_unfold[OF mono_union_L] L.simps(5))

lemma L_While_pfp: L c (L (WHILE b DO c) X) ⊆ L (WHILE b DO c)
```

X
using *L_While_unfold* **by** *blast*

lemma *L_While_vars*: *vars b* ⊆ *L* (*WHILE b DO c*) *X*
using *L_While_unfold* **by** *blast*

lemma *L_While_X*: *X* ⊆ *L* (*WHILE b DO c*) *X*
using *L_While_unfold* **by** *blast*

Disable *L WHILE* equation and reason only with *L WHILE* constraints:

declare *L.simps*(*5*)[*simp del*]

12.4.2 Correctness

theorem *L_correct*:
  (*c*,*s*) ⇒ *s′* ⟹ *s* = *t on L c X* ⟹
  ∃ *t′*. (*c*,*t*) ⇒ *t′* & *s′* = *t′ on X*
proof (*induction arbitrary*: *X t rule*: *big_step_induct*)
  case *Skip* **then show** *?case* **by** *auto*
next
  case *Assign* **then show** *?case*
    by (*auto simp*: *ball_Un*)
next
  case (*Seq c1 s1 s2 c2 s3 X t1*)
  from *Seq.IH*(*1*) *Seq.prems* **obtain** *t2* **where**
    t12: (*c1*, *t1*) ⇒ *t2* **and** *s2t2*: *s2* = *t2 on L c2 X*
    by *simp blast*
  from *Seq.IH*(*2*)[*OF s2t2*] **obtain** *t3* **where**
    t23: (*c2*, *t2*) ⇒ *t3* **and** *s3t3*: *s3* = *t3 on X*
    by *auto*
  show *?case* **using** *t12 t23 s3t3* **by** *auto*
next
  case (*IfTrue b s c1 s′ c2*)
  hence *s* = *t on vars b* **and** *s* = *t on L c1 X* **by** *auto*
  from *bval_eq_if_eq_on_vars*[*OF this*(*1*)] *IfTrue*(*1*) **have** *bval b t* **by** *simp*
  from *IfTrue.IH*[*OF* ‹*s* = *t on L c1 X*›] **obtain** *t′* **where**
    (*c1*, *t*) ⇒ *t′ s′* = *t′ on X* **by** *auto*
  thus *?case* **using** ‹*bval b t*› **by** *auto*
next
  case (*IfFalse b s c2 s′ c1*)
  hence *s* = *t on vars b s* = *t on L c2 X* **by** *auto*
  from *bval_eq_if_eq_on_vars*[*OF this*(*1*)] *IfFalse*(*1*) **have** ~*bval b t* **by** *simp*

from *IfFalse.IH*[*OF* ‹*s* = *t* *on* *L* *c2* *X*›] **obtain** *t'* **where**
  (*c2*, *t*) ⇒ *t'* *s'* = *t'* *on* *X* **by** *auto*
thus *?case* **using** ‹~*bval* *b* *t*› **by** *auto*
next
  case (*WhileFalse* *b* *s* *c*)
  hence ~ *bval* *b* *t*
    by (*metis* *L_While_vars* *bval_eq_if_eq_on_vars* *subsetD*)
  thus *?case* **using** *WhileFalse.prems* *L_While_X*[*of* *X* *b* *c*] **by** *auto*
next
  case (*WhileTrue* *b* *s1* *c* *s2* *s3* *X* *t1*)
  let *?w* = *WHILE* *b* *DO* *c*
  from ‹*bval* *b* *s1*› *WhileTrue.prems* **have** *bval* *b* *t1*
    by (*metis* *L_While_vars* *bval_eq_if_eq_on_vars* *subsetD*)
  have *s1* = *t1* *on* *L* *c* (*L* *?w* *X*) **using** *L_While_pfp* *WhileTrue.prems*
    by (*blast*)
  from *WhileTrue.IH*(*1*)[*OF* *this*] **obtain** *t2* **where**
    (*c*, *t1*) ⇒ *t2* *s2* = *t2* *on* *L* *?w* *X* **by** *auto*
  from *WhileTrue.IH*(*2*)[*OF* *this*(*2*)] **obtain** *t3* **where** (*?w*,*t2*) ⇒ *t3* *s3* = *t3* *on* *X*
    by *auto*
  with ‹*bval* *b* *t1*› ‹(*c*, *t1*) ⇒ *t2*› **show** *?case* **by** *auto*
qed

12.4.3 Executability

lemma *L_subset_vars*: *L* *c* *X* ⊆ *rvars* *c* ∪ *X*
proof(*induction* *c* *arbitrary*: *X*)
  case (*While* *b* *c*)
  have *lfp*(λ*Y*. *vars* *b* ∪ *X* ∪ *L* *c* *Y*) ⊆ *vars* *b* ∪ *rvars* *c* ∪ *X*
    using *While.IH*[*of* *vars* *b* ∪ *rvars* *c* ∪ *X*]
    by (*auto* *intro*!: *lfp_lowerbound*)
  thus *?case* **by** (*simp* *add*: *L.simps*(*5*))
qed *auto*

Make *L* executable by replacing *lfp* with the *while* combinator from theory *HOL−Library.While_Combinator*. The *while* combinator obeys the recursion equation

while *b* *c* *s* = (*if* *b* *s* *then* *while* *b* *c* (*c* *s*) *else* *s*)

and is thus executable.

lemma *L_While*: **fixes** *b* *c* *X*
assumes *finite* *X* **defines** *f* == λ*Y*. *vars* *b* ∪ *X* ∪ *L* *c* *Y*
shows *L* (*WHILE* *b* *DO* *c*) *X* = *while* (λ*Y*. *f* *Y* ≠ *Y*) *f* {} (**is** _ = *?r*)
proof −

let *?V* = *vars b* ∪ *rvars c* ∪ *X*
have *lfp f* = *?r*
proof(*rule lfp_while*[**where** *C* = *?V*])
show *mono f* **by**(*simp add*: *f_def mono_union_L*)
next
fix *Y* **show** *Y* ⊆ *?V* ⟹ *f Y* ⊆ *?V*
unfolding *f_def* **using** *L_subset_vars*[*of c*] **by** *blast*
next
show *finite ?V* **using** ‹*finite X*› **by** *simp*
qed
thus *?thesis* **by** (*simp add*: *f_def L.simps*(*5*))
qed

lemma *L_While_let*: *finite X* ⟹ *L* (*WHILE b DO c*) *X* =
(*let f* = (λ*Y*. *vars b* ∪ *X* ∪ *L c Y*)
in while (λ*Y*. *f Y* ≠ *Y*) *f* {})
by(*simp add*: *L_While*)

lemma *L_While_set*: *L* (*WHILE b DO c*) (*set xs*) =
(*let f* = (λ*Y*. *vars b* ∪ *set xs* ∪ *L c Y*)
in while (λ*Y*. *f Y* ≠ *Y*) *f* {})
by(*rule L_While_let*, *simp*)

Replace the equation for *L* (*WHILE* ...) by the executable *L_While_set*:

lemmas [*code*] = *L.simps*(*1*−*4*) *L_While_set*

Sorry, this syntax is odd.

A test:

lemma (*let b* = *Less* (*N 0*) (*V* ''*y*''); *c* = ''*y*'' ::= *V* ''*x*'';; ''*x*'' ::= *V* ''*z*''
in L (*WHILE b DO c*) {''*y*''}) = {''*x*'', ''*y*'', ''*z*''}
by *eval*

12.4.4 Limiting the number of iterations

The final parameter is the default value:

fun *iter* :: (′*a* ⇒ ′*a*) ⇒ *nat* ⇒ ′*a* ⇒ ′*a* ⇒ ′*a* **where**
iter f 0 p d = *d* |
iter f (*Suc n*) *p d* = (*if f p* = *p then p else iter f n* (*f p*) *d*)

A version of *L* with a bounded number of iterations (here: 2) in the WHILE case:

fun *Lb* :: *com* ⇒ *vname set* ⇒ *vname set* **where**
Lb SKIP X = *X* |
Lb (*x* ::= *a*) *X* = (*if x* ∈ *X then X* − {*x*} ∪ *vars a else X*) |

$Lb\ (c_1;;\ c_2)\ X = (Lb\ c_1 \circ Lb\ c_2)\ X\ |$
$Lb\ (IF\ b\ THEN\ c_1\ ELSE\ c_2)\ X = vars\ b \cup Lb\ c_1\ X \cup Lb\ c_2\ X\ |$
$Lb\ (WHILE\ b\ DO\ c)\ X = iter\ (\lambda A.\ vars\ b \cup X \cup Lb\ c\ A)\ 2\ \{\}\ (vars\ b \cup rvars\ c \cup X)$

Lb (and *iter*) is not monotone!

lemma *let w = WHILE Bc False DO (''x'' ::= V ''y'';; ''z'' ::= V ''x'')*
in ¬ *(Lb w {''z''}* ⊆ *Lb w {''y'',''z''})*
by *eval*

lemma *lfp_subset_iter*:
⟦ *mono f*; !!*X*. *f X* ⊆ *f' X*; *lfp f* ⊆ *D* ⟧ ⟹ *lfp f* ⊆ *iter f' n A D*
proof(*induction n arbitrary*: *A*)
case *0* **thus** *?case* **by** *simp*
next
case *Suc* **thus** *?case* **by** *simp* (*metis lfp_lowerbound*)
qed

lemma *L c X* ⊆ *Lb c X*
proof(*induction c arbitrary*: *X*)
case (*While b c*)
let *?f* = λ*A*. *vars b* ∪ *X* ∪ *L c A*
let *?fb* = λ*A*. *vars b* ∪ *X* ∪ *Lb c A*
show *?case*
proof (*simp add*: *L.simps*(*5*), *rule lfp_subset_iter*[*OF mono_union_L*])
show !!*X*. *?f X* ⊆ *?fb X* **using** *While.IH* **by** *blast*
show *lfp ?f* ⊆ *vars b* ∪ *rvars c* ∪ *X*
by (*metis* (*full_types*) *L.simps*(*5*) *L_subset_vars rvars.simps*(*5*))
qed
next
case *Seq* **thus** *?case* **by** *simp* (*metis* (*full_types*) *L_mono monoD subset_trans*)
qed *auto*

end

13 Hoare Logic

13.1 Hoare Logic for Partial Correctness

theory *Hoare* **imports** *Big_Step* **begin**

type_synonym *assn* = *state* ⇒ *bool*

definition
hoare_valid :: *assn* ⇒ *com* ⇒ *assn* ⇒ *bool* (⊨ {(*1*_)}/ (_)/ {(*1*_)} *50*)
where
$\models \{P\}c\{Q\} = (\forall s\ t.\ P\ s \wedge (c,s) \Rightarrow t \longrightarrow Q\ t)$

abbreviation *state_subst* :: *state* ⇒ *aexp* ⇒ *vname* ⇒ *state*
(_[_'/_] [*1000*,*0*,*0*] *999*)
where $s[a/x] == s(x := aval\ a\ s)$

inductive
hoare :: *assn* ⇒ *com* ⇒ *assn* ⇒ *bool* (⊢ ({(*1*_)}/ (_)/ {(*1*_)}) *50*)
where
Skip: $\vdash \{P\}\ SKIP\ \{P\}$ |

Assign: $\vdash \{\lambda s.\ P(s[a/x])\}\ x::=a\ \{P\}$ |

Seq: $[\![\vdash \{P\}\ c_1\ \{Q\};\ \vdash \{Q\}\ c_2\ \{R\}]\!]$
$\Longrightarrow \vdash \{P\}\ c_1;;c_2\ \{R\}$ |

If: $[\![\vdash \{\lambda s.\ P\ s \wedge bval\ b\ s\}\ c_1\ \{Q\};\ \vdash \{\lambda s.\ P\ s \wedge \neg\ bval\ b\ s\}\ c_2\ \{Q\}]\!]$
$\Longrightarrow \vdash \{P\}\ IF\ b\ THEN\ c_1\ ELSE\ c_2\ \{Q\}$ |

While: $\vdash \{\lambda s.\ P\ s \wedge bval\ b\ s\}\ c\ \{P\} \Longrightarrow$
$\vdash \{P\}\ WHILE\ b\ DO\ c\ \{\lambda s.\ P\ s \wedge \neg\ bval\ b\ s\}$ |

conseq: $[\![\forall s.\ P'\ s \longrightarrow P\ s;\ \vdash \{P\}\ c\ \{Q\};\ \forall s.\ Q\ s \longrightarrow Q'\ s]\!]$
$\Longrightarrow \vdash \{P'\}\ c\ \{Q'\}$

lemmas [*simp*] = *hoare.Skip hoare.Assign hoare.Seq If*

lemmas [*intro*!] = *hoare.Skip hoare.Assign hoare.Seq hoare.If*

lemma *strengthen_pre*:
$[\![\forall s.\ P'\ s \longrightarrow P\ s;\ \vdash \{P\}\ c\ \{Q\}]\!] \Longrightarrow \vdash \{P'\}\ c\ \{Q\}$
by (*blast intro*: *conseq*)

lemma *weaken_post*:
$[\![\vdash \{P\}\ c\ \{Q\};\ \forall s.\ Q\ s \longrightarrow Q'\ s]\!] \Longrightarrow \vdash \{P\}\ c\ \{Q'\}$
by (*blast intro*: *conseq*)

The assignment and While rule are awkward to use in actual proofs because their pre and postcondition are of a very special form and the actual goal would have to match this form exactly. Therefore we derive two variants with arbitrary pre and postconditions.

lemma *Assign'*: $\forall s.\ P\ s \longrightarrow Q(s[a/x]) \Longrightarrow \vdash \{P\}\ x ::= a\ \{Q\}$
by (*simp add*: *strengthen_pre*[*OF* _ *Assign*])

lemma *While'*:
assumes $\vdash \{\lambda s.\ P\ s \wedge bval\ b\ s\}\ c\ \{P\}$ **and** $\forall s.\ P\ s \wedge \neg\ bval\ b\ s \longrightarrow Q\ s$
shows $\vdash \{P\}\ WHILE\ b\ DO\ c\ \{Q\}$
by(*rule weaken_post*[*OF While*[*OF assms*(*1*)] *assms*(*2*)])

end

13.2 Examples

theory *Hoare_Examples* **imports** *Hoare* **begin**

hide_const (**open**) *sum*

Summing up the first x natural numbers in variable y.

fun *sum* :: *int* ⇒ *int* **where**
sum i = (*if i* ≤ *0 then 0 else sum* (*i* − *1*) + *i*)

lemma *sum_simps*[*simp*]:
  0 < *i* ⟹ *sum i* = *sum* (*i* − *1*) + *i*
  i ≤ *0* ⟹ *sum i* = *0*
by(*simp_all*)

declare *sum.simps*[*simp del*]

abbreviation *wsum* ==
  WHILE Less (*N 0*) (*V* ''*x*'')
  DO (''*y*'' ::= *Plus* (*V* ''*y*'') (*V* ''*x*'');;
      ''*x*'' ::= *Plus* (*V* ''*x*'') (*N* (− *1*)))

13.2.1 Proof by Operational Semantics

The behaviour of the loop is proved by induction:

lemma *while_sum*:
  (*wsum*, *s*) ⇒ *t* ⟹ *t* ''*y*'' = *s* ''*y*'' + *sum*(*s* ''*x*'')
apply(*induction wsum s t rule*: *big_step_induct*)
apply(*auto*)
done

We were lucky that the proof was automatic, except for the induction. In general, such proofs will not be so easy. The automation is partly due to the right inversion rules that we set up as automatic elimination rules that decompose big-step premises.

Now we prefix the loop with the necessary initialization:

```
lemma sum_via_bigstep:
  assumes (''y'' ::= N 0;; wsum, s) ⇒ t
  shows t ''y'' = sum (s ''x'')
proof –
  from assms have (wsum,s(''y'':=0)) ⇒ t by auto
  from while_sum[OF this] show ?thesis by simp
qed
```

13.2.2 Proof by Hoare Logic

Note that we deal with sequences of commands from right to left, pulling back the postcondition towards the precondition.

```
lemma ⊢ {λs. s ''x'' = n} ''y'' ::= N 0;; wsum {λs. s ''y'' = sum n}
apply(rule Seq)
 prefer 2
 apply(rule While' [where P = λs. (s ''y'' = sum n − sum(s ''x''))])
  apply(rule Seq)
   prefer 2
   apply(rule Assign)
  apply(rule Assign')
  apply simp
 apply simp
apply(rule Assign')
apply simp
done
```

The proof is intentionally an apply script because it merely composes the rules of Hoare logic. Of course, in a few places side conditions have to be proved. But since those proofs are 1-liners, a structured proof is overkill. In fact, we shall learn later that the application of the Hoare rules can be automated completely and all that is left for the user is to provide the loop invariants and prove the side-conditions.

end

13.3 Soundness and Completeness

```
theory Hoare_Sound_Complete
imports Hoare
begin
```

13.3.1 Soundness

lemma *hoare_sound*: $\vdash \{P\}c\{Q\} \implies \models \{P\}c\{Q\}$

proof(*induction rule*: *hoare.induct*)
  case (*While P b c*)
  have (*WHILE b DO c*,*s*) ⇒ *t* ⟹ *P s* ⟹ *P t* ∧ ¬ *bval b t* **for** *s t*
  proof(*induction WHILE b DO c s t rule*: *big_step_induct*)
    case *WhileFalse* **thus** *?case* **by** *blast*
  next
    case *WhileTrue* **thus** *?case*
      using *While.IH* **unfolding** *hoare_valid_def* **by** *blast*
  qed
  thus *?case* **unfolding** *hoare_valid_def* **by** *blast*
qed (*auto simp*: *hoare_valid_def*)

13.3.2 Weakest Precondition

definition *wp* :: *com* ⇒ *assn* ⇒ *assn* **where**
wp c Q = (λ*s*. ∀ *t*. (*c*,*s*) ⇒ *t* ⟶ *Q t*)

lemma *wp_SKIP*[*simp*]: *wp SKIP Q* = *Q*
by (*rule ext*) (*auto simp*: *wp_def*)

lemma *wp_Ass*[*simp*]: *wp* (*x*::=*a*) *Q* = (λ*s*. *Q*(*s*[*a*/*x*]))
by (*rule ext*) (*auto simp*: *wp_def*)

lemma *wp_Seq*[*simp*]: *wp* (c_1;;c_2) *Q* = *wp* c_1 (*wp* c_2 *Q*)
by (*rule ext*) (*auto simp*: *wp_def*)

lemma *wp_If*[*simp*]:
 wp (*IF b THEN* c_1 *ELSE* c_2) *Q* =
 (λ*s*. *if bval b s then wp* c_1 *Q s else wp* c_2 *Q s*)
by (*rule ext*) (*auto simp*: *wp_def*)

lemma *wp_While_If*:
 wp (*WHILE b DO c*) *Q s* =
  wp (*IF b THEN c*;;*WHILE b DO c ELSE SKIP*) *Q s*
unfolding *wp_def* **by** (*metis unfold_while*)

lemma *wp_While_True*[*simp*]: *bval b s* ⟹
  wp (*WHILE b DO c*) *Q s* = *wp* (*c*;; *WHILE b DO c*) *Q s*
by(*simp add*: *wp_While_If*)

lemma *wp_While_False*[*simp*]: ¬ *bval b s* ⟹ *wp* (*WHILE b DO c*) *Q s*
= *Q s*
by(*simp add*: *wp_While_If*)

13.3.3 Completeness

lemma *wp_is_pre*: $\vdash \{wp\ c\ Q\}\ c\ \{Q\}$
proof(*induction c arbitrary*: *Q*)
  case *If* **thus** *?case* **by**(*auto intro*: *conseq*)
next
  case (*While b c*)
  let *?w* = *WHILE b DO c*
  show $\vdash \{wp\ ?w\ Q\}\ ?w\ \{Q\}$
  proof(*rule While'*)
    show $\vdash \{\lambda s.\ wp\ ?w\ Q\ s \wedge bval\ b\ s\}\ c\ \{wp\ ?w\ Q\}$
    proof(*rule strengthen_pre*[*OF _ While.IH*])
      show $\forall s.\ wp\ ?w\ Q\ s \wedge bval\ b\ s \longrightarrow wp\ c\ (wp\ ?w\ Q)\ s$ **by** *auto*
    qed
    show $\forall s.\ wp\ ?w\ Q\ s \wedge \neg\ bval\ b\ s \longrightarrow Q\ s$ **by** *auto*
  qed
qed *auto*

lemma *hoare_complete*: **assumes** $\models \{P\}c\{Q\}$ **shows** $\vdash \{P\}c\{Q\}$
proof(*rule strengthen_pre*)
  show $\forall s.\ P\ s \longrightarrow wp\ c\ Q\ s$ **using** *assms*
    by (*auto simp*: *hoare_valid_def wp_def*)
  show $\vdash \{wp\ c\ Q\}\ c\ \{Q\}$ **by**(*rule wp_is_pre*)
qed

corollary *hoare_sound_complete*: $\vdash \{P\}c\{Q\} \longleftrightarrow\ \models \{P\}c\{Q\}$
by (*metis hoare_complete hoare_sound*)

end

13.4 Verification Condition Generation

theory *VCG* **imports** *Hoare* **begin**

13.4.1 Annotated Commands

Commands where loops are annotated with invariants.

datatype *acom* =
  Askip   (*SKIP*) |
  Aassign vname aexp   ((_ ::= _) [*1000*, *61*] *61*) |
  Aseq   acom acom   (_;;/ _ [*60*, *61*] *60*) |
  Aif bexp acom acom   ((*IF* _/ *THEN* _/ *ELSE* _) [*0*, *0*, *61*] *61*) |
  Awhile assn bexp acom (({_}/ *WHILE* _/ *DO* _) [*0*, *0*, *61*] *61*)

notation *com.SKIP* (*SKIP*)

Strip annotations:

fun *strip* :: *acom* ⇒ *com* **where**
strip SKIP = *SKIP* |
strip (*x* ::= *a*) = (*x* ::= *a*) |
strip (C_1;; C_2) = (*strip* C_1;; *strip* C_2) |
strip (*IF b THEN* C_1 *ELSE* C_2) = (*IF b THEN strip* C_1 *ELSE strip* C_2) |
strip ({_} *WHILE b DO C*) = (*WHILE b DO strip C*)

13.4.2 Weeakest Precondistion and Verification Condition

Weakest precondition:

fun *pre* :: *acom* ⇒ *assn* ⇒ *assn* **where**
pre SKIP Q = *Q* |
pre (*x* ::= *a*) *Q* = (λ*s*. *Q*(*s*(*x* := *aval a s*))) |
pre (C_1;; C_2) *Q* = *pre* C_1 (*pre* C_2 *Q*) |
pre (*IF b THEN* C_1 *ELSE* C_2) *Q* =
  (λ*s*. *if bval b s then pre* C_1 *Q s else pre* C_2 *Q s*) |
pre ({*I*} *WHILE b DO C*) *Q* = *I*

Verification condition:

fun *vc* :: *acom* ⇒ *assn* ⇒ *bool* **where**
vc SKIP Q = *True* |
vc (*x* ::= *a*) *Q* = *True* |
vc (C_1;; C_2) *Q* = (*vc* C_1 (*pre* C_2 *Q*) ∧ *vc* C_2 *Q*) |
vc (*IF b THEN* C_1 *ELSE* C_2) *Q* = (*vc* C_1 *Q* ∧ *vc* C_2 *Q*) |
vc ({*I*} *WHILE b DO C*) *Q* =
  ((∀ *s*. (*I s* ∧ *bval b s* ⟶ *pre C I s*) ∧
      (*I s* ∧ ¬ *bval b s* ⟶ *Q s*)) ∧
    vc C I)

13.4.3 Soundness

lemma *vc_sound*: *vc C Q* ⟹ ⊢ {*pre C Q*} *strip C* {*Q*}
proof(*induction C arbitrary*: *Q*)
  case (*Awhile I b C*)
  show *?case*
  proof(*simp*, *rule While′*)
    from ‹*vc* (*Awhile I b C*) *Q*›
    have *vc*: *vc C I* **and** *IQ*: ∀ *s*. *I s* ∧ ¬ *bval b s* ⟶ *Q s* **and**
        pre: ∀ *s*. *I s* ∧ *bval b s* ⟶ *pre C I s* **by** *simp_all*
    have ⊢ {*pre C I*} *strip C* {*I*} **by**(*rule Awhile.IH*[*OF vc*])
    with *pre* **show** ⊢ {λ*s*. *I s* ∧ *bval b s*} *strip C* {*I*}

by(*rule strengthen_pre*)
show ∀ *s*. *I s* ∧ ¬*bval b s* ⟶ *Q s* **by**(*rule IQ*)
qed
qed (*auto intro*: *hoare.conseq*)

corollary *vc_sound′*:
⟦ *vc C Q*; ∀ *s*. *P s* ⟶ *pre C Q s* ⟧ ⟹ ⊢ {*P*} *strip C* {*Q*}
by (*metis strengthen_pre vc_sound*)

13.4.4 Completeness

lemma *pre_mono*:
∀ *s*. *P s* ⟶ *P′ s* ⟹ *pre C P s* ⟹ *pre C P′ s*
proof (*induction C arbitrary*: *P P′ s*)
case *Aseq* **thus** *?case* **by** *simp metis*
qed *simp_all*

lemma *vc_mono*:
∀ *s*. *P s* ⟶ *P′ s* ⟹ *vc C P* ⟹ *vc C P′*
proof(*induction C arbitrary*: *P P′*)
case *Aseq* **thus** *?case* **by** *simp* (*metis pre_mono*)
qed *simp_all*

lemma *vc_complete*:
⊢ {*P*}*c*{*Q*} ⟹ ∃ *C*. *strip C* = *c* ∧ *vc C Q* ∧ (∀ *s*. *P s* ⟶ *pre C Q s*)
(**is** _ ⟹ ∃ *C*. *?G P c Q C*)
proof (*induction rule*: *hoare.induct*)
case *Skip*
show *?case* (**is** ∃ *C*. *?C C*)
proof show *?C Askip* **by** *simp* **qed**
next
case (*Assign P a x*)
show *?case* (**is** ∃ *C*. *?C C*)
proof show *?C*(*Aassign x a*) **by** *simp* **qed**
next
case (*Seq P c1 Q c2 R*)
from *Seq.IH* **obtain** *C1* **where** *ih1*: *?G P c1 Q C1* **by** *blast*
from *Seq.IH* **obtain** *C2* **where** *ih2*: *?G Q c2 R C2* **by** *blast*
show *?case* (**is** ∃ *C*. *?C C*)
proof
show *?C*(*Aseq C1 C2*)
using *ih1 ih2* **by** (*fastforce elim*!: *pre_mono vc_mono*)
qed
next

case (*If P b c1 Q c2*)
from *If.IH* **obtain** *C1* **where** *ih1*: *?G* ($\lambda s.$ *P s* $\wedge$ *bval b s*) *c1 Q C1*
by *blast*
from *If.IH* **obtain** *C2* **where** *ih2*: *?G* ($\lambda s.$ *P s* $\wedge$ $\neg$*bval b s*) *c2 Q C2*
by *blast*
show *?case* (**is** $\exists$ *C*. *?C C*)
proof
show *?C*(*Aif b C1 C2*) **using** *ih1 ih2* **by** *simp*
qed
next
case (*While P b c*)
from *While.IH* **obtain** *C* **where** *ih*: *?G* ($\lambda s.$ *P s* $\wedge$ *bval b s*) *c P C* **by** *blast*
show *?case* (**is** $\exists$ *C*. *?C C*)
proof show *?C*(*Awhile P b C*) **using** *ih* **by** *simp* **qed**
next
case *conseq* **thus** *?case* **by**(*fast elim*!: *pre_mono vc_mono*)
qed

end

13.5 Hoare Logic for Total Correctness

13.5.1 Separate Termination Relation

theory *Hoare_Total*
imports *Hoare_Examples*
begin

Note that this definition of total validity $\models_t$ only works if execution is deterministic (which it is in our case).

definition *hoare_tvalid* :: *assn* $\Rightarrow$ *com* $\Rightarrow$ *assn* $\Rightarrow$ *bool*
($\models_t$ {(*1*_)}/ (_)/ {(*1*_)} *50*) **where**
$\models_t$ {*P*}*c*{*Q*} $\longleftrightarrow$ ($\forall$ *s*. *P s* $\longrightarrow$ ($\exists$ *t*. (*c*,*s*) $\Rightarrow$ *t* $\wedge$ *Q t*))

Provability of Hoare triples in the proof system for total correctness is written $\vdash_t$ {*P*}*c*{*Q*} and defined inductively. The rules for $\vdash_t$ differ from those for $\vdash$ only in the one place where nontermination can arise: the *While*-rule.

inductive
hoaret :: *assn* $\Rightarrow$ *com* $\Rightarrow$ *assn* $\Rightarrow$ *bool* ($\vdash_t$ ({(*1*_)}/ (_)/ {(*1*_)}) *50*)
where

Skip: $\vdash_t$ {*P*} *SKIP* {*P*} |

Assign: $\vdash_t \{\lambda s.\ P(s[a/x])\}\ x{::=}a\ \{P\}\ \ |$

Seq: $[\![\ \vdash_t \{P_1\}\ c_1\ \{P_2\};\ \vdash_t \{P_2\}\ c_2\ \{P_3\}\ ]\!] \Longrightarrow \vdash_t \{P_1\}\ c_1;;c_2\ \{P_3\}\ \ |$

If: $[\![\ \vdash_t \{\lambda s.\ P\ s \wedge bval\ b\ s\}\ c_1\ \{Q\};\ \vdash_t \{\lambda s.\ P\ s \wedge \neg\ bval\ b\ s\}\ c_2\ \{Q\}\ ]\!]$
$\Longrightarrow \vdash_t \{P\}\ IF\ b\ THEN\ c_1\ ELSE\ c_2\ \{Q\}\ \ |$

While:
$(\bigwedge n{::}nat.$
$\vdash_t \{\lambda s.\ P\ s \wedge bval\ b\ s \wedge T\ s\ n\}\ c\ \{\lambda s.\ P\ s \wedge (\exists\, n'{<}n.\ T\ s\ n')\})$
$\Longrightarrow \vdash_t \{\lambda s.\ P\ s \wedge (\exists\, n.\ T\ s\ n)\}\ WHILE\ b\ DO\ c\ \{\lambda s.\ P\ s \wedge \neg bval\ b\ s\}\ \ |$

conseq: $[\![\ \forall\, s.\ P'\ s \longrightarrow P\ s;\ \vdash_t \{P\}c\{Q\};\ \forall\, s.\ Q\ s \longrightarrow Q'\ s\ \ ]\!] \Longrightarrow$
$\vdash_t \{P'\}c\{Q'\}$

The *While*-rule is like the one for partial correctness but it requires additionally that with every execution of the loop body some measure relation $T :: state \Rightarrow nat \Rightarrow bool$ decreases. The following functional version is more intuitive:

lemma *While_fun*:
$[\![\ \bigwedge n{::}nat.\ \vdash_t \{\lambda s.\ P\ s \wedge bval\ b\ s \wedge n = f\ s\}\ c\ \{\lambda s.\ P\ s \wedge f\ s < n\}]\!]$
$\Longrightarrow \vdash_t \{P\}\ WHILE\ b\ DO\ c\ \{\lambda s.\ P\ s \wedge \neg bval\ b\ s\}$
by (*rule While* [**where** $T{=}\lambda s\ n.\ n = f\ s$, *simplified*])

Building in the consequence rule:

lemma *strengthen_pre*:
$[\![\ \forall\, s.\ P'\ s \longrightarrow P\ s;\ \ \vdash_t \{P\}\ c\ \{Q\}\ ]\!] \Longrightarrow \vdash_t \{P'\}\ c\ \{Q\}$
by (*metis conseq*)

lemma *weaken_post*:
$[\![\ \vdash_t \{P\}\ c\ \{Q\};\ \ \forall\, s.\ Q\ s \longrightarrow Q'\ s\ ]\!] \Longrightarrow\ \vdash_t \{P\}\ c\ \{Q'\}$
by (*metis conseq*)

lemma *Assign'*: $\forall\, s.\ P\ s \longrightarrow Q(s[a/x]) \Longrightarrow \vdash_t \{P\}\ x ::= a\ \{Q\}$
by (*simp add*: *strengthen_pre*[*OF _ Assign*])

lemma *While_fun'*:
assumes $\bigwedge n{::}nat.\ \vdash_t \{\lambda s.\ P\ s \wedge bval\ b\ s \wedge n = f\ s\}\ c\ \{\lambda s.\ P\ s \wedge f\ s < n\}$
and $\forall\, s.\ P\ s \wedge \neg\ bval\ b\ s \longrightarrow Q\ s$
shows $\vdash_t \{P\}\ WHILE\ b\ DO\ c\ \{Q\}$
by(*blast intro*: *assms*(*1*) *weaken_post*[*OF While_fun assms*(*2*)])

Our standard example:

lemma $\vdash_t \{\lambda s.\ s\ ''x'' = i\}\ ''y'' ::= N\ 0;;\ wsum\ \{\lambda s.\ s\ ''y'' = sum\ i\}$

```
apply(rule Seq)
 prefer 2
 apply(rule While_fun' [where P = λs. (s ''y'' = sum i − sum(s ''x''))
    and f = λs. nat(s ''x'')])
   apply(rule Seq)
   prefer 2
   apply(rule Assign)
  apply(rule Assign')
  apply simp
 apply(simp)
apply(rule Assign')
apply simp
done
```

Nested loops. This poses a problem for VCGs because the proof of the inner loop needs to refer to outer loops. This works here because the invariant is not written down statically but created in the context of a proof that has already introduced/fixed outer ns that can be referred to.

```
lemma
 ⊢t {λ_. True}
  WHILE Less (N 0) (V ''x'')
  DO (''x'' ::= Plus (V ''x'') (N(−1));;
      ''y'' ::= V ''x'';;
      WHILE Less (N 0) (V ''y'') DO ''y'' ::= Plus (V ''y'') (N(−1)))
  {λ_. True}
apply(rule While_fun'[where f = λs. nat(s ''x'')])
 prefer 2 apply simp
apply(rule_tac P2 = λs. nat(s ''x'') < n in Seq)
 apply(rule_tac P2 = λs. nat(s ''x'') < n in Seq)
  apply(rule Assign')
  apply simp
 apply(rule Assign')
 apply simp

apply(rule While_fun'[where f = λs. nat(s ''y'')])
 prefer 2 apply simp
apply(rule Assign')
apply simp
done
```

The soundness theorem:

```
theorem hoaret_sound: ⊢t {P}c{Q} ⟹ ⊨t {P}c{Q}
proof(unfold hoare_tvalid_def, induction rule: hoaret.induct)
  case (While P b T c)
```

have $\llbracket$ *P s*; *T s n* $\rrbracket \Longrightarrow \exists t.$ (*WHILE b DO c*, *s*) $\Rightarrow$ *t* $\wedge$ *P t* $\wedge$ $\neg$ *bval b t*
for *s n*
proof(*induction n arbitrary*: *s rule*: *less_induct*)
case (*less n*) **thus** *?case* **by** (*metis While.IH WhileFalse WhileTrue*)
qed
thus *?case* **by** *auto*
next
case *If* **thus** *?case* **by** *auto blast*
qed *fastforce*+

The completeness proof proceeds along the same lines as the one for partial correctness. First we have to strengthen our notion of weakest precondition to take termination into account:

definition *wpt* :: *com* $\Rightarrow$ *assn* $\Rightarrow$ *assn* (wp_t) **where**
wp_t *c Q* = ($\lambda s.\ \exists t.$ (*c*,*s*) $\Rightarrow$ *t* $\wedge$ *Q t*)

lemma [*simp*]: wp_t *SKIP Q* = *Q*
by(*auto intro*!: *ext simp*: *wpt_def*)

lemma [*simp*]: wp_t (*x* ::= *e*) *Q* = ($\lambda s.$ *Q*(*s*(*x* := *aval e s*)))
by(*auto intro*!: *ext simp*: *wpt_def*)

lemma [*simp*]: wp_t (c_1;;c_2) *Q* = wp_t c_1 (wp_t c_2 *Q*)
unfolding *wpt_def*
apply(*rule ext*)
apply *auto*
done

lemma [*simp*]:
wp_t (*IF b THEN* c_1 *ELSE* c_2) *Q* = ($\lambda s.$ wp_t (*if bval b s then* c_1 *else* c_2) *Q s*)
apply(*unfold wpt_def*)
apply(*rule ext*)
apply *auto*
done

Now we define the number of iterations *WHILE b DO c* needs to terminate when started in state *s*. Because this is a truly partial function, we define it as an (inductive) relation first:

inductive *Its* :: *bexp* $\Rightarrow$ *com* $\Rightarrow$ *state* $\Rightarrow$ *nat* $\Rightarrow$ *bool* **where**
Its_0: $\neg$ *bval b s* $\Longrightarrow$ *Its b c s 0* |
Its_Suc: $\llbracket$ *bval b s*; (*c*,*s*) $\Rightarrow$ s'; *Its b c* s' *n* $\rrbracket \Longrightarrow$ *Its b c s* (*Suc n*)

The relation is in fact a function:

lemma *Its_fun*: *Its b c s n* $\Longrightarrow$ *Its b c s* n' $\Longrightarrow$ $n=n'$

proof(*induction arbitrary*: *n′ rule*:*Its.induct*)
  case *Its_0* **thus** *?case* **by**(*metis Its.cases*)
next
  case *Its_Suc* **thus** *?case* **by**(*metis Its.cases big_step_determ*)
qed

For all terminating loops, *Its* yields a result:

lemma *WHILE_Its*: (*WHILE b DO c*,*s*) ⇒ *t* ⟹ ∃ *n*. *Its b c s n*
proof(*induction WHILE b DO c s t rule*: *big_step_induct*)
  case *WhileFalse* **thus** *?case* **by** (*metis Its_0*)
next
  case *WhileTrue* **thus** *?case* **by** (*metis Its_Suc*)
qed

lemma *wpt_is_pre*: $\vdash_t$ {wp_t *c Q*} *c* {*Q*}
proof (*induction c arbitrary*: *Q*)
  case *SKIP* **show** *?case* **by** (*auto intro*:*hoaret.Skip*)
next
  case *Assign* **show** *?case* **by** (*auto intro*:*hoaret.Assign*)
next
  case *Seq* **thus** *?case* **by** (*auto intro*:*hoaret.Seq*)
next
  case *If* **thus** *?case* **by** (*auto intro*:*hoaret.If hoaret.conseq*)
next
  case (*While b c*)
  let *?w* = *WHILE b DO c*
  let *?T* = *Its b c*
  have *1*: ∀ *s*. wp_t *?w Q s* ⟶ wp_t *?w Q s* ∧ (∃ *n*. *Its b c s n*)
    unfolding *wpt_def* **by** (*metis WHILE_Its*)
  let *?R* = λ*n s′*. wp_t *?w Q s′* ∧ (∃ *n′*<*n*. *?T s′ n′*)
  have ∀ *s*. wp_t *?w Q s* ∧ *bval b s* ∧ *?T s n* ⟶ wp_t *c* (*?R n*) *s* **for** *n*
  proof −
    have wp_t *c* (*?R n*) *s* **if** *bval b s* **and** *?T s n* **and** (*?w*, *s*) ⇒ *t* **and** *Q t*
for *s t*
    proof −
      from ‹*bval b s*› **and** ‹(*?w*, *s*) ⇒ *t*› **obtain** *s′* **where**
        (*c*,*s*) ⇒ *s′* (*?w*,*s′*) ⇒ *t* **by** *auto*
      from ‹(*?w*, *s′*) ⇒ *t*› **obtain** *n′* **where** *?T s′ n′*
        by (*blast dest*: *WHILE_Its*)
      with ‹*bval b s*› **and** ‹(*c*, *s*) ⇒ *s′*› **have** *?T s* (*Suc n′*) **by** (*rule Its_Suc*)
      with ‹*?T s n*› **have** *n* = *Suc n′* **by** (*rule Its_fun*)
      with ‹(*c*,*s*) ⇒ *s′*› **and** ‹(*?w*,*s′*) ⇒ *t*› **and** ‹*Q t*› **and** ‹*?T s′ n′*›
      show *?thesis* **by** (*auto simp*: *wpt_def*)
    qed

thus *?thesis*
unfolding *wpt_def* **by** *auto*

qed
note *2* = *hoaret.While*[*OF strengthen_pre*[*OF this While.IH*]]
have $\forall$ *s*. wp_t *?w Q s* $\wedge$ $\neg$ *bval b s* $\longrightarrow$ *Q s*
by (*auto simp add*:*wpt_def*)
with *1 2* **show** *?case* **by** (*rule conseq*)
qed

In the *While*-case, *Its* provides the obvious termination argument.

The actual completeness theorem follows directly, in the same manner as for partial correctness:

theorem *hoaret_complete*: $\models_t$ {*P*}*c*{*Q*} $\Longrightarrow$ $\vdash_t$ {*P*}*c*{*Q*}
apply(*rule strengthen_pre*[*OF _ wpt_is_pre*])
apply(*auto simp*: *hoare_tvalid_def wpt_def*)
done

corollary *hoaret_sound_complete*: $\vdash_t$ {*P*}*c*{*Q*} $\longleftrightarrow$ $\models_t$ {*P*}*c*{*Q*}
by (*metis hoaret_sound hoaret_complete*)

end

13.5.2 *nat*-Indexed Invariant

theory *Hoare_Total_EX*
imports *Hoare*
begin

This is the standard set of rules that you find in many publications. The While-rule is different from the one in Concrete Semantics in that the invariant is indexed by natural numbers and goes down by 1 with every iteration. The completeness proof is easier but the rule is harder to apply in program proofs.

definition *hoare_tvalid* :: *assn* $\Rightarrow$ *com* $\Rightarrow$ *assn* $\Rightarrow$ *bool*
($\models_t$ {(*1_*)}/ (_)/ {(*1_*)} *50*) **where**
$\models_t$ {*P*}*c*{*Q*} $\longleftrightarrow$ ($\forall$ *s*. *P s* $\longrightarrow$ ($\exists$ *t*. (*c*,*s*) $\Rightarrow$ *t* $\wedge$ *Q t*))

inductive
hoaret :: *assn* $\Rightarrow$ *com* $\Rightarrow$ *assn* $\Rightarrow$ *bool* ($\vdash_t$ ({(*1_*)}/ (_)/ {(*1_*)}) *50*)
where

Skip: $\vdash_t$ {*P*} *SKIP* {*P*} |

Assign: $\vdash_t \{\lambda s.\ P(s[a/x])\}\ x{::=}a\ \{P\}$ |

Seq: $[\![\vdash_t \{P_1\}\ c_1\ \{P_2\};\ \vdash_t \{P_2\}\ c_2\ \{P_3\}]\!] \Longrightarrow \vdash_t \{P_1\}\ c_1;;c_2\ \{P_3\}$ |

If: $[\![\vdash_t \{\lambda s.\ P\ s \wedge bval\ b\ s\}\ c_1\ \{Q\};\ \vdash_t \{\lambda s.\ P\ s \wedge \neg\ bval\ b\ s\}\ c_2\ \{Q\}]\!]$
$\Longrightarrow \vdash_t \{P\}\ IF\ b\ THEN\ c_1\ ELSE\ c_2\ \{Q\}$ |

While:
$[\![\bigwedge n{::}nat.\ \vdash_t \{P\ (Suc\ n)\}\ c\ \{P\ n\};$
$\forall n\ s.\ P\ (Suc\ n)\ s \longrightarrow bval\ b\ s;\ \ \forall s.\ P\ 0\ s \longrightarrow \neg\ bval\ b\ s]\!]$
$\Longrightarrow \vdash_t \{\lambda s.\ \exists n.\ P\ n\ s\}\ WHILE\ b\ DO\ c\ \{P\ 0\}$ |

conseq: $[\![\forall s.\ P'\ s \longrightarrow P\ s;\ \vdash_t \{P\}c\{Q\};\ \forall s.\ Q\ s \longrightarrow Q'\ s\ ]\!] \Longrightarrow$
$\vdash_t \{P'\}c\{Q'\}$

Building in the consequence rule:

lemma *strengthen_pre*:
$[\![\forall s.\ P'\ s \longrightarrow P\ s;\ \ \vdash_t \{P\}\ c\ \{Q\}]\!] \Longrightarrow \vdash_t \{P'\}\ c\ \{Q\}$
by (*metis conseq*)

lemma *weaken_post*:
$[\![\vdash_t \{P\}\ c\ \{Q\};\ \ \forall s.\ Q\ s \longrightarrow Q'\ s]\!] \Longrightarrow\ \vdash_t \{P\}\ c\ \{Q'\}$
by (*metis conseq*)

lemma *Assign'*: $\forall s.\ P\ s \longrightarrow Q(s[a/x]) \Longrightarrow \vdash_t \{P\}\ x ::= a\ \{Q\}$
by (*simp add*: *strengthen_pre*[*OF _ Assign*])

The soundness theorem:

theorem *hoaret_sound*: $\vdash_t \{P\}c\{Q\}\ \Longrightarrow\ \models_t \{P\}c\{Q\}$
proof(*unfold hoare_tvalid_def*, *induction rule*: *hoaret.induct*)
  case (*While P c b*)
  have $P\ n\ s \Longrightarrow \exists t.\ (WHILE\ b\ DO\ c,\ s) \Rightarrow t \wedge P\ 0\ t$ **for** *n s*
  proof(*induction n arbitrary*: *s*)
    case *0* **thus** *?case* **using** *While.hyps*(*3*) *WhileFalse* **by** *blast*
  next
    case *Suc*
    thus *?case* **by** (*meson While.IH While.hyps*(*2*) *WhileTrue*)
  qed
  thus *?case* **by** *auto*
next
  case *If* **thus** *?case* **by** *auto blast*
qed *fastforce+*

definition *wpt* :: *com* ⇒ *assn* ⇒ *assn* (wp_t) **where**
wp_t *c Q* = (λ*s*. ∃ *t*. (*c*,*s*) ⇒ *t* ∧ *Q t*)

lemma [*simp*]: wp_t *SKIP Q* = *Q*
by(*auto intro*!: *ext simp*: *wpt_def*)

lemma [*simp*]: wp_t (*x* ::= *e*) *Q* = (λ*s*. *Q*(*s*(*x* := *aval e s*)))
by(*auto intro*!: *ext simp*: *wpt_def*)

lemma [*simp*]: wp_t (c_1;;c_2) *Q* = wp_t c_1 (wp_t c_2 *Q*)
unfolding *wpt_def*
apply(*rule ext*)
apply *auto*
done

lemma [*simp*]:
wp_t (*IF b THEN* c_1 *ELSE* c_2) *Q* = (λ*s*. wp_t (*if bval b s then* c_1 *else* c_2) *Q* *s*)
apply(*unfold wpt_def*)
apply(*rule ext*)
apply *auto*
done

Function *wpw* computes the weakest precondition of a While-loop that is unfolded a fixed number of times.

fun *wpw* :: *bexp* ⇒ *com* ⇒ *nat* ⇒ *assn* ⇒ *assn* **where**
wpw b c 0 Q s = (¬ *bval b s* ∧ *Q s*) |
wpw b c (*Suc n*) *Q s* = (*bval b s* ∧ (∃ *s'*. (*c*,*s*) ⇒ *s'* ∧ *wpw b c n Q s'*))

lemma *WHILE_Its*: (*WHILE b DO c*,*s*) ⇒ *t* ⟹ *Q t* ⟹ ∃ *n*. *wpw b c n Q s*
proof(*induction WHILE b DO c s t rule*: *big_step_induct*)
case *WhileFalse* **thus** *?case* **using** *wpw.simps*(*1*) **by** *blast*
next
case *WhileTrue* **thus** *?case* **using** *wpw.simps*(*2*) **by** *blast*
qed

lemma *wpt_is_pre*: $\vdash_t$ {wp_t *c Q*} *c* {*Q*}
proof (*induction c arbitrary*: *Q*)
case *SKIP* **show** *?case* **by** (*auto intro*:*hoaret.Skip*)
next
case *Assign* **show** *?case* **by** (*auto intro*:*hoaret.Assign*)
next
case *Seq* **thus** *?case* **by** (*auto intro*:*hoaret.Seq*)

next
  case *If* **thus** *?case* **by** (*auto intro*:*hoaret.If hoaret.conseq*)
next
  case (*While b c*)
  let *?w* = *WHILE b DO c*
  have *c1*: $\forall s.\ wp_t\ ?w\ Q\ s \longrightarrow (\exists n.\ wpw\ b\ c\ n\ Q\ s)$
    unfolding *wpt_def* **by** (*metis WHILE_Its*)
  have *c3*: $\forall s.\ wpw\ b\ c\ 0\ Q\ s \longrightarrow Q\ s$ **by** *simp*
  have *w2*: $\forall n\ s.\ wpw\ b\ c\ (Suc\ n)\ Q\ s \longrightarrow bval\ b\ s$ **by** *simp*
  have *w3*: $\forall s.\ wpw\ b\ c\ 0\ Q\ s \longrightarrow \neg\ bval\ b\ s$ **by** *simp*
  have $\vdash_t \{wpw\ b\ c\ (Suc\ n)\ Q\}\ c\ \{wpw\ b\ c\ n\ Q\}$ **for** *n*
  proof −
    have ∗: $\forall s.\ wpw\ b\ c\ (Suc\ n)\ Q\ s \longrightarrow (\exists t.\ (c, s) \Rightarrow t \land wpw\ b\ c\ n\ Q\ t)$
by *simp*
     show *?thesis* **by**(*rule strengthen_pre*[*OF* ∗ *While.IH*[*of wpw b c n Q*,
unfolded wpt_def]])
  qed
  from *conseq*[*OF c1 hoaret.While*[*OF this w2 w3*] *c3*]
  show *?case* **.**
qed

theorem *hoaret_complete*: $\models_t \{P\}c\{Q\} \Longrightarrow \vdash_t \{P\}c\{Q\}$
apply(*rule strengthen_pre*[*OF _ wpt_is_pre*])
apply(*auto simp*: *hoare_tvalid_def wpt_def*)
done

corollary *hoaret_sound_complete*: $\vdash_t \{P\}c\{Q\} \longleftrightarrow\ \models_t \{P\}c\{Q\}$
by (*metis hoaret_sound hoaret_complete*)

Two examples:

lemma $\vdash_t$
$\{\lambda s.\ \exists n.\ n = nat(s\ ''x'')\}$
 WHILE Less (*N 0*) (*V ''x''*) *DO ''x''* ::= *Plus* (*V ''x''*) (*N* (−*1*))
$\{\lambda s.\ s\ ''x'' \le 0\}$
apply(*rule weaken_post*)
 apply(*rule While*)
   apply(*rule Assign'*)
   apply *auto*
done

lemma $\vdash_t$
$\{\lambda s.\ \exists n.\ n = nat(s\ ''x'')\}$
 WHILE Less (*N 0*) (*V ''x''*)
 DO (*''x''* ::= *Plus* (*V ''x''*) (*N* (−*1*));;

```
  (''y'' ::= V ''x'';;
   WHILE Less (N 0) (V ''y'') DO ''y'' ::= Plus (V ''y'') (N (−1))))
{λs. s ''x'' ≤ 0}
apply(rule weaken_post)
 apply(rule While)
   defer
   apply auto[3]
apply(rule Seq)
 prefer 2
 apply(rule Seq)
  prefer 2
  apply(rule weaken_post)
   apply(rule_tac P = λm s. n = nat(s ''x'') ∧ m = nat(s ''y'') in While)
     apply(rule Assign')
     apply auto[4]
 apply(rule Assign)
apply(rule Assign')
apply auto
done

end
```

13.6 Verification Conditions for Total Correctness

13.6.1 The Standard Approach

```
theory VCG_Total_EX
imports Hoare_Total_EX
begin
```

Annotated commands: commands where loops are annotated with invariants.

```
datatype acom =
  Askip                  (SKIP) |
  Aassign vname aexp     ((_ ::= _) [1000, 61] 61) |
  Aseq   acom acom       (_;;/ _  [60, 61] 60) |
  Aif bexp acom acom     ((IF _/ THEN _/ ELSE _)  [0, 0, 61] 61) |
  Awhile nat ⇒ assn bexp acom
   (({_}/ WHILE _/ DO _)  [0, 0, 61] 61)

notation com.SKIP (SKIP)
```

Strip annotations:

```
fun strip :: acom ⇒ com where
strip SKIP = SKIP |
```

$strip\ (x ::= a) = (x ::= a)\ |$
$strip\ (C_1;;\ C_2) = (strip\ C_1;;\ strip\ C_2)\ |$
$strip\ (IF\ b\ THEN\ C_1\ ELSE\ C_2) = (IF\ b\ THEN\ strip\ C_1\ ELSE\ strip\ C_2)\ |$
$strip\ (\{\_\}\ WHILE\ b\ DO\ C) = (WHILE\ b\ DO\ strip\ C)$

Weakest precondition from annotated commands:

fun $pre :: acom \Rightarrow assn \Rightarrow assn$ **where**
$pre\ SKIP\ Q = Q\ |$
$pre\ (x ::= a)\ Q = (\lambda s.\ Q(s(x := aval\ a\ s)))\ |$
$pre\ (C_1;;\ C_2)\ Q = pre\ C_1\ (pre\ C_2\ Q)\ |$
$pre\ (IF\ b\ THEN\ C_1\ ELSE\ C_2)\ Q =$
$(\lambda s.\ if\ bval\ b\ s\ then\ pre\ C_1\ Q\ s\ else\ pre\ C_2\ Q\ s)\ |$
$pre\ (\{I\}\ WHILE\ b\ DO\ C)\ Q = (\lambda s.\ \exists n.\ I\ n\ s)$

Verification condition:

fun $vc :: acom \Rightarrow assn \Rightarrow bool$ **where**
$vc\ SKIP\ Q = True\ |$
$vc\ (x ::= a)\ Q = True\ |$
$vc\ (C_1;;\ C_2)\ Q = (vc\ C_1\ (pre\ C_2\ Q) \wedge vc\ C_2\ Q)\ |$
$vc\ (IF\ b\ THEN\ C_1\ ELSE\ C_2)\ Q = (vc\ C_1\ Q \wedge vc\ C_2\ Q)\ |$
$vc\ (\{I\}\ WHILE\ b\ DO\ C)\ Q =$
$(\forall s\ n.\ (I\ (Suc\ n)\ s \longrightarrow pre\ C\ (I\ n)\ s) \wedge$
$(I\ (Suc\ n)\ s \longrightarrow bval\ b\ s) \wedge$
$(I\ 0\ s \longrightarrow \neg\ bval\ b\ s \wedge Q\ s) \wedge$
$vc\ C\ (I\ n))$

lemma *vc_sound*: $vc\ C\ Q \Longrightarrow\ \vdash_t \{pre\ C\ Q\}\ strip\ C\ \{Q\}$
proof(*induction C arbitrary*: *Q*)
  case (*Awhile I b C*)
  show *?case*
  proof(*simp*, *rule conseq*[*OF _ While*[*of I*]], *goal_cases*)
    case (*2 n*) **show** *?case*
      using *Awhile.IH*[*of I n*] *Awhile.prems*
      by (*auto intro*: *strengthen_pre*)
  qed (*insert Awhile.prems*, *auto*)
qed (*auto intro*: *conseq Seq If simp*: *Skip Assign*)

When trying to extend the completeness proof of the VCG for partial correctness to total correctness one runs into the following problem. In the case of the while-rule, the universally quantified n in the first premise means that for that premise the induction hypothesis does not yield a single annotated command C but merely that for every n such a C exists.

end

13.6.2 Hoare Logic for Total Correctness With Logical Variables

theory *Hoare_Total_EX2*
imports *Hoare*
begin

This is the standard set of rules that you find in many publications. In the while-rule, a logical variable is needed to remember the pre-value of the variant (an expression that decreases by one with each iteration). In this theory, logical variables are modeled explicitly. A simpler (but not quite as flexible) approach is found in theory *Hoare_Total_EX*: pre and post-condition are connected via a universally quantified HOL variable.

type_synonym *lvname* = *string*
type_synonym *assn2* = (*lvname* ⇒ *nat*) ⇒ *state* ⇒ *bool*

definition *hoare_tvalid* :: *assn2* ⇒ *com* ⇒ *assn2* ⇒ *bool*
 ($\models_t$ {(*1*_)}/ (_)/ {(*1*_)} *50*) **where**
$\models_t \{P\}c\{Q\} \longleftrightarrow (\forall l\ s.\ P\ l\ s \longrightarrow (\exists t.\ (c,s) \Rightarrow t \land Q\ l\ t))$

inductive
 hoaret :: *assn2* ⇒ *com* ⇒ *assn2* ⇒ *bool* ($\vdash_t$ ({(*1*_)}/ (_)/ {(*1*_)}) *50*)
where

Skip: $\vdash_t \{P\}\ SKIP\ \{P\}$ |

Assign: $\vdash_t \{\lambda l\ s.\ P\ l\ (s[a/x])\}\ x::=a\ \{P\}$ |

Seq: $[\![\vdash_t \{P_1\}\ c_1\ \{P_2\}; \vdash_t \{P_2\}\ c_2\ \{P_3\}]\!] \Longrightarrow \vdash_t \{P_1\}\ c_1;;c_2\ \{P_3\}$ |

If: $[\![\vdash_t \{\lambda l\ s.\ P\ l\ s \land bval\ b\ s\}\ c_1\ \{Q\}; \vdash_t \{\lambda l\ s.\ P\ l\ s \land \neg\ bval\ b\ s\}\ c_2\ \{Q\}$
$]\!]$
 $\Longrightarrow \vdash_t \{P\}\ IF\ b\ THEN\ c_1\ ELSE\ c_2\ \{Q\}$ |

While:
 $[\![\vdash_t \{\lambda l.\ P\ (l(x := Suc(l(x))))\}\ c\ \{P\};$
 $\forall l\ s.\ l\ x > 0 \land P\ l\ s \longrightarrow bval\ b\ s;$
 $\forall l\ s.\ l\ x = 0 \land P\ l\ s \longrightarrow \neg\ bval\ b\ s\ ]\!]$
 $\Longrightarrow \vdash_t \{\lambda l\ s.\ \exists n.\ P\ (l(x:=n))\ s\}\ WHILE\ b\ DO\ c\ \{\lambda l\ s.\ P\ (l(x := 0))\ s\}$
|

conseq: $[\![\forall l\ s.\ P'\ l\ s \longrightarrow P\ l\ s; \vdash_t \{P\}c\{Q\}; \forall l\ s.\ Q\ l\ s \longrightarrow Q'\ l\ s\ ]\!] \Longrightarrow$
 $\vdash_t \{P'\}c\{Q'\}$

Building in the consequence rule:

lemma *strengthen_pre*:

$\llbracket\ \forall\, l\ s.\ P'\ l\ s \longrightarrow P\ l\ s;\ \ \vdash_t \{P\}\ c\ \{Q\}\ \rrbracket \Longrightarrow \vdash_t \{P'\}\ c\ \{Q\}$
by (*metis conseq*)

lemma *weaken_post*:
$\llbracket \vdash_t \{P\}\ c\ \{Q\};\ \ \forall\, l\ s.\ Q\ l\ s \longrightarrow Q'\ l\ s\ \rrbracket \Longrightarrow\ \vdash_t \{P\}\ c\ \{Q'\}$
by (*metis conseq*)

lemma *Assign'*: $\forall\, l\ s.\ P\ l\ s \longrightarrow Q\ l\ (s[a/x]) \Longrightarrow \vdash_t \{P\}\ x ::= a\ \{Q\}$
by (*simp add*: *strengthen_pre*[*OF _ Assign*])

The soundness theorem:

theorem *hoaret_sound*: $\vdash_t \{P\}c\{Q\} \implies\ \models_t \{P\}c\{Q\}$
proof(*unfold hoare_tvalid_def*, *induction rule*: *hoaret.induct*)
  case (*While P x c b*)
  have $\llbracket\ l\ x = n;\ P\ l\ s\ \rrbracket \Longrightarrow \exists\, t.\ (WHILE\ b\ DO\ c,\ s) \Rightarrow t \wedge P\ (l(x := 0))\ t$ **for** *n l s*
  proof(*induction n arbitrary*: *l s*)
    case *0* **thus** *?case* **using** *While.hyps*(*3*) *WhileFalse*
      by (*metis fun_upd_triv*)
  next
    case *Suc*
    thus *?case* **using** *While.IH While.hyps*(*2*) *WhileTrue*
      by (*metis fun_upd_same fun_upd_triv fun_upd_upd zero_less_Suc*)
  qed
  thus *?case* **by** *fastforce*
next
  case *If* **thus** *?case* **by** *auto blast*
qed *fastforce+*

definition *wpt* :: *com* ⇒ *assn2* ⇒ *assn2* (wp_t) **where**
$wp_t\ c\ Q\ =\ (\lambda l\ s.\ \exists\, t.\ (c,s) \Rightarrow t \wedge Q\ l\ t)$

lemma [*simp*]: $wp_t\ SKIP\ Q = Q$
by(*auto intro*!: *ext simp*: *wpt_def*)

lemma [*simp*]: $wp_t\ (x ::= e)\ Q = (\lambda l\ s.\ Q\ l\ (s(x := aval\ e\ s)))$
by(*auto intro*!: *ext simp*: *wpt_def*)

lemma *wpt_Seq*[*simp*]: $wp_t\ (c_1;;c_2)\ Q = wp_t\ c_1\ (wp_t\ c_2\ Q)$
by (*auto simp*: *wpt_def fun_eq_iff*)

lemma [*simp*]:
$wp_t\ (IF\ b\ THEN\ c_1\ ELSE\ c_2)\ Q = (\lambda l\ s.\ wp_t\ (if\ bval\ b\ s\ then\ c_1\ else\ c_2)$

Q l s)
by (*auto simp*: *wpt_def fun_eq_iff*)

Function *wpw* computes the weakest precondition of a While-loop that is unfolded a fixed number of times.

fun *wpw* :: *bexp* ⇒ *com* ⇒ *nat* ⇒ *assn2* ⇒ *assn2* **where**
wpw b c 0 Q l s = (¬ *bval b s* ∧ *Q l s*) |
wpw b c (*Suc n*) *Q l s* = (*bval b s* ∧ (∃ *s'*. (*c*,*s*) ⇒ *s'* ∧ *wpw b c n Q l s'*))

lemma *WHILE_Its*:
(*WHILE b DO c*,*s*) ⇒ *t* ⟹ *Q l t* ⟹ ∃ *n*. *wpw b c n Q l s*
proof(*induction WHILE b DO c s t arbitrary*: *l rule*: *big_step_induct*)
case *WhileFalse* **thus** *?case* **using** *wpw.simps*(*1*) **by** *blast*
next
case *WhileTrue* **show** *?case*
using *wpw.simps*(*2*) *WhileTrue*(*1*,*2*) *WhileTrue*(*5*)[*OF WhileTrue*(*6*)]
by *blast*
qed

definition *support* :: *assn2* ⇒ *string set* **where**
support P = {*x*. ∃ *l1 l2 s*. (∀ *y*. *y* ≠ *x* ⟶ *l1 y* = *l2 y*) ∧ *P l1 s* ≠ *P l2 s*}

lemma *support_wpt*: *support* (wp_t *c Q*) ⊆ *support Q*
by(*simp add*: *support_def wpt_def*) *blast*

lemma *support_wpw0*: *support* (*wpw b c n Q*) ⊆ *support Q*
proof(*induction n*)
case *0* **show** *?case* **by** (*simp add*: *support_def*) *blast*
next
case *Suc*
have *1*: *support* (λ*l s*. *A s* ∧ *B l s*) ⊆ *support B* **for** *A B*
by(*auto simp*: *support_def*)
have *2*: *support* (λ*l s*. ∃ *s'*. *A s s'* ∧ *B l s'*) ⊆ *support B* **for** *A B*
by(*auto simp*: *support_def*) *blast*+
from *Suc 1 2* **show** *?case* **by** *simp* (*meson order_trans*)
qed

lemma *support_wpw_Un*:
support (%*l*. *wpw b c* (*l x*) *Q l*) ⊆ *insert x* (*UN n*. *support*(*wpw b c n Q*))
using *support_wpw0*[*of b c _ Q*]
apply(*auto simp add*: *support_def subset_iff*)
apply *metis*
apply *metis*

done

lemma *support_wpw*: *support* (%*l*. *wpw b c* (*l x*) *Q l*) ⊆ *insert x* (*support Q*)
using *support_wpw0*[*of b c _ Q*] *support_wpw_Un*[*of b c _ Q*]
by *blast*

lemma *assn2_lupd*: $x \notin$ *support Q* ⟹ *Q* (*l*(*x*:=*n*)) = *Q l*
by(*simp add*: *support_def fun_upd_other fun_eq_iff*)
(*metis* (*no_types*, *lifting*) *fun_upd_def*)

abbreviation *new Q* ≡ *SOME x*. $x \notin$ *support Q*

lemma *wpw_lupd*: $x \notin$ *support Q* ⟹ *wpw b c n Q* (*l*(*x* := *u*)) = *wpw b c n Q l*
by(*induction n*) (*auto simp*: *assn2_lupd fun_eq_iff*)

lemma *wpt_is_pre*: *finite*(*support Q*) ⟹ $\vdash_t$ {wp_t *c Q*} *c* {*Q*}
proof (*induction c arbitrary*: *Q*)
 case *SKIP* **show** *?case* **by** (*auto intro*:*hoaret.Skip*)
next
 case *Assign* **show** *?case* **by** (*auto intro*:*hoaret.Assign*)
next
 case (*Seq c1 c2*) **show** *?case*
  by (*auto intro*:*hoaret.Seq Seq finite_subset*[*OF support_wpt*])
next
 case *If* **thus** *?case* **by** (*auto intro*:*hoaret.If hoaret.conseq*)
next
 case (*While b c*)
 let *?x* = *new Q*
 have ∃ *x*. $x \notin$ *support Q* **using** *While.prems infinite_UNIV_listI*
  using *ex_new_if_finite* **by** *blast*
 hence [*simp*]: *?x* ∉ *support Q* **by** (*rule someI_ex*)
 let *?w* = *WHILE b DO c*
 have *fsup*: *finite* (*support* (λ*l*. *wpw b c* (*l x*) *Q l*)) **for** *x*
  using *finite_subset*[*OF support_wpw*] *While.prems* **by** *simp*
 have *c1*: ∀ *l s*. wp_t *?w Q l s* ⟶ (∃ *n*. *wpw b c n Q l s*)
  unfolding *wpt_def* **by** (*metis WHILE_Its*)
 have *c2*: ∀ *l s*. *l ?x* = *0* ∧ *wpw b c* (*l ?x*) *Q l s* ⟶ ¬ *bval b s*
  by (*simp cong*: *conj_cong*)
 have *w2*: ∀ *l s*. *0* < *l ?x* ∧ *wpw b c* (*l ?x*) *Q l s* ⟶ *bval b s*
  by (*auto simp*: *gr0_conv_Suc cong*: *conj_cong*)
 have *1*: ∀ *l s*. *wpw b c* (*Suc*(*l ?x*)) *Q l s* ⟶
     (∃ *t*. (*c*, *s*) ⇒ *t* ∧ *wpw b c* (*l ?x*) *Q l t*)

```
    by simp
  have *: ⊢t {λl. wpw b c (Suc (l ?x)) Q l} c {λl. wpw b c (l ?x) Q l}
    by(rule strengthen_pre[OF 1
          While.IH[of λl. wpw b c (l ?x) Q l, unfolded wpt_def, OF fsup]])
  show ?case
  apply(rule conseq[OF _ hoaret.While[OF _ w2 c2]])
    apply (simp_all add: c1 * assn2_lupd wpw_lupd del: wpw.simps(2))
  done
qed

theorem hoaret_complete: finite(support Q) ⟹ ⊨t {P}c{Q} ⟹ ⊢t {P}c{Q}
apply(rule strengthen_pre[OF _ wpt_is_pre])
apply(auto simp: hoare_tvalid_def wpt_def)
done
```

Two examples:

```
lemma ⊢t
{λl s. l ''x'' = nat(s ''x'')}
 WHILE Less (N 0) (V ''x'') DO ''x'' ::= Plus (V ''x'') (N (−1))
{λl s. s ''x'' ≤ 0}
apply(rule conseq)
prefer 2
 apply(rule While[where P = λl s. l ''x'' = nat(s ''x'') and x = ''x''])
   apply(rule Assign')
   apply auto
done

lemma ⊢t
{λl s. l ''x'' = nat(s ''x'')}
 WHILE Less (N 0) (V ''x'')
 DO (''x'' ::= Plus (V ''x'') (N (−1));;
    (''y'' ::= V ''x'';;
     WHILE Less (N 0) (V ''y'') DO ''y'' ::= Plus (V ''y'') (N (−1))))
{λl s. s ''x'' ≤ 0}
apply(rule conseq)
prefer 2
 apply(rule While[where P = λl s. l ''x'' = nat(s ''x'') and x = ''x''])
   defer
   apply auto
apply(rule Seq)
 prefer 2
 apply(rule Seq)
  prefer 2
  apply(rule weaken_post)
```

apply(*rule_tac P* = λ*l s. l ''x'' = nat*(*s ''x''*) $\wedge$ *l ''y'' = nat*(*s ''y''*) **and** *x = ''y''* **in** *While*)
apply(*rule Assign'*)
apply *auto*[*4*]
apply(*rule Assign*)
apply(*rule Assign'*)
apply *auto*
done

end

13.6.3 VCG for Total Correctness With Logical Variables

theory *VCG_Total_EX2*
imports *Hoare_Total_EX2*
begin

Theory *VCG_Total_EX* conatins a VCG built on top of a Hoare logic without logical variables. As a result the completeness proof runs into a problem. This theory uses a Hoare logic with logical variables and proves soundness and completeness.

Annotated commands: commands where loops are annotated with invariants.

datatype *acom* =
Askip (*SKIP*) |
Aassign vname aexp ((_ ::= _) [*1000*, *61*] *61*) |
Aseq acom acom (_;;/ _ [*60*, *61*] *60*) |
Aif bexp acom acom ((*IF* _/ *THEN* _/ *ELSE* _) [*0*, *0*, *61*] *61*) |
Awhile assn2 lvname bexp acom
(({_'/_}/ *WHILE* _/ *DO* _) [*0*, *0*, *0*, *61*] *61*)

notation *com.SKIP* (*SKIP*)

Strip annotations:

fun *strip* :: *acom* $\Rightarrow$ *com* **where**
strip SKIP = SKIP |
strip (*x* ::= *a*) = (*x* ::= *a*) |
strip (C_1;; C_2) = (*strip* C_1;; *strip* C_2) |
strip (*IF b THEN* C_1 *ELSE* C_2) = (*IF b THEN strip* C_1 *ELSE strip* C_2) |
strip ({_/_} *WHILE b DO C*) = (*WHILE b DO strip C*)

Weakest precondition from annotated commands:

fun *pre* :: *acom* $\Rightarrow$ *assn2* $\Rightarrow$ *assn2* **where**
pre SKIP Q = Q |

$pre\ (x ::= a)\ Q = (\lambda l\ s.\ Q\ l\ (s(x := aval\ a\ s)))\ |$
$pre\ (C_1;;\ C_2)\ Q = pre\ C_1\ (pre\ C_2\ Q)\ |$
$pre\ (IF\ b\ THEN\ C_1\ ELSE\ C_2)\ Q =$
$(\lambda l\ s.\ if\ bval\ b\ s\ then\ pre\ C_1\ Q\ l\ s\ else\ pre\ C_2\ Q\ l\ s)\ |$
$pre\ (\{I/x\}\ WHILE\ b\ DO\ C)\ Q = (\lambda l\ s.\ \exists\, n.\ I\ (l(x{:=}n))\ s)$

Verification condition:

fun $vc :: acom \Rightarrow assn2 \Rightarrow bool$ **where**
$vc\ SKIP\ Q = True\ |$
$vc\ (x ::= a)\ Q = True\ |$
$vc\ (C_1;;\ C_2)\ Q = (vc\ C_1\ (pre\ C_2\ Q) \wedge vc\ C_2\ Q)\ |$
$vc\ (IF\ b\ THEN\ C_1\ ELSE\ C_2)\ Q = (vc\ C_1\ Q \wedge vc\ C_2\ Q)\ |$
$vc\ (\{I/x\}\ WHILE\ b\ DO\ C)\ Q =$
$(\forall\, l\ s.\ (I\ (l(x{:=}Suc(l\ x)))\ s \longrightarrow pre\ C\ I\ l\ s) \wedge$
$(l\ x > 0 \wedge I\ l\ s \longrightarrow bval\ b\ s) \wedge$
$(I\ (l(x := 0))\ s \longrightarrow \neg\ bval\ b\ s \wedge Q\ l\ s) \wedge$
$vc\ C\ I)$

lemma *vc_sound*: $vc\ C\ Q \Longrightarrow \vdash_t \{pre\ C\ Q\}\ strip\ C\ \{Q\}$
proof(*induction C arbitrary*: *Q*)
case (*Awhile I x b C*)
show *?case*
proof(*simp*, *rule weaken_post*[*OF While*[*of I x*]], *goal_cases*)
case *1* **show** *?case*
using *Awhile.IH*[*of I*] *Awhile.prems* **by** (*auto intro*: *strengthen_pre*)
next
case *3* **show** *?case*
using *Awhile.prems* **by** (*simp*) (*metis fun_upd_triv*)
qed (*insert Awhile.prems*, *auto*)
qed (*auto intro*: *conseq Seq If simp*: *Skip Assign*)

Completeness:

lemma *pre_mono*:
$\forall\, l\ s.\ P\ l\ s \longrightarrow P'\ l\ s \Longrightarrow pre\ C\ P\ l\ s \Longrightarrow pre\ C\ P'\ l\ s$
proof (*induction C arbitrary*: $P\ P'\ l\ s$)
case *Aseq* **thus** *?case* **by** *simp metis*
qed *simp_all*

lemma *vc_mono*:
$\forall\, l\ s.\ P\ l\ s \longrightarrow P'\ l\ s \Longrightarrow vc\ C\ P \Longrightarrow vc\ C\ P'$
proof(*induction C arbitrary*: $P\ P'$)
case *Aseq* **thus** *?case* **by** *simp* (*metis pre_mono*)
qed *simp_all*

```
lemma vc_complete:
⊢t {P}c{Q} ⟹ ∃ C. strip C = c ∧ vc C Q ∧ (∀ l s. P l s ⟶ pre C Q l
s)
  (is _ ⟹ ∃ C. ?G P c Q C)
proof (induction rule: hoaret.induct)
  case Skip
  show ?case (is ∃ C. ?C C)
  proof show ?C Askip by simp qed
next
  case (Assign P a x)
  show ?case (is ∃ C. ?C C)
  proof show ?C(Aassign x a) by simp qed
next
  case (Seq P c1 Q c2 R)
  from Seq.IH obtain C1 where ih1: ?G P c1 Q C1 by blast
  from Seq.IH obtain C2 where ih2: ?G Q c2 R C2 by blast
  show ?case (is ∃ C. ?C C)
  proof
    show ?C(Aseq C1 C2)
      using ih1 ih2 by (fastforce elim!: pre_mono vc_mono)
  qed
next
  case (If P b c1 Q c2)
  from If.IH obtain C1 where ih1: ?G (λl s. P l s ∧ bval b s) c1 Q C1
    by blast
  from If.IH obtain C2 where ih2: ?G (λl s. P l s ∧ ¬bval b s) c2 Q C2
    by blast
  show ?case (is ∃ C. ?C C)
  proof
    show ?C(Aif b C1 C2) using ih1 ih2 by simp
  qed
next
  case (While P x c b)
  from While.IH obtain C where
    ih: ?G (λl s. P (l(x:=Suc(l x))) s ∧ bval b s) c P C
    by blast
  show ?case (is ∃ C. ?C C)
  proof
    have vc ({P/x} WHILE b DO C) (λl. P (l(x := 0)))
      using ih While.hyps(2,3)
      by simp (metis fun_upd_same zero_less_Suc)
    thus ?C(Awhile P x b C) using ih by simp
 qed
next
```

case *conseq* **thus** *?case* **by**(*fast elim*!: *pre_mono vc_mono*)
qed

Two examples:

lemma *vc1*: *vc*
({λ*l s*. *l* ''*x*'' = *nat*(*s* ''*x*'') / ''*x*''} *WHILE Less* (*N 0*) (*V* ''*x*'') *DO* ''*x*''
::= *Plus* (*V* ''*x*'') (*N* (−*1*)))
(λ*l s*. *s* ''*x*'' ≤ *0*)
by *auto*

thm *vc_sound*[*OF vc1*, *simplified*]

lemma *vc2*: *vc*
({λ*l s*. *l* ''*x*'' = *nat*(*s* ''*x*'') / ''*x*''} *WHILE Less* (*N 0*) (*V* ''*x*'')
DO (''*x*'' ::= *Plus* (*V* ''*x*'') (*N* (−*1*));;
(''*y*'' ::= *V* ''*x*'';;
{λ*l s*. *l* ''*x*'' = *nat*(*s* ''*x*'') ∧ *l* ''*y*'' = *nat*(*s* ''*y*'') / ''*y*''}
WHILE Less (*N 0*) (*V* ''*y*'') *DO* ''*y*'' ::= *Plus* (*V* ''*y*'') (*N* (−*1*)))))
(λ*l s*. *s* ''*x*'' ≤ *0*)
by *auto*

thm *vc_sound*[*OF vc2*, *simplified*]

end

14 Abstract Interpretation

14.1 Complete Lattice

theory *Complete_Lattice*
imports *Main*
begin

locale *Complete_Lattice* =
fixes *L* :: *'a*::*order set* **and** *Glb* :: *'a set* ⇒ *'a*
assumes *Glb_lower*: *A* ⊆ *L* ⟹ *a* ∈ *A* ⟹ *Glb A* ≤ *a*
and *Glb_greatest*: *b* ∈ *L* ⟹ ∀ *a*∈*A*. *b* ≤ *a* ⟹ *b* ≤ *Glb A*
and *Glb_in_L*: *A* ⊆ *L* ⟹ *Glb A* ∈ *L*
begin

definition *lfp* :: (*'a* ⇒ *'a*) ⇒ *'a* **where**
lfp f = *Glb* {*a* : *L*. *f a* ≤ *a*}

lemma *index_lfp*: *lfp f* ∈ *L*

by(*auto simp*: *lfp_def intro*: *Glb_in_L*)

lemma *lfp_lowerbound*:
  ⟦ $a \in L$; $f\ a \le a$ ⟧ $\Longrightarrow$ *lfp* $f \le a$
by (*auto simp add*: *lfp_def intro*: *Glb_lower*)

lemma *lfp_greatest*:
  ⟦ $a \in L$; $\bigwedge u.$ ⟦ $u \in L$; $f\ u \le u$⟧ $\Longrightarrow a \le u$ ⟧ $\Longrightarrow a \le$ *lfp* f
by (*auto simp add*: *lfp_def intro*: *Glb_greatest*)

lemma *lfp_unfold*: **assumes** $\bigwedge x.\ f\ x \in L \longleftrightarrow x \in L$
and *mono*: *mono f* **shows** *lfp f* = *f* (*lfp f*)
proof−
  note *assms*(*1*)[*simp*] *index_lfp*[*simp*]
  have *1*: *f* (*lfp f*) ≤ *lfp f*
    apply(*rule lfp_greatest*)
    apply *simp*
    by (*blast intro*: *lfp_lowerbound monoD*[*OF mono*] *order_trans*)
  have *lfp f* ≤ *f* (*lfp f*)
    by (*fastforce intro*: *1 monoD*[*OF mono*] *lfp_lowerbound*)
  with *1* **show** *?thesis* **by**(*blast intro*: *order_antisym*)
qed

end

end

14.2 Annotated Commands

theory *ACom*
imports *Com*
begin

datatype *'a acom* =
  SKIP 'a (*SKIP* {_} *61*) |
  Assign vname aexp 'a ((_ ::= _/ {_}) [*1000*, *61*, *0*] *61*) |
  Seq (*'a acom*) (*'a acom*) (_;;//_ [*60*, *61*] *60*) |
  If bexp 'a (*'a acom*) *'a* (*'a acom*) *'a*
    ((*IF* _/ *THEN* ({_}/ _)/ *ELSE* ({_}/ _)//{_}) [*0*, *0*, *0*, *61*, *0*, *0*]
61) |
  While 'a bexp 'a (*'a acom*) *'a*
    (({_}//*WHILE* _//*DO* ({_}//_)//{_}) [*0*, *0*, *0*, *61*, *0*] *61*)

notation *com.SKIP* (*SKIP*)

fun $strip :: {}'a\ acom \Rightarrow com$ **where**
$strip\ (SKIP\ \{P\}) = SKIP\ |$
$strip\ (x ::= e\ \{P\}) = x ::= e\ |$
$strip\ (C_1;;C_2) = strip\ C_1;;\ strip\ C_2\ |$
$strip\ (IF\ b\ THEN\ \{P_1\}\ C_1\ ELSE\ \{P_2\}\ C_2\ \{P\}) =$
$\quad IF\ b\ THEN\ strip\ C_1\ ELSE\ strip\ C_2\ |$
$strip\ (\{I\}\ WHILE\ b\ DO\ \{P\}\ C\ \{Q\}) = WHILE\ b\ DO\ strip\ C$

fun $asize :: com \Rightarrow nat$ **where**
$asize\ SKIP = 1\ |$
$asize\ (x ::= e) = 1\ |$
$asize\ (C_1;;C_2) = asize\ C_1 + asize\ C_2\ |$
$asize\ (IF\ b\ THEN\ C_1\ ELSE\ C_2) = asize\ C_1 + asize\ C_2 + 3\ |$
$asize\ (WHILE\ b\ DO\ C) = asize\ C + 3$

definition $shift :: (nat \Rightarrow {}'a) \Rightarrow nat \Rightarrow nat \Rightarrow {}'a$ **where**
$shift\ f\ n = (\lambda p.\ f(p+n))$

fun $annotate :: (nat \Rightarrow {}'a) \Rightarrow com \Rightarrow {}'a\ acom$ **where**
$annotate\ f\ SKIP = SKIP\ \{f\ 0\}\ |$
$annotate\ f\ (x ::= e) = x ::= e\ \{f\ 0\}\ |$
$annotate\ f\ (c_1;;c_2) = annotate\ f\ c_1;;\ annotate\ (shift\ f\ (asize\ c_1))\ c_2\ |$
$annotate\ f\ (IF\ b\ THEN\ c_1\ ELSE\ c_2) =$
$\quad IF\ b\ THEN\ \{f\ 0\}\ annotate\ (shift\ f\ 1)\ c_1$
$\quad ELSE\ \{f(asize\ c_1 + 1)\}\ annotate\ (shift\ f\ (asize\ c_1 + 2))\ c_2$
$\quad \{f(asize\ c_1 + asize\ c_2 + 2)\}\ |$
$annotate\ f\ (WHILE\ b\ DO\ c) =$
$\quad \{f\ 0\}\ WHILE\ b\ DO\ \{f\ 1\}\ annotate\ (shift\ f\ 2)\ c\ \{f(asize\ c + 2)\}$

fun $annos :: {}'a\ acom \Rightarrow {}'a\ list$ **where**
$annos\ (SKIP\ \{P\}) = [P]\ |$
$annos\ (x ::= e\ \{P\}) = [P]\ |$
$annos\ (C_1;;C_2) = annos\ C_1\ @\ annos\ C_2\ |$
$annos\ (IF\ b\ THEN\ \{P_1\}\ C_1\ ELSE\ \{P_2\}\ C_2\ \{Q\}) =$
$\quad P_1\ \#\ annos\ C_1\ @\ \ P_2\ \#\ annos\ C_2\ @\ [Q]\ |$
$annos\ (\{I\}\ WHILE\ b\ DO\ \{P\}\ C\ \{Q\}) = I\ \#\ P\ \#\ annos\ C\ @\ [Q]$
definition $anno :: {}'a\ acom \Rightarrow nat \Rightarrow {}'a$ **where**
$anno\ C\ p = annos\ C\ !\ p$

definition $post :: {}'a\ acom \Rightarrow {}'a$ **where**
$post\ C = last(annos\ C)$
fun $map\_acom :: ({}'a \Rightarrow {}'b) \Rightarrow {}'a\ acom \Rightarrow {}'b\ acom$ **where**
$map\_acom\ f\ (SKIP\ \{P\}) = SKIP\ \{f\ P\}\ |$
$map\_acom\ f\ (x ::= e\ \{P\}) = x ::= e\ \{f\ P\}\ |$

map_acom f $(C_1;;C_2)$ = *map_acom f* C_1;; *map_acom f* C_2 |
map_acom f (*IF b THEN* $\{P_1\}$ C_1 *ELSE* $\{P_2\}$ C_2 $\{Q\}$) =
  IF b THEN $\{f\ P_1\}$ *map_acom f* C_1 *ELSE* $\{f\ P_2\}$ *map_acom f* C_2
  $\{f\ Q\}$ |
map_acom f ($\{I\}$ *WHILE b DO* $\{P\}$ *C* $\{Q\}$) =
  $\{f\ I\}$ *WHILE b DO* $\{f\ P\}$ *map_acom f C* $\{f\ Q\}$

lemma *annos_ne*: *annos C* $\neq$ []
by(*induction C*) *auto*

lemma *strip_annotate*[*simp*]: *strip*(*annotate f c*) = *c*
by(*induction c arbitrary*: *f*) *auto*

lemma *length_annos_annotate*[*simp*]: *length* (*annos* (*annotate f c*)) = *asize*
c
by(*induction c arbitrary*: *f*) *auto*

lemma *size_annos*: *size*(*annos C*) = *asize*(*strip C*)
by(*induction C*)(*auto*)

lemma *size_annos_same*: *strip C1* = *strip C2* $\Longrightarrow$ *size*(*annos C1*) =
size(*annos C2*)
apply(*induct C2 arbitrary*: *C1*)
apply(*case_tac C1*, *simp_all*)+
done

lemmas *size_annos_same2* = *eqTrueI*[*OF size_annos_same*]

lemma *anno_annotate*[*simp*]: *p* < *asize c* $\Longrightarrow$ *anno* (*annotate f c*) *p* = *f p*
apply(*induction c arbitrary*: *f p*)
apply (*auto simp*: *anno_def nth_append nth_Cons numeral_eq_Suc shift_def*
  split: *nat.split*)
  apply (*metis add_Suc_right add_diff_inverse add.commute*)
 apply(*rule_tac f=f* **in** *arg_cong*)
 apply *arith*
apply (*metis less_Suc_eq*)
done

lemma *eq_acom_iff_strip_annos*:
  C1 = *C2* $\longleftrightarrow$ *strip C1* = *strip C2* $\wedge$ *annos C1* = *annos C2*
apply(*induction C1 arbitrary*: *C2*)
apply(*case_tac C2*, *auto simp*: *size_annos_same2*)+
done

lemma *eq_acom_iff_strip_anno*:
  $C1=C2 \longleftrightarrow strip\ C1 = strip\ C2 \wedge (\forall p<size(annos\ C1).\ anno\ C1\ p = anno\ C2\ p)$
by(*auto simp add*: *eq_acom_iff_strip_annos anno_def*
  list_eq_iff_nth_eq size_annos_same2)

lemma *post_map_acom*[*simp*]: $post(map\_acom\ f\ C) = f(post\ C)$
by (*induction C*) (*auto simp*: *post_def last_append annos_ne*)

lemma *strip_map_acom*[*simp*]: $strip\ (map\_acom\ f\ C) = strip\ C$
by (*induction C*) *auto*

lemma *anno_map_acom*: $p < size(annos\ C) \Longrightarrow anno\ (map\_acom\ f\ C)\ p = f(anno\ C\ p)$
apply(*induction C arbitrary*: *p*)
apply(*auto simp*: *anno_def nth_append nth_Cons' size_annos*)
done

lemma *strip_eq_SKIP*:
  $strip\ C = SKIP \longleftrightarrow (\exists P.\ C = SKIP\ \{P\})$
by (*cases C*) *simp_all*

lemma *strip_eq_Assign*:
  $strip\ C = x{::=}e \longleftrightarrow (\exists P.\ C = x{::=}e\ \{P\})$
by (*cases C*) *simp_all*

lemma *strip_eq_Seq*:
  $strip\ C = c1;;c2 \longleftrightarrow (\exists C1\ C2.\ C = C1;;C2\ \&\ strip\ C1 = c1\ \&\ strip\ C2 = c2)$
by (*cases C*) *simp_all*

lemma *strip_eq_If*:
  $strip\ C = IF\ b\ THEN\ c1\ ELSE\ c2 \longleftrightarrow$
  $(\exists P1\ P2\ C1\ C2\ Q.\ C = IF\ b\ THEN\ \{P1\}\ C1\ ELSE\ \{P2\}\ C2\ \{Q\}\ \&\ strip\ C1 = c1\ \&\ strip\ C2 = c2)$
by (*cases C*) *simp_all*

lemma *strip_eq_While*:
  $strip\ C = WHILE\ b\ DO\ c1 \longleftrightarrow$
  $(\exists I\ P\ C1\ Q.\ C = \{I\}\ WHILE\ b\ DO\ \{P\}\ C1\ \{Q\}\ \&\ strip\ C1 = c1)$
by (*cases C*) *simp_all*

lemma [*simp*]: $shift\ (\lambda p.\ a)\ n = (\lambda p.\ a)$
by(*simp add*:*shift_def*)

```
lemma set_annos_anno[simp]: set (annos (annotate (λp. a) c)) = {a}
by(induction c) simp_all

lemma post_in_annos: post C ∈ set(annos C)
by(auto simp: post_def annos_ne)

lemma post_anno_asize: post C = anno C (size(annos C) − 1)
by(simp add: post_def last_conv_nth[OF annos_ne] anno_def)

end
```

14.3 Collecting Semantics of Commands

```
theory Collecting
imports Complete_Lattice Big_Step ACom
begin
```

14.3.1 The generic Step function

```
notation
  sup (infixl ⊔ 65) and
  inf (infixl ⊓ 70) and
  bot (⊥) and
  top (⊤)

context
  fixes f :: vname ⇒ aexp ⇒ 'a ⇒ 'a::sup
  fixes g :: bexp ⇒ 'a ⇒ 'a
begin
fun Step :: 'a ⇒ 'a acom ⇒ 'a acom where
Step S (SKIP {Q}) = (SKIP {S}) |
Step S (x ::= e {Q}) =
  x ::= e {f x e S} |
Step S (C1;; C2) = Step S C1;; Step (post C1) C2 |
Step S (IF b THEN {P1} C1 ELSE {P2} C2 {Q}) =
  IF b THEN {g b S} Step P1 C1 ELSE {g (Not b) S} Step P2 C2
  {post C1 ⊔ post C2} |
Step S ({I} WHILE b DO {P} C {Q}) =
  {S ⊔ post C} WHILE b DO {g b I} Step P C {g (Not b) I}
end

lemma strip_Step[simp]: strip(Step f g S C) = strip C
by(induct C arbitrary: S) auto
```

14.3.2 Annotated commands as a complete lattice

instantiation *acom* :: (*order*) *order*
begin

definition *less_eq_acom* :: ('*a*::*order*)*acom* ⇒ '*a acom* ⇒ *bool* **where**
C1 ≤ *C2* ⟷ *strip C1* = *strip C2* ∧ (∀ *p*<*size*(*annos C1*). *anno C1 p* ≤ *anno C2 p*)

definition *less_acom* :: '*a acom* ⇒ '*a acom* ⇒ *bool* **where**
less_acom x y = (*x* ≤ *y* ∧ ¬ *y* ≤ *x*)

instance
proof (*standard*, *goal_cases*)
  case *1* **show** *?case* **by**(*simp add*: *less_acom_def*)
next
  case *2* **thus** *?case* **by**(*auto simp*: *less_eq_acom_def*)
next
  case *3* **thus** *?case* **by**(*fastforce simp*: *less_eq_acom_def size_annos*)
next
  case *4* **thus** *?case*
    by(*fastforce simp*: *le_antisym less_eq_acom_def size_annos eq_acom_iff_strip_anno*)
qed

end

lemma *less_eq_acom_annos*:
  C1 ≤ *C2* ⟷ *strip C1* = *strip C2* ∧ *list_all2* (≤) (*annos C1*) (*annos C2*)
by(*auto simp add*: *less_eq_acom_def anno_def list_all2_conv_all_nth size_annos_same2*)

lemma *SKIP_le*[*simp*]: *SKIP* {*S*} ≤ *c* ⟷ (∃ *S*'. *c* = *SKIP* {*S*'} ∧ *S* ≤ *S*')
by (*cases c*) (*auto simp*:*less_eq_acom_def anno_def*)

lemma *Assign_le*[*simp*]: *x* ::= *e* {*S*} ≤ *c* ⟷ (∃ *S*'. *c* = *x* ::= *e* {*S*'} ∧ *S* ≤ *S*')
by (*cases c*) (*auto simp*:*less_eq_acom_def anno_def*)

lemma *Seq_le*[*simp*]: *C1*;;*C2* ≤ *C* ⟷ (∃ *C1*' *C2*'. *C* = *C1*';;*C2*' ∧ *C1* ≤ *C1*' ∧ *C2* ≤ *C2*')
apply (*cases C*)
apply(*auto simp*: *less_eq_acom_annos list_all2_append size_annos_same2*)

done

lemma *If_le*[*simp*]: *IF b THEN* {*p1*} *C1 ELSE* {*p2*} *C2* {*S*} ≤ *C* ⟷
  (∃ *p1' p2' C1' C2' S'*. *C* = *IF b THEN* {*p1'*} *C1' ELSE* {*p2'*} *C2'* {*S'*}
∧
    p1 ≤ *p1'* ∧ *p2* ≤ *p2'* ∧ *C1* ≤ *C1'* ∧ *C2* ≤ *C2'* ∧ *S* ≤ *S'*)
apply (*cases C*)
apply(*auto simp*: *less_eq_acom_annos list_all2_append size_annos_same2*)
done

lemma *While_le*[*simp*]: {*I*} *WHILE b DO* {*p*} *C* {*P*} ≤ *W* ⟷
  (∃ *I' p' C' P'*. *W* = {*I'*} *WHILE b DO* {*p'*} *C'* {*P'*} ∧ *C* ≤ *C'* ∧ *p* ≤ *p'*
∧ *I* ≤ *I'* ∧ *P* ≤ *P'*)
apply (*cases W*)
apply(*auto simp*: *less_eq_acom_annos list_all2_append size_annos_same2*)
done

lemma *mono_post*: *C* ≤ *C'* ⟹ *post C* ≤ *post C'*
using *annos_ne*[*of C'*]
by(*auto simp*: *post_def less_eq_acom_def last_conv_nth*[*OF annos_ne*]
anno_def
    dest: *size_annos_same*)

definition *Inf_acom* :: *com* ⇒ *'a*::*complete_lattice acom set* ⇒ *'a acom*
where
Inf_acom c M = *annotate* (λ*p*. *INF C*∈*M*. *anno C p*) *c*

global_interpretation
  Complete_Lattice {*C*. *strip C* = *c*} *Inf_acom c* **for** *c*
proof (*standard*, *goal_cases*)
  case *1* **thus** *?case*
   by(*auto simp*: *Inf_acom_def less_eq_acom_def size_annos intro*:*INF_lower*)
next
  case *2* **thus** *?case*
   by(*auto simp*: *Inf_acom_def less_eq_acom_def size_annos intro*:*INF_greatest*)
next
  case *3* **thus** *?case* **by**(*auto simp*: *Inf_acom_def*)
qed

14.3.3 Collecting semantics

definition *step* = *Step* (λ*x e S*. {*s*(*x* := *aval e s*) |*s*. *s* ∈ *S*}) (λ*b S*. {*s*:*S*.
bval b s})

```
definition CS :: com ⇒ state set acom where
CS c = lfp c (step UNIV)

lemma mono2_Step: fixes C1 C2 :: 'a::semilattice_sup acom
  assumes !!x e S1 S2. S1 ≤ S2 ⟹ f x e S1 ≤ f x e S2
          !!b S1 S2. S1 ≤ S2 ⟹ g b S1 ≤ g b S2
  shows C1 ≤ C2 ⟹ S1 ≤ S2 ⟹ Step f g S1 C1 ≤ Step f g S2 C2
proof(induction S1 C1 arbitrary: C2 S2 rule: Step.induct)
  case 1 thus ?case by(auto)
next
  case 2 thus ?case by (auto simp: assms(1))
next
  case 3 thus ?case by(auto simp: mono_post)
next
  case 4 thus ?case
    by(auto simp: subset_iff assms(2))
      (metis mono_post le_supI1 le_supI2)+
next
  case 5 thus ?case
    by(auto simp: subset_iff assms(2))
      (metis mono_post le_supI1 le_supI2)+
qed

lemma mono2_step: C1 ≤ C2 ⟹ S1 ⊆ S2 ⟹ step S1 C1 ≤ step S2 C2
unfolding step_def by(rule mono2_Step) auto

lemma mono_step: mono (step S)
by(blast intro: monoI mono2_step)

lemma strip_step: strip(step S C) = strip C
by (induction C arbitrary: S) (auto simp: step_def)

lemma lfp_cs_unfold: lfp c (step S) = step S (lfp c (step S))
apply(rule lfp_unfold[OF _  mono_step])
apply(simp add: strip_step)
done

lemma CS_unfold: CS c = step UNIV (CS c)
by (metis CS_def lfp_cs_unfold)

lemma strip_CS[simp]: strip(CS c) = c
by(simp add: CS_def index_lfp[simplified])
```

14.3.4 Relation to big-step semantics

lemma *asize_nz*: *asize*(*c*::*com*) ≠ *0*
by (*metis length_0_conv length_annos_annotate annos_ne*)

lemma *post_Inf_acom*:
∀ *C*∈*M*. *strip C* = *c* ⟹ *post* (*Inf_acom c M*) = ⋂(*post* ' *M*)
apply(*subgoal_tac* ∀ *C*∈*M*. *size*(*annos C*) = *asize c*)
apply(*simp add*: *post_anno_asize Inf_acom_def asize_nz neq0_conv*[*symmetric*])
apply(*simp add*: *size_annos*)
done

lemma *post_lfp*: *post*(*lfp c f*) = (⋂{*post C*|*C*. *strip C* = *c* ∧ *f C* ≤ *C*})
by(*auto simp add*: *lfp_def post_Inf_acom*)

lemma *big_step_post_step*:
⟦ (*c*, *s*) ⇒ *t*; *strip C* = *c*; *s* ∈ *S*; *step S C* ≤ *C* ⟧ ⟹ *t* ∈ *post C*
proof(*induction arbitrary*: *C S rule*: *big_step_induct*)
case *Skip* **thus** *?case* **by**(*auto simp*: *strip_eq_SKIP step_def post_def*)
next
case *Assign* **thus** *?case*
by(*fastforce simp*: *strip_eq_Assign step_def post_def*)
next
case *Seq* **thus** *?case*
by(*fastforce simp*: *strip_eq_Seq step_def post_def last_append annos_ne*)
next
case *IfTrue* **thus** *?case* **apply**(*auto simp*: *strip_eq_If step_def post_def*)
by (*metis* (*lifting*,*full_types*) *mem_Collect_eq subsetD*)
next
case *IfFalse* **thus** *?case* **apply**(*auto simp*: *strip_eq_If step_def post_def*)
by (*metis* (*lifting*,*full_types*) *mem_Collect_eq subsetD*)
next
case (*WhileTrue b s1 c' s2 s3*)
from *WhileTrue.prems*(*1*) **obtain** *I P C' Q* **where** *C* = {*I*} *WHILE b DO* {*P*} *C'* {*Q*} *strip C'* = *c'*
by(*auto simp*: *strip_eq_While*)
from *WhileTrue.prems*(*3*) ‹*C* = _›
have *step P C'* ≤ *C'* {*s* ∈ *I*. *bval b s*} ≤ *P* *S* ≤ *I* *step* (*post C'*) *C* ≤ *C*
by (*auto simp*: *step_def post_def*)
have *step* {*s* ∈ *I*. *bval b s*} *C'* ≤ *C'*
by (*rule order_trans*[*OF mono2_step*[*OF order_refl* ‹{*s* ∈ *I*. *bval b s*} ≤ *P*›] ‹*step P C'* ≤ *C'*›])
have *s1* ∈ {*s*∈*I*. *bval b s*} **using** ‹*s1* ∈ *S*› ‹*S* ⊆ *I*› ‹*bval b s1*› **by** *auto*

note *s2_in_post_C′* = *WhileTrue.IH*(*1*)[*OF* ‹*strip* *C′* = *c′*› *this* ‹*step* {*s* ∈ *I*. *bval* *b* *s*} *C′* ≤ *C′*›]
from *WhileTrue.IH*(*2*)[*OF* *WhileTrue.prems*(*1*) *s2_in_post_C′* ‹*step* (*post* *C′*) *C* ≤ *C*›]
show *?case* **.**
next
case (*WhileFalse* *b* *s1* *c′*) **thus** *?case*
by (*force* *simp*: *strip_eq_While* *step_def* *post_def*)
qed

lemma *big_step_lfp*: ⟦ (*c*,*s*) ⇒ *t*; *s* ∈ *S* ⟧ ⟹ *t* ∈ *post*(*lfp* *c* (*step* *S*))
by(*auto* *simp* *add*: *post_lfp* *intro*: *big_step_post_step*)

lemma *big_step_CS*: (*c*,*s*) ⇒ *t* ⟹ *t* ∈ *post*(*CS* *c*)
by(*simp* *add*: *CS_def* *big_step_lfp*)

end

14.4 A small step semantics on annotated commands

theory *Collecting1*
imports *Collecting*
begin

The idea: the state is propagated through the annotated command as an annotation {*s*}, all other annotations are {}. It is easy to show that this semantics approximates the collecting semantics.

lemma *step_preserves_le*:
⟦ *step* *S* *cs* = *cs*; *S′* ⊆ *S*; *cs′* ≤ *cs* ⟧ ⟹
step *S′* *cs′* ≤ *cs*
by (*metis* *mono2_step*)

lemma *steps_empty_preserves_le*: **assumes** *step* *S* *cs* = *cs*
shows *cs′* ≤ *cs* ⟹ (*step* {} ^^ *n*) *cs′* ≤ *cs*
proof(*induction* *n* *arbitrary*: *cs′*)
case *0* **thus** *?case* **by** *simp*
next
case (*Suc* *n*) **thus** *?case*
using *Suc.IH*[*OF* *step_preserves_le*[*OF* *assms* *empty_subsetI* *Suc.prems*]]
by(*simp* *add*:*funpow_swap1*)
qed

definition *steps* :: *state* ⇒ *com* ⇒ *nat* ⇒ *state* *set* *acom* **where**

steps s c n = ((*step* {})^^*n*) (*step* {*s*} (*annotate* (λ*p*. {}) *c*))

lemma *steps_approx_fix_step*: **assumes** *step S cs* = *cs* **and** *s* ∈ *S*
shows *steps s* (*strip cs*) *n* ≤ *cs*
proof−
  let *?bot* = *annotate* (λ*p*. {}) (*strip cs*)
  have *?bot* ≤ *cs* **by**(*induction cs*) *auto*
  from *step_preserves_le*[*OF assms*(*1*)_ *this*, *of* {*s*}] ‹*s* ∈ *S*›
  have *1*: *step* {*s*} *?bot* ≤ *cs* **by** *simp*
  from *steps_empty_preserves_le*[*OF assms*(*1*) *1*]
  show *?thesis* **by**(*simp add*: *steps_def*)
qed

theorem *steps_approx_CS*: *steps s c n* ≤ *CS c*
by (*metis CS_unfold UNIV_I steps_approx_fix_step strip_CS*)

end

14.5 Collecting Semantics Examples

theory *Collecting_Examples*
imports *Collecting Vars*
begin

14.5.1 Pretty printing state sets

Tweak code generation to work with sets of non-equality types:

declare *insert_code*[*code del*] *union_coset_filter*[*code del*]
lemma *insert_code* [*code*]: *insert x* (*set xs*) = *set* (*x*#*xs*)
by *simp*

Compensate for the fact that sets may now have duplicates:

definition *compact* :: '*a set* ⇒ '*a set* **where**
compact X = *X*

lemma [*code*]: *compact*(*set xs*) = *set*(*remdups xs*)
by(*simp add*: *compact_def*)

definition *vars_acom* = *compact o vars o strip*

In order to display commands annotated with state sets, states must be translated into a printable format as sets of variable-state pairs, for the variables in the command:

definition *show_acom* :: *state set acom* ⇒ (*vname*∗*val*)*set set acom* **where**

show_acom C =
  annotate (λ*p*. (λ*s*. (λ*x*. (*x*, *s x*)) ‘ (*vars_acom C*)) ‘ *anno C p*) (*strip C*)

14.5.2 Examples

definition *c0* = *WHILE Less* (*V ''x''*) (*N 3*)
    DO ''x'' ::= *Plus* (*V ''x''*) (*N 2*)

definition *C0* :: *state set acom* **where** *C0* = *annotate* (λ*p*. {}) *c0*

Collecting semantics:

value *show_acom* (((*step* {<>}) ^^ *0*) *C0*)
value *show_acom* (((*step* {<>}) ^^ *1*) *C0*)
value *show_acom* (((*step* {<>}) ^^ *2*) *C0*)
value *show_acom* (((*step* {<>}) ^^ *3*) *C0*)
value *show_acom* (((*step* {<>}) ^^ *4*) *C0*)
value *show_acom* (((*step* {<>}) ^^ *5*) *C0*)
value *show_acom* (((*step* {<>}) ^^ *6*) *C0*)
value *show_acom* (((*step* {<>}) ^^ *7*) *C0*)
value *show_acom* (((*step* {<>}) ^^ *8*) *C0*)

Small-step semantics:

value *show_acom* (((*step* {}) ^^ *0*) (*step* {<>} *C0*))
value *show_acom* (((*step* {}) ^^ *1*) (*step* {<>} *C0*))
value *show_acom* (((*step* {}) ^^ *2*) (*step* {<>} *C0*))
value *show_acom* (((*step* {}) ^^ *3*) (*step* {<>} *C0*))
value *show_acom* (((*step* {}) ^^ *4*) (*step* {<>} *C0*))
value *show_acom* (((*step* {}) ^^ *5*) (*step* {<>} *C0*))
value *show_acom* (((*step* {}) ^^ *6*) (*step* {<>} *C0*))
value *show_acom* (((*step* {}) ^^ *7*) (*step* {<>} *C0*))
value *show_acom* (((*step* {}) ^^ *8*) (*step* {<>} *C0*))

end

14.6 Abstract Interpretation Test Programs

theory *Abs_Int_Tests*
imports *Com*
begin

For constant propagation:

Straight line code:

definition *test1_const* =
 ''y'' ::= *N 7*;;

''z'' ::= Plus (V ''y'') (N 2);;
''y'' ::= Plus (V ''x'') (N 0)

Conditional:

definition *test2_const* =
IF Less (N 41) (V ''x'') THEN ''x'' ::= N 5 ELSE ''x'' ::= N 5

Conditional, test is relevant:

definition *test3_const* =
''x'' ::= N 42;;
IF Less (N 41) (V ''x'') THEN ''x'' ::= N 5 ELSE ''x'' ::= N 6

While:

definition *test4_const* =
''x'' ::= N 0;; WHILE Bc True DO ''x'' ::= N 0

While, test is relevant:

definition *test5_const* =
''x'' ::= N 0;; WHILE Less (V ''x'') (N 1) DO ''x'' ::= N 1

Iteration is needed:

definition *test6_const* =
''x'' ::= N 0;; ''y'' ::= N 0;; ''z'' ::= N 2;;
WHILE Less (V ''x'') (N 1) DO (''x'' ::= V ''y'';; ''y'' ::= V ''z'')

For intervals:

definition *test1_ivl* =
''y'' ::= N 7;;
IF Less (V ''x'') (V ''y'')
THEN ''y'' ::= Plus (V ''y'') (V ''x'')
ELSE ''x'' ::= Plus (V ''x'') (V ''y'')

definition *test2_ivl* =
WHILE Less (V ''x'') (N 100)
DO ''x'' ::= Plus (V ''x'') (N 1)

definition *test3_ivl* =
''x'' ::= N 0;;
WHILE Less (V ''x'') (N 100)
DO ''x'' ::= Plus (V ''x'') (N 1)

definition *test4_ivl* =
''x'' ::= N 0;; ''y'' ::= N 0;;
WHILE Less (V ''x'') (N 11)
DO (''x'' ::= Plus (V ''x'') (N 1);; ''y'' ::= Plus (V ''y'') (N 1))

definition *test5_ivl* =
''x'' ::= N 0;; ''y'' ::= N 0;;
WHILE Less (V ''x'') (N 100)
DO (''y'' ::= V ''x'';; ''x'' ::= Plus (V ''x'') (N 1))

definition *test6_ivl* =
''x'' ::= N 0;;
WHILE Less (N (− 1)) (V ''x'') DO ''x'' ::= Plus (V ''x'') (N 1)

end
theory *Abs_Int_init*
imports *HOL−Library.While_Combinator*
HOL−Library.Extended
Vars Collecting Abs_Int_Tests
begin

hide_const (**open**) *top bot dom* — to avoid qualified names

end

14.7 Abstract Interpretation

theory *Abs_Int0*
imports *Abs_Int_init*
begin

14.7.1 Orderings

The basic type classes *order*, *semilattice_sup* and *order_top* are defined in *Main*, more precisely in theories *HOL.Orderings* and *HOL.Lattices*. If you view this theory with jedit, just click on the names to get there.

class *semilattice_sup_top* = *semilattice_sup* + *order_top*

instance *fun* :: (*type*, *semilattice_sup_top*) *semilattice_sup_top* **..**

instantiation *option* :: (*order*)*order*
begin

fun *less_eq_option* **where**
Some x ≤ Some y = (x ≤ y) |
None ≤ y = True |
Some _ ≤ None = False

definition *less_option* **where** $x < (y{::}'a\ option) = (x \leq y \wedge \neg\ y \leq x)$

lemma *le_None*[*simp*]: $(x \leq None) = (x = None)$
by (*cases x*) *simp_all*

lemma *Some_le*[*simp*]: $(Some\ x \leq u) = (\exists\, y.\ u = Some\ y \wedge x \leq y)$
by (*cases u*) *auto*

instance
proof (*standard*, *goal_cases*)
  case *1* **show** *?case* **by**(*rule less_option_def*)
next
  case (*2 x*) **show** *?case* **by**(*cases x*, *simp_all*)
next
  case (*3 x y z*) **thus** *?case* **by**(*cases z*, *simp*, *cases y*, *simp*, *cases x*, *auto*)
next
  case (*4 x y*) **thus** *?case* **by**(*cases y*, *simp*, *cases x*, *auto*)
qed

end

instantiation *option* :: (*sup*)*sup*
begin

fun *sup_option* **where**
$Some\ x \sqcup Some\ y = Some(x \sqcup y)$ |
$None \sqcup y = y$ |
$x \sqcup None = x$

lemma *sup_None2*[*simp*]: $x \sqcup None = x$
by (*cases x*) *simp_all*

instance ..

end

instantiation *option* :: (*semilattice_sup_top*)*semilattice_sup_top*
begin

definition *top_option* **where** $\top = Some\ \top$

instance
proof (*standard*, *goal_cases*)

case (*4 a*) **show** *?case* **by**(*cases a*, *simp_all add*: *top_option_def*)
next
case (*1 x y*) **thus** *?case* **by**(*cases x*, *simp*, *cases y*, *simp_all*)
next
case (*2 x y*) **thus** *?case* **by**(*cases y*, *simp*, *cases x*, *simp_all*)
next
case (*3 x y z*) **thus** *?case* **by**(*cases z*, *simp*, *cases y*, *simp*, *cases x*, *simp_all*)
qed

end

lemma [*simp*]: (*Some x* < *Some y*) = (*x* < *y*)
by(*auto simp*: *less_le*)

instantiation *option* :: (*order*)*order_bot*
begin

definition *bot_option* :: *'a option* **where**
⊥ = *None*

instance
proof (*standard*, *goal_cases*)
case *1* **thus** *?case* **by**(*auto simp*: *bot_option_def*)
qed

end

definition *bot* :: *com* ⇒ *'a option acom* **where**
bot c = *annotate* (λ*p*. *None*) *c*

lemma *bot_least*: *strip C* = *c* ⟹ *bot c* ≤ *C*
by(*auto simp*: *bot_def less_eq_acom_def*)

lemma *strip_bot*[*simp*]: *strip*(*bot c*) = *c*
by(*simp add*: *bot_def*)

14.7.2 Pre-fixpoint iteration

definition *pfp* :: ((*'a*::*order*) ⇒ *'a*) ⇒ *'a* ⇒ *'a option* **where**
pfp f = *while_option* (λ*x*. ¬ *f x* ≤ *x*) *f*

lemma *pfp_pfp*: **assumes** *pfp f x0* = *Some x* **shows** *f x* ≤ *x*

using *while_option_stop*[*OF assms*[*simplified pfp_def*]] **by** *simp*

lemma *while_least*:
fixes q :: *'a::order*
assumes $\forall x{\in}L.\forall y{\in}L.\ x \le y \longrightarrow f\,x \le f\,y$ **and** $\forall x.\ x \in L \longrightarrow f\,x \in L$
and $\forall x \in L.\ b \le x$ **and** $b \in L$ **and** $f\,q \le q$ **and** $q \in L$
and *while_option* $P\ f\ b = Some\ p$
shows $p \le q$
using *while_option_rule*[*OF* _ *assms*(*7*)[*unfolded pfp_def*],
where $P = \%x.\ x \in L \wedge x \le q$]
by (*metis assms*(*1−6*) *order_trans*)

lemma *pfp_bot_least*:
assumes $\forall x{\in}\{C.\ strip\ C = c\}.\forall y{\in}\{C.\ strip\ C = c\}.\ x \le y \longrightarrow f\,x \le f\,y$
and $\forall C.\ C \in \{C.\ strip\ C = c\} \longrightarrow f\,C \in \{C.\ strip\ C = c\}$
and $f\,C' \le C'$ $strip\ C' = c$ *pfp* f (*bot* c) = *Some* C
shows $C \le C'$
by(*rule while_least*[*OF assms*(*1*,*2*) _ _ *assms*(*3*) _ *assms*(*5*)[*unfolded pfp_def*]])
(*simp_all add*: *assms*(*4*) *bot_least*)

lemma *pfp_inv*:
pfp $f\ x = Some\ y \Longrightarrow (\bigwedge x.\ P\ x \Longrightarrow P(f\ x)) \Longrightarrow P\ x \Longrightarrow P\ y$
unfolding *pfp_def* **by** (*blast intro*: *while_option_rule*)

lemma *strip_pfp*:
assumes $\bigwedge x.\ g(f\ x) = g\ x$ **and** *pfp* $f\ x0 = Some\ x$ **shows** $g\ x = g\ x0$
using *pfp_inv*[*OF assms*(*2*), **where** $P = \%x.\ g\ x = g\ x0$] *assms*(*1*) **by** *simp*

14.7.3 Abstract Interpretation

definition γ*_fun* :: $('a \Rightarrow 'b\ set) \Rightarrow ('c \Rightarrow 'a) \Rightarrow ('c \Rightarrow 'b)set$ **where**
γ*_fun* $\gamma\ F = \{f.\ \forall x.\ f\ x \in \gamma(F\ x)\}$

fun γ*_option* :: $('a \Rightarrow 'b\ set) \Rightarrow 'a\ option \Rightarrow 'b\ set$ **where**
γ*_option* γ *None* = {} |
γ*_option* γ (*Some* a) = $\gamma\ a$

The interface for abstract values:

locale *Val_semilattice* =
fixes γ :: *'av::semilattice_sup_top* $\Rightarrow$ *val set*
assumes *mono_gamma*: $a \le b \Longrightarrow \gamma\ a \le \gamma\ b$
and *gamma_Top*[*simp*]: $\gamma\ \top = UNIV$
fixes *num'* :: *val* $\Rightarrow$ *'av*

and *plus'* :: *'av* ⇒ *'av* ⇒ *'av*
  assumes *gamma_num'*: $i \in \gamma(num'\ i)$
  and *gamma_plus'*: $i1 \in \gamma\ a1 \Longrightarrow i2 \in \gamma\ a2 \Longrightarrow i1{+}i2 \in \gamma(plus'\ a1\ a2)$

type_synonym *'av st* = (*vname* ⇒ *'av*)

The for-clause (here and elsewhere) only serves the purpose of fixing the name of the type parameter *'av* which would otherwise be renamed to *'a*.

locale *Abs_Int_fun* = *Val_semilattice* **where** $\gamma{=}\gamma$
  for γ :: *'av::semilattice_sup_top* ⇒ *val set*
begin

fun *aval'* :: *aexp* ⇒ *'av st* ⇒ *'av* **where**
aval' (*N i*) *S* = *num' i* |
aval' (*V x*) *S* = *S x* |
aval' (*Plus a1 a2*) *S* = *plus'* (*aval' a1 S*) (*aval' a2 S*)

definition *asem x e S* = (*case S of None* ⇒ *None* | *Some S* ⇒ *Some*(*S*(*x* := *aval' e S*)))

definition *step'* = *Step asem* (λb *S*. *S*)

lemma *strip_step'*[*simp*]: *strip*(*step' S C*) = *strip C*
by(*simp add*: *step'_def*)

definition *AI* :: *com* ⇒ *'av st option acom option* **where**
AI c = *pfp* (*step'* ⊤) (*bot c*)

abbreviation γ_s :: *'av st* ⇒ *state set*
where γ_s == γ*_fun* γ

abbreviation γ_o :: *'av st option* ⇒ *state set*
where γ_o == γ*_option* γ_s

abbreviation γ_c :: *'av st option acom* ⇒ *state set acom*
where γ_c == *map_acom* γ_o

lemma *gamma_s_Top*[*simp*]: γ_s ⊤ = *UNIV*
by(*simp add*: *top_fun_def* γ*_fun_def*)

lemma *gamma_o_Top*[*simp*]: γ_o ⊤ = *UNIV*
by (*simp add*: *top_option_def*)

lemma *mono_gamma_s*: $f1 \le f2 \Longrightarrow \gamma_s\ f1 \subseteq \gamma_s\ f2$
by(*auto simp*: *le_fun_def γ_fun_def dest*: *mono_gamma*)

lemma *mono_gamma_o*:
  $S1 \le S2 \Longrightarrow \gamma_o\ S1 \subseteq \gamma_o\ S2$
by(*induction S1 S2 rule*: *less_eq_option.induct*)(*simp_all add*: *mono_gamma_s*)

lemma *mono_gamma_c*: $C1 \le C2 \Longrightarrow \gamma_c\ C1 \le \gamma_c\ C2$
by (*simp add*: *less_eq_acom_def mono_gamma_o size_annos anno_map_acom size_annos_same*[*of C1 C2*])

Correctness:

lemma *aval′_correct*: $s \in \gamma_s\ S \Longrightarrow aval\ a\ s \in \gamma(aval'\ a\ S)$
by (*induct a*) (*auto simp*: *gamma_num′ gamma_plus′ γ_fun_def*)

lemma *in_gamma_update*: $[\![\ s \in \gamma_s\ S;\ i \in \gamma\ a\ ]\!] \Longrightarrow s(x := i) \in \gamma_s(S(x := a))$
by(*simp add*: *γ_fun_def*)

lemma *gamma_Step_subcomm*:
  assumes $\bigwedge x\ e\ S.\ f1\ x\ e\ (\gamma_o\ S) \subseteq \gamma_o\ (f2\ x\ e\ S)$ $\bigwedge b\ S.\ g1\ b\ (\gamma_o\ S) \subseteq \gamma_o\ (g2\ b\ S)$
  shows $Step\ f1\ g1\ (\gamma_o\ S)\ (\gamma_c\ C) \le \gamma_c\ (Step\ f2\ g2\ S\ C)$
by (*induction C arbitrary*: *S*) (*auto simp*: *mono_gamma_o assms*)

lemma *step_step′*: $step\ (\gamma_o\ S)\ (\gamma_c\ C) \le \gamma_c\ (step'\ S\ C)$
unfolding *step_def step′_def*
by(*rule gamma_Step_subcomm*)
 (*auto simp*: *aval′_correct in_gamma_update asem_def split*: *option.splits*)

lemma *AI_correct*: $AI\ c = Some\ C \Longrightarrow CS\ c \le \gamma_c\ C$
proof(*simp add*: *CS_def AI_def*)
  assume *1*: $pfp\ (step'\ \top)\ (bot\ c) = Some\ C$
  have *pfp′*: $step'\ \top\ C \le C$ **by**(*rule pfp_pfp*[*OF 1*])
  have *2*: $step\ (\gamma_o\ \top)\ (\gamma_c\ C) \le \gamma_c\ C$ — transfer the pfp'
  proof(*rule order_trans*)
    show $step\ (\gamma_o\ \top)\ (\gamma_c\ C) \le \gamma_c\ (step'\ \top\ C)$ **by**(*rule step_step′*)
    show ... $\le \gamma_c\ C$ **by** (*metis mono_gamma_c*[*OF pfp′*])
  qed
  have *3*: $strip\ (\gamma_c\ C) = c$ **by**(*simp add*: *strip_pfp*[*OF _ 1*] *step′_def*)
  have $lfp\ c\ (step\ (\gamma_o\ \top)) \le \gamma_c\ C$
    by(*rule lfp_lowerbound*[*simplified*,**where** *f=step* $(\gamma_o\ \top)$, *OF 3 2*])
  thus $lfp\ c\ (step\ UNIV) \le \gamma_c\ C$ **by** *simp*
qed

end

14.7.4 Monotonicity

locale *Abs_Int_fun_mono* = *Abs_Int_fun* +
assumes *mono_plus'*: *a1* ≤ *b1* ⟹ *a2* ≤ *b2* ⟹ *plus' a1 a2* ≤ *plus' b1 b2*
begin

lemma *mono_aval'*: *S* ≤ *S'* ⟹ *aval' e S* ≤ *aval' e S'*
by(*induction e*)(*auto simp*: *le_fun_def mono_plus'*)

lemma *mono_update*: *a* ≤ *a'* ⟹ *S* ≤ *S'* ⟹ *S*(*x* := *a*) ≤ *S'*(*x* := *a'*)
by(*simp add*: *le_fun_def*)

lemma *mono_step'*: *S1* ≤ *S2* ⟹ *C1* ≤ *C2* ⟹ *step' S1 C1* ≤ *step' S2 C2*
unfolding *step'_def*
by(*rule mono2_Step*)
  (*auto simp*: *mono_update mono_aval' asem_def split*: *option.split*)

lemma *mono_step'_top*: *C* ≤ *C'* ⟹ *step'* ⊤ *C* ≤ *step'* ⊤ *C'*
by (*metis mono_step' order_refl*)

lemma *AI_least_pfp*: **assumes** *AI c* = *Some C step'* ⊤ *C'* ≤ *C' strip C'* = *c*
shows *C* ≤ *C'*
by(*rule pfp_bot_least*[*OF* _ _ *assms*(*2*,*3*) *assms*(*1*)[*unfolded AI_def*]])
  (*simp_all add*: *mono_step'_top*)

end

instantiation *acom* :: (*type*) *vars*
begin

definition *vars_acom* = *vars o strip*

instance ..

end

lemma *finite_Cvars*: *finite*(*vars*(*C*::'*a acom*))
by(*simp add*: *vars_acom_def*)

14.7.5 Termination

lemma *pfp_termination*:
fixes *x0* :: *'a::order* **and** *m* :: *'a ⇒ nat*
assumes *mono*: ⋀*x y. I x ⟹ I y ⟹ x ≤ y ⟹ f x ≤ f y*
and *m*: ⋀*x y. I x ⟹ I y ⟹ x < y ⟹ m x > m y*
and *I*: ⋀*x y. I x ⟹ I*(*f x*) **and** *I x0* **and** *x0 ≤ f x0*
shows ∃*x. pfp f x0 = Some x*
proof(*simp add*: *pfp_def*, *rule wf_while_option_Some*[**where** *P* = %*x. I x* & *x ≤ f x*])
  show *wf* {(*y*,*x*). ((*I x* ∧ *x ≤ f x*) ∧ ¬ *f x ≤ x*) ∧ *y = f x*}
    by(*rule wf_subset*[*OF wf_measure*[*of m*]]) (*auto simp*: *m I*)
next
  show *I x0* ∧ *x0 ≤ f x0* **using** ‹*I x0*› ‹*x0 ≤ f x0*› **by** *blast*
next
  fix *x* **assume** *I x* ∧ *x ≤ f x* **thus** *I*(*f x*) ∧ *f x ≤ f*(*f x*)
    by (*blast intro*: *I mono*)
qed

lemma *le_iff_le_annos*: *C1 ≤ C2* ⟷
  strip C1 = strip C2 ∧ (∀ *i<size*(*annos C1*). *annos C1 ! i ≤ annos C2 ! i*)
by(*simp add*: *less_eq_acom_def anno_def*)

locale *Measure1_fun* =
fixes *m* :: *'av::top ⇒ nat*
fixes *h* :: *nat*
assumes *h*: *m x ≤ h*
begin

definition *m_s* :: *'av st ⇒ vname set ⇒ nat* (m_s) **where**
m_s S X = (∑ *x* ∈ *X*. *m*(*S x*))

lemma *m_s_h*: *finite X ⟹ m_s S X ≤ h * card X*
by(*simp add*: *m_s_def*) (*metis mult.commute of_nat_id sum_bounded_above*[*OF h*])

fun *m_o* :: *'av st option ⇒ vname set ⇒ nat* (m_o) **where**
m_o (*Some S*) *X = m_s S X* |
*m_o None X = h * card X + 1*

lemma *m_o_h*: *finite X ⟹ m_o opt X ≤* (*h*card X + 1*)
by(*cases opt*)(*auto simp add*: *m_s_h le_SucI dest*: *m_s_h*)

definition *m_c* :: *'av st option acom* ⇒ *nat* (m_c) **where**
m_c C = *sum_list* (*map* (λ*a*. *m_o a* (*vars C*)) (*annos C*))

Upper complexity bound:

lemma *m_c_h*: *m_c C* ≤ *size*(*annos C*) ∗ (*h* ∗ *card*(*vars C*) + *1*)
proof−
  let *?X* = *vars C* **let** *?n* = *card ?X* **let** *?a* = *size*(*annos C*)
  have *m_c C* = (∑ *i*<*?a*. *m_o* (*annos C* ! *i*) *?X*)
    by(*simp add*: *m_c_def sum_list_sum_nth atLeast0LessThan*)
  also have . . . ≤ (∑ *i*<*?a*. *h* ∗ *?n* + *1*)
    apply(*rule sum_mono*) **using** *m_o_h*[*OF finite_Cvars*] **by** *simp*
  also have . . . = *?a* ∗ (*h* ∗ *?n* + *1*) **by** *simp*
  finally show *?thesis* **.**
qed

end

locale *Measure_fun* = *Measure1_fun* **where** *m*=*m*
  for *m* :: *'av*::*semilattice_sup_top* ⇒ *nat* +
assumes *m2*: *x* < *y* ⟹ *m x* > *m y*
begin

The predicates *top_on_ty a X* that follow describe that any abstract state in *a* maps all variables in *X* to ⊤. This is an important invariant for the termination proof where we argue that only the finitely many variables in the program change. That the others do not change follows because they remain ⊤.

fun *top_on_st* :: *'av st* ⇒ *vname set* ⇒ *bool* (top'_on_s) **where**
top_on_st S X = (∀ *x*∈*X*. *S x* = ⊤)

fun *top_on_opt* :: *'av st option* ⇒ *vname set* ⇒ *bool* (top'_on_o) **where**
top_on_opt (*Some S*) *X* = *top_on_st S X* |
top_on_opt None X = *True*

definition *top_on_acom* :: *'av st option acom* ⇒ *vname set* ⇒ *bool* (top'_on_c)
where
top_on_acom C X = (∀ *a* ∈ *set*(*annos C*). *top_on_opt a X*)

lemma *top_on_top*: *top_on_opt* ⊤ *X*
by(*auto simp*: *top_option_def*)

lemma *top_on_bot*: *top_on_acom* (*bot c*) *X*
by(*auto simp add*: *top_on_acom_def bot_def*)

lemma *top_on_post*: *top_on_acom C X* $\Longrightarrow$ *top_on_opt* (*post C*) *X*
by(*simp add*: *top_on_acom_def post_in_annos*)

lemma *top_on_acom_simps*:
top_on_acom (*SKIP* {*Q*}) *X* = *top_on_opt Q X*
top_on_acom (*x* ::= *e* {*Q*}) *X* = *top_on_opt Q X*
top_on_acom (*C1*;;*C2*) *X* = (*top_on_acom C1 X* ∧ *top_on_acom C2 X*)
top_on_acom (*IF b THEN* {*P1*} *C1 ELSE* {*P2*} *C2* {*Q*}) *X* =
(*top_on_opt P1 X* ∧ *top_on_acom C1 X* ∧ *top_on_opt P2 X* ∧ *top_on_acom C2 X* ∧ *top_on_opt Q X*)
top_on_acom ({*I*} *WHILE b DO* {*P*} *C* {*Q*}) *X* =
(*top_on_opt I X* ∧ *top_on_acom C X* ∧ *top_on_opt P X* ∧ *top_on_opt Q X*)
by(*auto simp add*: *top_on_acom_def*)

lemma *top_on_sup*:
top_on_opt o1 X $\Longrightarrow$ *top_on_opt o2 X* $\Longrightarrow$ *top_on_opt* (*o1* ⊔ *o2*) *X*
apply(*induction o1 o2 rule*: *sup_option.induct*)
apply(*auto*)
done

lemma *top_on_Step*: **fixes** *C* :: *'av st option acom*
assumes !!*x e S*. ⟦*top_on_opt S X*; *x* ∉ *X*; *vars e* ⊆ −*X*⟧ $\Longrightarrow$ *top_on_opt* (*f x e S*) *X*
!!*b S*. *top_on_opt S X* $\Longrightarrow$ *vars b* ⊆ −*X* $\Longrightarrow$ *top_on_opt* (*g b S*) *X*
shows ⟦ *vars C* ⊆ −*X*; *top_on_opt S X*; *top_on_acom C X* ⟧ $\Longrightarrow$ *top_on_acom* (*Step f g S C*) *X*
proof(*induction C arbitrary*: *S*)
qed (*auto simp*: *top_on_acom_simps vars_acom_def top_on_post top_on_sup assms*)

lemma *m1*: *x* ≤ *y* $\Longrightarrow$ *m x* ≥ *m y*
by(*auto simp*: *le_less m2*)

lemma *m_s2_rep*: **assumes** *finite*(*X*) **and** *S1* = *S2 on* −*X* **and** ∀ *x*. *S1 x* ≤ *S2 x* **and** *S1* ≠ *S2*
shows (∑ *x*∈*X*. *m* (*S2 x*)) < (∑ *x*∈*X*. *m* (*S1 x*))
proof−
from *assms*(*3*) **have** *1*: ∀ *x*∈*X*. *m*(*S1 x*) ≥ *m*(*S2 x*) **by** (*simp add*: *m1*)
from *assms*(*2*,*3*,*4*) **have** ∃ *x*∈*X*. *S1 x* < *S2 x*
by(*simp add*: *fun_eq_iff*) (*metis Compl_iff le_neq_trans*)
hence *2*: ∃ *x*∈*X*. *m*(*S1 x*) > *m*(*S2 x*) **by** (*metis m2*)

from *sum_strict_mono_ex1*[*OF* ‹*finite X*› *1 2*]
show ($\sum x \in X$. *m* (*S2 x*)) < ($\sum x \in X$. *m* (*S1 x*)) **.**
qed

lemma *m_s2*: *finite*(*X*) ⟹ *S1* = *S2 on* −*X* ⟹ *S1* < *S2* ⟹ *m_s S1 X* > *m_s S2 X*
apply(*auto simp add*: *less_fun_def m_s_def*)
apply(*simp add*: *m_s2_rep le_fun_def*)
done

lemma *m_o2*: *finite X* ⟹ *top_on_opt o1* (−*X*) ⟹ *top_on_opt o2* (−*X*) ⟹
o1 < *o2* ⟹ *m_o o1 X* > *m_o o2 X*
proof(*induction o1 o2 rule*: *less_eq_option.induct*)
case *1* **thus** *?case* **by** (*auto simp*: *m_s2 less_option_def*)
next
case *2* **thus** *?case* **by**(*auto simp*: *less_option_def le_imp_less_Suc m_s_h*)
next
case *3* **thus** *?case* **by** (*auto simp*: *less_option_def*)
qed

lemma *m_o1*: *finite X* ⟹ *top_on_opt o1* (−*X*) ⟹ *top_on_opt o2* (−*X*) ⟹
o1 ≤ *o2* ⟹ *m_o o1 X* ≥ *m_o o2 X*
by(*auto simp*: *le_less m_o2*)

lemma *m_c2*: *top_on_acom C1* (−*vars C1*) ⟹ *top_on_acom C2* (−*vars C2*) ⟹
C1 < *C2* ⟹ *m_c C1* > *m_c C2*
proof(*auto simp add*: *le_iff_le_annos size_annos_same*[*of C1 C2*] *vars_acom_def less_acom_def*)
let *?X* = *vars*(*strip C2*)
assume *top*: *top_on_acom C1* (− *vars*(*strip C2*)) *top_on_acom C2* (− *vars*(*strip C2*))
and *strip_eq*: *strip C1* = *strip C2*
and *0*: ∀ *i*<*size*(*annos C2*). *annos C1* ! *i* ≤ *annos C2* ! *i*
hence *1*: ∀ *i*<*size*(*annos C2*). *m_o* (*annos C1* ! *i*) *?X* ≥ *m_o* (*annos C2* ! *i*) *?X*
apply (*auto simp*: *all_set_conv_all_nth vars_acom_def top_on_acom_def*)
by (*metis* (*lifting*, *no_types*) *finite_cvars m_o1 size_annos_same2*)
fix *i* **assume** *i*: *i* < *size*(*annos C2*) ¬ *annos C2* ! *i* ≤ *annos C1* ! *i*
have *topo1*: *top_on_opt* (*annos C1* ! *i*) (− *?X*)
using *i*(*1*) *top*(*1*) **by**(*simp add*: *top_on_acom_def size_annos_same*[*OF*

strip_eq])
  have *topo2*: *top_on_opt* (*annos C2* ! *i*) (− *?X*)
   using *i*(*1*) *top*(*2*) **by**(*simp add*: *top_on_acom_def size_annos_same*[*OF strip_eq*])
  from *i* **have** *m_o* (*annos C1* ! *i*) *?X* > *m_o* (*annos C2* ! *i*) *?X* (**is** *?P i*)
    by (*metis 0 less_option_def m_o2*[*OF finite_cvars topo1*] *topo2*)
  hence *2*: ∃ *i* < *size*(*annos C2*). *?P i* **using** ‹*i* < *size*(*annos C2*)› **by** *blast*
  have (∑ *i*<*size*(*annos C2*). *m_o* (*annos C2* ! *i*) *?X*)
        < (∑ *i*<*size*(*annos C2*). *m_o* (*annos C1* ! *i*) *?X*)
    apply(*rule sum_strict_mono_ex1*) **using** *1 2* **by** (*auto*)
  thus *?thesis*
     by(*simp add*: *m_c_def vars_acom_def strip_eq sum_list_sum_nth atLeast0LessThan size_annos_same*[*OF strip_eq*])
qed

end

locale *Abs_Int_fun_measure* =
  Abs_Int_fun_mono **where** *γ*=*γ* + *Measure_fun* **where** *m*=*m*
  for *γ* :: *'av*::*semilattice_sup_top* ⇒ *val set* **and** *m* :: *'av* ⇒ *nat*
begin

lemma *top_on_step'*: *top_on_acom C* (−*vars C*) ⟹ *top_on_acom* (*step'* ⊤ *C*) (−*vars C*)
unfolding *step'_def*
by(*rule top_on_Step*)
  (*auto simp add*: *top_option_def asem_def split*: *option.splits*)

lemma *AI_Some_measure*: ∃ *C*. *AI c* = *Some C*
unfolding *AI_def*
apply(*rule pfp_termination*[**where** *I* = *λC*. *top_on_acom C* (− *vars C*) **and** *m*=*m_c*])
apply(*simp_all add*: *m_c2 mono_step'_top bot_least top_on_bot*)
using *top_on_step'* **apply**(*auto simp add*: *vars_acom_def*)
done

end

Problem: not executable because of the comparison of abstract states, i.e. functions, in the pre-fixpoint computation.

end

14.8 Computable State

theory *Abs_State*
imports *Abs_Int0*
begin

type_synonym *'a st_rep* = (*vname* ∗ *'a*) *list*

fun *fun_rep* :: (*'a*::*top*) *st_rep* ⇒ *vname* ⇒ *'a* **where**
fun_rep [] = (λ*x*. ⊤) |
fun_rep ((*x*,*a*)#*ps*) = (*fun_rep ps*) (*x* := *a*)

lemma *fun_rep_map_of*[*code*]: — original def is too slow
  fun_rep ps = (%*x*. *case map_of ps x of None* ⇒ ⊤ | *Some a* ⇒ *a*)
by(*induction ps rule*: *fun_rep.induct*) *auto*

definition *eq_st* :: (*'a*::*top*) *st_rep* ⇒ *'a st_rep* ⇒ *bool* **where**
eq_st S1 S2 = (*fun_rep S1* = *fun_rep S2*)

hide_type *st* — hide previous def to avoid long names
declare [[*typedef_overloaded*]] — allow quotient types to depend on classes

quotient_type *'a st* = (*'a*::*top*) *st_rep* / *eq_st*
morphisms *rep_st St*
by (*metis eq_st_def equivpI reflpI sympI transpI*)

lift_definition *update* :: (*'a*::*top*) *st* ⇒ *vname* ⇒ *'a* ⇒ *'a st*
  is λ*ps x a*. (*x*,*a*)#*ps*
by(*auto simp*: *eq_st_def*)

lift_definition *fun* :: (*'a*::*top*) *st* ⇒ *vname* ⇒ *'a* **is** *fun_rep*
by(*simp add*: *eq_st_def*)

definition *show_st* :: *vname set* ⇒ (*'a*::*top*) *st* ⇒ (*vname* ∗ *'a*)*set* **where**
show_st X S = (λ*x*. (*x*, *fun S x*)) ‘ *X*

definition *show_acom C* = *map_acom* (*map_option* (*show_st* (*vars*(*strip C*)))) *C*
definition *show_acom_opt* = *map_option show_acom*

lemma *fun_update*[*simp*]: *fun* (*update S x y*) = (*fun S*)(*x*:=*y*)
by *transfer auto*

definition γ*_st* :: ((*'a*::*top*) ⇒ *'b set*) ⇒ *'a st* ⇒ (*vname* ⇒ *'b*) *set* **where**

γ_st γ F = {*f*. ∀ *x*. *f x* ∈ *γ*(*fun F x*)}

instantiation *st* :: (*order_top*) *order*
begin

definition *less_eq_st_rep* :: *'a st_rep* ⇒ *'a st_rep* ⇒ *bool* **where**
less_eq_st_rep ps1 ps2 =
((∀ *x* ∈ *set*(*map fst ps1*) ∪ *set*(*map fst ps2*). *fun_rep ps1 x* ≤ *fun_rep ps2 x*))

lemma *less_eq_st_rep_iff*:
less_eq_st_rep r1 r2 = (∀ *x*. *fun_rep r1 x* ≤ *fun_rep r2 x*)
apply(*auto simp*: *less_eq_st_rep_def fun_rep_map_of split*: *option.split*)
apply (*metis Un_iff map_of_eq_None_iff option.distinct*(*1*))
apply (*metis Un_iff map_of_eq_None_iff option.distinct*(*1*))
done

corollary *less_eq_st_rep_iff_fun*:
less_eq_st_rep r1 r2 = (*fun_rep r1* ≤ *fun_rep r2*)
by (*metis less_eq_st_rep_iff le_fun_def*)

lift_definition *less_eq_st* :: *'a st* ⇒ *'a st* ⇒ *bool* **is** *less_eq_st_rep*
by(*auto simp add*: *eq_st_def less_eq_st_rep_iff*)

definition *less_st* **where** *F* < (*G*::*'a st*) = (*F* ≤ *G* ∧ ¬ *G* ≤ *F*)

instance
proof (*standard*, *goal_cases*)
case *1* **show** *?case* **by**(*rule less_st_def*)
next
case *2* **show** *?case* **by** *transfer* (*auto simp*: *less_eq_st_rep_def*)
next
case *3* **thus** *?case* **by** *transfer* (*metis less_eq_st_rep_iff order_trans*)
next
case *4* **thus** *?case*
by *transfer* (*metis less_eq_st_rep_iff eq_st_def fun_eq_iff antisym*)
qed

end

lemma *le_st_iff*: (*F* ≤ *G*) = (∀ *x*. *fun F x* ≤ *fun G x*)
by *transfer* (*rule less_eq_st_rep_iff*)

fun *map2_st_rep* :: (*'a*::*top* ⇒ *'a* ⇒ *'a*) ⇒ *'a st_rep* ⇒ *'a st_rep* ⇒ *'a*

st_rep **where**
map2_st_rep f [] *ps2* = *map* (%(*x*,*y*). (*x*, *f* ⊤ *y*)) *ps2* |
map2_st_rep f ((*x*,*y*)#*ps1*) *ps2* =
  (*let y2* = *fun_rep ps2 x*
   in (*x*,*f y y2*) # *map2_st_rep f ps1 ps2*)

lemma *fun_rep_map2_rep*[*simp*]: *f* ⊤ ⊤ = ⊤ ⟹
  fun_rep (*map2_st_rep f ps1 ps2*) = (λ*x*. *f* (*fun_rep ps1 x*) (*fun_rep ps2 x*))
apply(*induction f ps1 ps2 rule*: *map2_st_rep.induct*)
apply(*simp add*: *fun_rep_map_of map_of_map fun_eq_iff split*: *option.split*)
apply(*fastforce simp*: *fun_rep_map_of fun_eq_iff split*:*option.splits*)
done

instantiation *st* :: (*semilattice_sup_top*) *semilattice_sup_top*
begin

lift_definition *sup_st* :: ′*a st* ⇒ ′*a st* ⇒ ′*a st* **is** *map2_st_rep* (⊔)
by (*simp add*: *eq_st_def*)

lift_definition *top_st* :: ′*a st* **is** [] **.**

instance
proof (*standard*, *goal_cases*)
  case *1* **show** *?case* **by** *transfer* (*simp add*:*less_eq_st_rep_iff*)
next
  case *2* **show** *?case* **by** *transfer* (*simp add*:*less_eq_st_rep_iff*)
next
  case *3* **thus** *?case* **by** *transfer* (*simp add*:*less_eq_st_rep_iff*)
next
  case *4* **show** *?case* **by** *transfer* (*simp add*:*less_eq_st_rep_iff fun_rep_map_of*)
qed

end

lemma *fun_top*: *fun* ⊤ = (λ*x*. ⊤)
by *transfer simp*

lemma *mono_update*[*simp*]:
  a1 ≤ *a2* ⟹ *S1* ≤ *S2* ⟹ *update S1 x a1* ≤ *update S2 x a2*
by *transfer* (*auto simp add*: *less_eq_st_rep_def*)

lemma *mono_fun*: *S1* ≤ *S2* ⟹ *fun S1 x* ≤ *fun S2 x*
by *transfer* (*simp add*: *less_eq_st_rep_iff*)

locale *Gamma_semilattice* = *Val_semilattice* **where** γ=γ
for γ :: *'av::semilattice_sup_top* ⇒ *val set*
begin

abbreviation γ_s :: *'av st* ⇒ *state set*
where γ_s == γ*_st* γ

abbreviation γ_o :: *'av st option* ⇒ *state set*
where γ_o == γ*_option* γ_s

abbreviation γ_c :: *'av st option acom* ⇒ *state set acom*
where γ_c == *map_acom* γ_o

lemma *gamma_s_top*[*simp*]: γ_s ⊤ = *UNIV*
by(*auto simp*: γ*_st_def fun_top*)

lemma *gamma_o_Top*[*simp*]: γ_o ⊤ = *UNIV*
by (*simp add*: *top_option_def*)

lemma *mono_gamma_s*: *f* ≤ *g* ⟹ γ_s *f* ⊆ γ_s *g*
by(*simp add*:γ*_st_def le_st_iff subset_iff*) (*metis mono_gamma subsetD*)

lemma *mono_gamma_o*:
S1 ≤ *S2* ⟹ γ_o *S1* ⊆ γ_o *S2*
by(*induction S1 S2 rule*: *less_eq_option.induct*)(*simp_all add*: *mono_gamma_s*)

lemma *mono_gamma_c*: *C1* ≤ *C2* ⟹ γ_c *C1* ≤ γ_c *C2*
by (*simp add*: *less_eq_acom_def mono_gamma_o size_annos anno_map_acom size_annos_same*[*of C1 C2*])

lemma *in_gamma_option_iff*:
x ∈ γ*_option r u* ⟷ (∃ *u'*. *u* = *Some u'* ∧ *x* ∈ *r u'*)
by (*cases u*) *auto*

end

end

14.9 Computable Abstract Interpretation

theory *Abs_Int1*
imports *Abs_State*
begin

Abstract interpretation over type *st* instead of functions.

context *Gamma_semilattice*
begin

fun $aval'$ :: $aexp \Rightarrow {}'av\ st \Rightarrow {}'av$ **where**
$aval'\ (N\ i)\ S = num'\ i\ |$
$aval'\ (V\ x)\ S = fun\ S\ x\ |$
$aval'\ (Plus\ a1\ a2)\ S = plus'\ (aval'\ a1\ S)\ (aval'\ a2\ S)$

lemma *aval'_correct*: $s \in \gamma_s\ S \Longrightarrow aval\ a\ s \in \gamma(aval'\ a\ S)$
by (*induction a*) (*auto simp*: *gamma_num' gamma_plus' γ_st_def*)

lemma *gamma_Step_subcomm*: **fixes** *C1 C2* :: *'a::semilattice_sup acom*
  assumes $!!x\ e\ S.\ f1\ x\ e\ (\gamma_o\ S) \subseteq \gamma_o\ (f2\ x\ e\ S)$
  $!!b\ S.\ g1\ b\ (\gamma_o\ S) \subseteq \gamma_o\ (g2\ b\ S)$
  shows $Step\ f1\ g1\ (\gamma_o\ S)\ (\gamma_c\ C) \le \gamma_c\ (Step\ f2\ g2\ S\ C)$
proof(*induction C arbitrary*: *S*)
qed (*auto simp*: *assms intro*!: *mono_gamma_o sup_ge1 sup_ge2*)

lemma *in_gamma_update*: $[\![\ s \in \gamma_s\ S;\ i \in \gamma\ a\ ]\!] \Longrightarrow s(x := i) \in \gamma_s(update\ S\ x\ a)$
by(*simp add*: *γ_st_def*)

end

locale *Abs_Int* = *Gamma_semilattice* **where** $\gamma{=}\gamma$
  for γ :: *'av::semilattice_sup_top* $\Rightarrow$ *val set*
begin

definition $step' = Step$
  $(\lambda x\ e\ S.\ case\ S\ of\ None \Rightarrow None\ |\ Some\ S \Rightarrow Some(update\ S\ x\ (aval'\ e\ S)))$
  $(\lambda b\ S.\ S)$

definition *AI* :: *com* $\Rightarrow$ *'av st option acom option* **where**
$AI\ c = pfp\ (step'\ \top)\ (bot\ c)$

lemma *strip_step'*[*simp*]: $strip(step'\ S\ C) = strip\ C$
by(*simp add*: *step'_def*)

Correctness:

lemma *step_step'*: $step\ (\gamma_o\ S)\ (\gamma_c\ C) \le \gamma_c\ (step'\ S\ C)$

unfolding *step_def step′_def*
by(*rule gamma_Step_subcomm*)
 (*auto simp*: *intro*!: *aval′_correct in_gamma_update split*: *option.splits*)

lemma *AI_correct*: $AI\ c = Some\ C \Longrightarrow CS\ c \le \gamma_c\ C$
proof(*simp add*: *CS_def AI_def*)
 assume *1*: $pfp\ (step'\ \top)\ (bot\ c) = Some\ C$
 have *pfp′*: $step'\ \top\ C \le C$ **by**(*rule pfp_pfp*[*OF 1*])
 have *2*: $step\ (\gamma_o\ \top)\ (\gamma_c\ C) \le \gamma_c\ C$ — transfer the pfp'
 proof(*rule order_trans*)
  show $step\ (\gamma_o\ \top)\ (\gamma_c\ C) \le \gamma_c\ (step'\ \top\ C)$ **by**(*rule step_step′*)
  show ... $\le \gamma_c\ C$ **by** (*metis mono_gamma_c*[*OF pfp′*])
 qed
 have *3*: $strip\ (\gamma_c\ C) = c$ **by**(*simp add*: *strip_pfp*[*OF _ 1*] *step′_def*)
 have $lfp\ c\ (step\ (\gamma_o\ \top)) \le \gamma_c\ C$
  by(*rule lfp_lowerbound*[*simplified*,**where** $f=step\ (\gamma_o\ \top)$, *OF 3 2*])
 thus $lfp\ c\ (step\ UNIV) \le \gamma_c\ C$ **by** *simp*
qed

end

14.9.1 Monotonicity

locale *Abs_Int_mono* = *Abs_Int* +
assumes *mono_plus′*: $a1 \le b1 \Longrightarrow a2 \le b2 \Longrightarrow plus'\ a1\ a2 \le plus'\ b1\ b2$
begin

lemma *mono_aval′*: $S1 \le S2 \Longrightarrow aval'\ e\ S1 \le aval'\ e\ S2$
by(*induction e*) (*auto simp*: *mono_plus′ mono_fun*)

theorem *mono_step′*: $S1 \le S2 \Longrightarrow C1 \le C2 \Longrightarrow step'\ S1\ C1 \le step'\ S2\ C2$
unfolding *step′_def*
by(*rule mono2_Step*) (*auto simp*: *mono_aval′ split*: *option.split*)

lemma *mono_step′_top*: $C \le C' \Longrightarrow step'\ \top\ C \le step'\ \top\ C'$
by (*metis mono_step′ order_refl*)

lemma *AI_least_pfp*: **assumes** $AI\ c = Some\ C$ $step'\ \top\ C' \le C'$ $strip\ C' = c$
shows $C \le C'$
by(*rule pfp_bot_least*[*OF _ _ assms*(*2*,*3*) *assms*(*1*)[*unfolded AI_def*]])
 (*simp_all add*: *mono_step′_top*)

end

14.9.2 Termination

locale *Measure1* =
fixes *m* :: *'av::order_top* ⇒ *nat*
fixes *h* :: *nat*
assumes *h*: *m x* ≤ *h*
begin

definition *m_s* :: *'av st* ⇒ *vname set* ⇒ *nat* (m_s) **where**
m_s S X = (∑ *x* ∈ *X*. *m*(*fun S x*))

lemma *m_s_h*: *finite X* ⟹ *m_s S X* ≤ *h* ∗ *card X*
by(*simp add*: *m_s_def*) (*metis mult.commute of_nat_id sum_bounded_above*[*OF h*])

definition *m_o* :: *'av st option* ⇒ *vname set* ⇒ *nat* (m_o) **where**
m_o opt X = (*case opt of None* ⇒ *h* ∗ *card X* + *1* | *Some S* ⇒ *m_s S X*)

lemma *m_o_h*: *finite X* ⟹ *m_o opt X* ≤ (*h*∗*card X* + *1*)
by(*auto simp add*: *m_o_def m_s_h le_SucI split*: *option.split dest*:*m_s_h*)

definition *m_c* :: *'av st option acom* ⇒ *nat* (m_c) **where**
m_c C = *sum_list* (*map* (λ*a*. *m_o a* (*vars C*)) (*annos C*))

Upper complexity bound:

lemma *m_c_h*: *m_c C* ≤ *size*(*annos C*) ∗ (*h* ∗ *card*(*vars C*) + *1*)
proof−
  let *?X* = *vars C* **let** *?n* = *card ?X* **let** *?a* = *size*(*annos C*)
  have *m_c C* = (∑ *i*<*?a*. *m_o* (*annos C* ! *i*) *?X*)
    by(*simp add*: *m_c_def sum_list_sum_nth atLeast0LessThan*)
  also have . . . ≤ (∑ *i*<*?a*. *h* ∗ *?n* + *1*)
    apply(*rule sum_mono*) **using** *m_o_h*[*OF finite_Cvars*] **by** *simp*
  also have . . . = *?a* ∗ (*h* ∗ *?n* + *1*) **by** *simp*
  finally show *?thesis* **.**
qed

end

fun *top_on_st* :: *'a::order_top st* ⇒ *vname set* ⇒ *bool* (top'_on_s) **where**
top_on_st S X = (∀ *x*∈*X*. *fun S x* = ⊤)

fun *top_on_opt* :: *'a::order_top st option* ⇒ *vname set* ⇒ *bool* (top'_on_o)

where
top_on_opt (*Some S*) *X* = *top_on_st S X* |
top_on_opt None X = *True*

definition *top_on_acom* :: *'a::order_top st option acom* ⇒ *vname set* ⇒ *bool* (top'_on_c) **where**
top_on_acom C X = (∀ *a* ∈ *set*(*annos C*). *top_on_opt a X*)

lemma *top_on_top*: *top_on_opt* (⊤::_ *st option*) *X*
by(*auto simp*: *top_option_def fun_top*)

lemma *top_on_bot*: *top_on_acom* (*bot c*) *X*
by(*auto simp add*: *top_on_acom_def bot_def*)

lemma *top_on_post*: *top_on_acom C X* ⟹ *top_on_opt* (*post C*) *X*
by(*simp add*: *top_on_acom_def post_in_annos*)

lemma *top_on_acom_simps*:
  top_on_acom (*SKIP* {*Q*}) *X* = *top_on_opt Q X*
  top_on_acom (*x* ::= *e* {*Q*}) *X* = *top_on_opt Q X*
  top_on_acom (*C1*;;*C2*) *X* = (*top_on_acom C1 X* ∧ *top_on_acom C2 X*)
  top_on_acom (*IF b THEN* {*P1*} *C1 ELSE* {*P2*} *C2* {*Q*}) *X* =
    (*top_on_opt P1 X* ∧ *top_on_acom C1 X* ∧ *top_on_opt P2 X* ∧ *top_on_acom C2 X* ∧ *top_on_opt Q X*)
  top_on_acom ({*I*} *WHILE b DO* {*P*} *C* {*Q*}) *X* =
  (*top_on_opt I X* ∧ *top_on_acom C X* ∧ *top_on_opt P X* ∧ *top_on_opt Q X*)
by(*auto simp add*: *top_on_acom_def*)

lemma *top_on_sup*:
  top_on_opt o1 X ⟹ *top_on_opt o2 X* ⟹ *top_on_opt* (*o1* ⊔ *o2* :: _ *st option*) *X*
apply(*induction o1 o2 rule*: *sup_option.induct*)
apply(*auto*)
by *transfer simp*

lemma *top_on_Step*: **fixes** *C* :: (*'a::semilattice_sup_top*)*st option acom*
assumes !!*x e S*. ⟦*top_on_opt S X*; *x* ∉ *X*; *vars e* ⊆ −*X*⟧ ⟹ *top_on_opt* (*f x e S*) *X*
        !!*b S*. *top_on_opt S X* ⟹ *vars b* ⊆ −*X* ⟹ *top_on_opt* (*g b S*) *X*
shows ⟦ *vars C* ⊆ −*X*; *top_on_opt S X*; *top_on_acom C X* ⟧ ⟹ *top_on_acom* (*Step f g S C*) *X*
proof(*induction C arbitrary*: *S*)

qed (*auto simp*: *top_on_acom_simps vars_acom_def top_on_post top_on_sup assms*)

locale *Measure* = *Measure1* +
assumes *m2*: $x < y \Longrightarrow m\ x > m\ y$
begin

lemma *m1*: $x \le y \Longrightarrow m\ x \ge m\ y$
by(*auto simp*: *le_less m2*)

lemma *m_s2_rep*: **assumes** *finite*(*X*) **and** *S1* = *S2 on* −*X* **and** ∀ *x*. *S1 x* ≤ *S2 x* **and** *S1* ≠ *S2*
shows ($\sum$ *x*∈*X*. *m* (*S2 x*)) < ($\sum$ *x*∈*X*. *m* (*S1 x*))
proof−
  from *assms*(*3*) **have** *1*: ∀ *x*∈*X*. *m*(*S1 x*) ≥ *m*(*S2 x*) **by** (*simp add*: *m1*)
  from *assms*(*2*,*3*,*4*) **have** ∃ *x*∈*X*. *S1 x* < *S2 x*
    by(*simp add*: *fun_eq_iff*) (*metis Compl_iff le_neq_trans*)
  hence *2*: ∃ *x*∈*X*. *m*(*S1 x*) > *m*(*S2 x*) **by** (*metis m2*)
  from *sum_strict_mono_ex1*[*OF* ‹*finite X*› *1 2*]
  show ($\sum$ *x*∈*X*. *m* (*S2 x*)) < ($\sum$ *x*∈*X*. *m* (*S1 x*)) **.**
qed

lemma *m_s2*: *finite*(*X*) ⟹ *fun S1* = *fun S2 on* −*X*
  ⟹ *S1* < *S2* ⟹ *m_s S1 X* > *m_s S2 X*
apply(*auto simp add*: *less_st_def m_s_def*)
apply (*transfer fixing*: *m*)
apply(*simp add*: *less_eq_st_rep_iff eq_st_def m_s2_rep*)
done

lemma *m_o2*: *finite X* ⟹ *top_on_opt o1* (−*X*) ⟹ *top_on_opt o2* (−*X*) ⟹
  o1 < *o2* ⟹ *m_o o1 X* > *m_o o2 X*
proof(*induction o1 o2 rule*: *less_eq_option.induct*)
  case *1* **thus** *?case* **by** (*auto simp*: *m_o_def m_s2 less_option_def*)
next
  case *2* **thus** *?case* **by**(*auto simp*: *m_o_def less_option_def le_imp_less_Suc m_s_h*)
next
  case *3* **thus** *?case* **by** (*auto simp*: *less_option_def*)
qed

lemma *m_o1*: *finite X* ⟹ *top_on_opt o1* (−*X*) ⟹ *top_on_opt o2* (−*X*) ⟹

o1 ≤ *o2* ⟹ *m_o o1 X* ≥ *m_o o2 X*
by(*auto simp*: *le_less m_o2*)

lemma *m_c2*: *top_on_acom C1* (−*vars C1*) ⟹ *top_on_acom C2* (−*vars C2*) ⟹
C1 < *C2* ⟹ *m_c C1* > *m_c C2*
proof(*auto simp add*: *le_iff_le_annos size_annos_same*[*of C1 C2*] *vars_acom_def less_acom_def*)
let *?X* = *vars*(*strip C2*)
assume *top*: *top_on_acom C1* (− *vars*(*strip C2*)) *top_on_acom C2* (− *vars*(*strip C2*))
and *strip_eq*: *strip C1* = *strip C2*
and *0*: ∀ *i*<*size*(*annos C2*). *annos C1* ! *i* ≤ *annos C2* ! *i*
hence *1*: ∀ *i*<*size*(*annos C2*). *m_o* (*annos C1* ! *i*) *?X* ≥ *m_o* (*annos C2* ! *i*) *?X*
apply (*auto simp*: *all_set_conv_all_nth vars_acom_def top_on_acom_def*)
by (*metis finite_cvars m_o1 size_annos_same2*)
fix *i* **assume** *i*: *i* < *size*(*annos C2*) ¬ *annos C2* ! *i* ≤ *annos C1* ! *i*
have *topo1*: *top_on_opt* (*annos C1* ! *i*) (− *?X*)
using *i*(*1*) *top*(*1*) **by**(*simp add*: *top_on_acom_def size_annos_same*[*OF strip_eq*])
have *topo2*: *top_on_opt* (*annos C2* ! *i*) (− *?X*)
using *i*(*1*) *top*(*2*) **by**(*simp add*: *top_on_acom_def size_annos_same*[*OF strip_eq*])
from *i* **have** *m_o* (*annos C1* ! *i*) *?X* > *m_o* (*annos C2* ! *i*) *?X* (**is** *?P i*)
by (*metis 0 less_option_def m_o2*[*OF finite_cvars topo1*] *topo2*)
hence *2*: ∃ *i* < *size*(*annos C2*). *?P i* **using** ‹*i* < *size*(*annos C2*)› **by** *blast*
have (∑ *i*<*size*(*annos C2*). *m_o* (*annos C2* ! *i*) *?X*)
< (∑ *i*<*size*(*annos C2*). *m_o* (*annos C1* ! *i*) *?X*)
apply(*rule sum_strict_mono_ex1*) **using** *1 2* **by** (*auto*)
thus *?thesis*
by(*simp add*: *m_c_def vars_acom_def strip_eq sum_list_sum_nth atLeast0LessThan size_annos_same*[*OF strip_eq*])
qed

end

locale *Abs_Int_measure* =
Abs_Int_mono **where** *γ*=*γ* + *Measure* **where** *m*=*m*
for *γ* :: *'av*::*semilattice_sup_top* ⇒ *val set* **and** *m* :: *'av* ⇒ *nat*
begin

lemma *top_on_step'*: ⟦ *top_on_acom C* (−*vars C*) ⟧ ⟹ *top_on_acom* (*step'* ⊤ *C*) (−*vars C*)
unfolding *step'_def*
by(*rule top_on_Step*)
  (*auto simp add*: *top_option_def fun_top split*: *option.splits*)

lemma *AI_Some_measure*: ∃ *C*. *AI c* = *Some C*
unfolding *AI_def*
apply(*rule pfp_termination*[**where** *I* = λ*C*. *top_on_acom C* (− *vars C*) **and** *m*=*m_c*])
apply(*simp_all add*: *m_c2 mono_step'_top bot_least top_on_bot*)
using *top_on_step'* **apply**(*auto simp add*: *vars_acom_def*)
done

end

end

14.10 Parity Analysis

theory *Abs_Int1_parity*
imports *Abs_Int1*
begin

datatype *parity* = *Even* | *Odd* | *Either*

Instantiation of class *order* with type *parity*:

instantiation *parity* :: *order*
begin

First the definition of the interface function ≤. Note that the header of the definition must refer to the ascii name (≤) of the constants as *less_eq_parity* and the definition is named *less_eq_parity_def*. Inside the definition the symbolic names can be used.

definition *less_eq_parity* **where**
$x \le y = (y = Either \lor x{=}y)$

We also need <, which is defined canonically:

definition *less_parity* **where**
$x < y = (x \le y \land \neg\ y \le (x{::}parity))$

(The type annotation is necessary to fix the type of the polymorphic predicates.)

Now the instance proof, i.e. the proof that the definition fulfills the axioms (assumptions) of the class. The initial proof-step generates the necessary proof obligations.

instance
proof
  fix *x*::*parity* **show** $x \leq x$ **by**(*auto simp*: *less_eq_parity_def*)
next
  fix $x\ y\ z$ :: *parity* **assume** $x \leq y$ $y \leq z$ **thus** $x \leq z$
    by(*auto simp*: *less_eq_parity_def*)
next
  fix $x\ y$ :: *parity* **assume** $x \leq y$ $y \leq x$ **thus** $x = y$
    by(*auto simp*: *less_eq_parity_def*)
next
  fix $x\ y$ :: *parity* **show** $(x < y) = (x \leq y \wedge \neg\ y \leq x)$ **by**(*rule less_parity_def*)
qed

end

Instantiation of class *semilattice_sup_top* with type *parity*:

instantiation *parity* :: *semilattice_sup_top*
begin

definition *sup_parity* **where**
$x \sqcup y = ($*if* $x = y$ *then* x *else Either*$)$

definition *top_parity* **where**
$\top =$ *Either*

Now the instance proof. This time we take a shortcut with the help of proof method *goal_cases*: it creates cases 1 ... n for the subgoals 1 ... n; in case i, i is also the name of the assumptions of subgoal i and *case?* refers to the conclusion of subgoal i. The class axioms are presented in the same order as in the class definition.

instance
proof (*standard*, *goal_cases*)
  case *1* **show** *?case* **by**(*auto simp*: *less_eq_parity_def sup_parity_def*)
next
  case *2* **show** *?case* **by**(*auto simp*: *less_eq_parity_def sup_parity_def*)
next
  case *3* **thus** *?case* **by**(*auto simp*: *less_eq_parity_def sup_parity_def*)
next
  case *4* **show** *?case* **by**(*auto simp*: *less_eq_parity_def top_parity_def*)
qed

end

Now we define the functions used for instantiating the abstract interpretation locales. Note that the Isabelle terminology is *interpretation*, not *instantiation* of locales, but we use instantiation to avoid confusion with abstract interpretation.

fun *γ_parity* :: *parity* ⇒ *val set* **where**
γ_parity Even = {*i. i mod 2* = *0*} |
γ_parity Odd = {*i. i mod 2* = *1*} |
γ_parity Either = *UNIV*

fun *num_parity* :: *val* ⇒ *parity* **where**
num_parity i = (*if i mod 2* = *0 then Even else Odd*)

fun *plus_parity* :: *parity* ⇒ *parity* ⇒ *parity* **where**
plus_parity Even Even = *Even* |
plus_parity Odd Odd = *Even* |
plus_parity Even Odd = *Odd* |
plus_parity Odd Even = *Odd* |
plus_parity Either y = *Either* |
plus_parity x Either = *Either*

First we instantiate the abstract value interface and prove that the functions on type *parity* have all the necessary properties:

global_interpretation *Val_semilattice*
where γ = *γ_parity* **and** *num′* = *num_parity* **and** *plus′* = *plus_parity*
proof (*standard*, *goal_cases*)

subgoals are the locale axioms

  case *1* **thus** *?case* **by**(*auto simp*: *less_eq_parity_def*)
next
  case *2* **show** *?case* **by**(*auto simp*: *top_parity_def*)
next
  case *3* **show** *?case* **by** *auto*
next
  case (*4* _ *a1* _ *a2*) **thus** *?case*
    by (*induction a1 a2 rule*: *plus_parity.induct*)
      (*auto simp add*: *mod_add_eq* [*symmetric*])
qed

In case 4 we needed to refer to particular variables. Writing (i x y z) fixes the names of the variables in case i to be x, y and z in the left-to-right order in which the variables occur in the subgoal. Underscores are anonymous placeholders for variable names we don't care to fix.

Instantiating the abstract interpretation locale requires no more proofs (they happened in the instatiation above) but delivers the instantiated abstract interpreter which we call *AI_parity*:

```
global_interpretation Abs_Int
where γ = γ_parity and num' = num_parity and plus' = plus_parity
defines aval_parity = aval' and step_parity = step' and AI_parity = AI
..
```

14.10.1 Tests

```
definition test1_parity =
  ''x'' ::= N 1;;
  WHILE Less (V ''x'') (N 100) DO ''x'' ::= Plus (V ''x'') (N 2)
value show_acom (the(AI_parity test1_parity))

definition test2_parity =
  ''x'' ::= N 1;;
  WHILE Less (V ''x'') (N 100) DO ''x'' ::= Plus (V ''x'') (N 3)

definition steps c i = ((step_parity ⊤) ^^ i) (bot c)

value show_acom (steps test2_parity 0)
value show_acom (steps test2_parity 1)
value show_acom (steps test2_parity 2)
value show_acom (steps test2_parity 3)
value show_acom (steps test2_parity 4)
value show_acom (steps test2_parity 5)
value show_acom (steps test2_parity 6)
value show_acom (the(AI_parity test2_parity))
```

14.10.2 Termination

```
global_interpretation Abs_Int_mono
where γ = γ_parity and num' = num_parity and plus' = plus_parity
proof (standard, goal_cases)
  case (1 _ a1 _ a2) thus ?case
    by(induction a1 a2 rule: plus_parity.induct)
      (auto simp add:less_eq_parity_def)
qed

definition m_parity :: parity ⇒ nat where
m_parity x = (if x = Either then 0 else 1)

global_interpretation Abs_Int_measure
```

where $\gamma = \gamma\_parity$ **and** $num' = num\_parity$ **and** $plus' = plus\_parity$
and $m = m\_parity$ **and** $h = 1$
proof (*standard*, *goal_cases*)
  case *1* **thus** *?case* **by**(*auto simp add*: *m_parity_def less_eq_parity_def*)
next
  case *2* **thus** *?case* **by**(*auto simp add*: *m_parity_def less_eq_parity_def less_parity_def*)
qed

thm *AI_Some_measure*

end

14.11 Constant Propagation

theory *Abs_Int1_const*
imports *Abs_Int1*
begin

datatype *const* = *Const val* | *Any*

fun γ*_const* **where**
γ*_const* (*Const i*) = {*i*} |
γ*_const* (*Any*) = *UNIV*

fun *plus_const* **where**
plus_const (*Const i*) (*Const j*) = *Const*(*i*+*j*) |
plus_const _ _ = *Any*

lemma *plus_const_cases*: *plus_const a1 a2* =
  (*case* (*a1*,*a2*) *of* (*Const i*, *Const j*) ⇒ *Const*(*i*+*j*) | _ ⇒ *Any*)
by(*auto split*: *prod.split const.split*)

instantiation *const* :: *semilattice_sup_top*
begin

fun *less_eq_const* **where** $x \le y = (y = Any \mid x{=}y)$

definition $x < (y{::}const) = (x \le y \;\&\; \neg\, y \le x)$

fun *sup_const* **where** $x \sqcup y = (if\ x{=}y\ then\ x\ else\ Any)$

definition $\top = Any$

```
instance
proof (standard, goal_cases)
  case 1 thus ?case by (rule less_const_def)
next
  case (2 x) show ?case by (cases x) simp_all
next
  case (3 x y z) thus ?case by(cases z, cases y, cases x, simp_all)
next
  case (4 x y) thus ?case by(cases x, cases y, simp_all, cases y, simp_all)
next
  case (6 x y) thus ?case by(cases x, cases y, simp_all)
next
  case (5 x y) thus ?case by(cases y, cases x, simp_all)
next
  case (7 x y z) thus ?case by(cases z, cases y, cases x, simp_all)
next
  case 8 thus ?case by(simp add: top_const_def)
qed

end


global_interpretation Val_semilattice
where γ = γ_const and num' = Const and plus' = plus_const
proof (standard, goal_cases)
  case (1 a b) thus ?case
    by(cases a, cases b, simp, simp, cases b, simp, simp)
next
  case 2 show ?case by(simp add: top_const_def)
next
  case 3 show ?case by simp
next
  case 4 thus ?case by(auto simp: plus_const_cases split: const.split)
qed

global_interpretation Abs_Int
where γ = γ_const and num' = Const and plus' = plus_const
defines AI_const = AI and step_const = step' and aval'_const = aval'
..
```

14.11.1 Tests

```
definition steps c i = (step_const ⊤ ^^ i) (bot c)
```

value *show_acom* (*steps test1_const 0*)
value *show_acom* (*steps test1_const 1*)
value *show_acom* (*steps test1_const 2*)
value *show_acom* (*steps test1_const 3*)
value *show_acom* (*the*(*AI_const test1_const*))

value *show_acom* (*the*(*AI_const test2_const*))
value *show_acom* (*the*(*AI_const test3_const*))

value *show_acom* (*steps test4_const 0*)
value *show_acom* (*steps test4_const 1*)
value *show_acom* (*steps test4_const 2*)
value *show_acom* (*steps test4_const 3*)
value *show_acom* (*steps test4_const 4*)
value *show_acom* (*the*(*AI_const test4_const*))

value *show_acom* (*steps test5_const 0*)
value *show_acom* (*steps test5_const 1*)
value *show_acom* (*steps test5_const 2*)
value *show_acom* (*steps test5_const 3*)
value *show_acom* (*steps test5_const 4*)
value *show_acom* (*steps test5_const 5*)
value *show_acom* (*steps test5_const 6*)
value *show_acom* (*the*(*AI_const test5_const*))

value *show_acom* (*steps test6_const 0*)
value *show_acom* (*steps test6_const 1*)
value *show_acom* (*steps test6_const 2*)
value *show_acom* (*steps test6_const 3*)
value *show_acom* (*steps test6_const 4*)
value *show_acom* (*steps test6_const 5*)
value *show_acom* (*steps test6_const 6*)
value *show_acom* (*steps test6_const 7*)
value *show_acom* (*steps test6_const 8*)
value *show_acom* (*steps test6_const 9*)
value *show_acom* (*steps test6_const 10*)
value *show_acom* (*steps test6_const 11*)
value *show_acom* (*steps test6_const 12*)
value *show_acom* (*steps test6_const 13*)
value *show_acom* (*the*(*AI_const test6_const*))

Monotonicity:

global_interpretation *Abs_Int_mono*
where $\gamma = \gamma$*_const* **and** $num' =$ *Const* **and** $plus' =$ *plus_const*

proof (*standard*, *goal_cases*)
  case *1* **thus** *?case* **by**(*auto simp*: *plus_const_cases split*: *const.split*)
qed

Termination:

definition *m_const* :: *const* ⇒ *nat* **where**
m_const x = (*if x* = *Any then 0 else 1*)

global_interpretation *Abs_Int_measure*
where γ = γ*_const* **and** *num'* = *Const* **and** *plus'* = *plus_const*
and *m* = *m_const* **and** *h* = *1*
proof (*standard*, *goal_cases*)
  case *1* **thus** *?case* **by**(*auto simp*: *m_const_def split*: *const.splits*)
next
  case *2* **thus** *?case* **by**(*auto simp*: *m_const_def less_const_def split*: *const.splits*)
qed

thm *AI_Some_measure*

end

14.12 Backward Analysis of Expressions

theory *Abs_Int2*
imports *Abs_Int1*
begin

instantiation *prod* :: (*order*,*order*) *order*
begin

definition *less_eq_prod p1 p2* = (*fst p1* ≤ *fst p2* ∧ *snd p1* ≤ *snd p2*)
definition *less_prod p1 p2* = (*p1* ≤ *p2* ∧ ¬ *p2* ≤ (*p1*::*'a***'b*))

instance
proof (*standard*, *goal_cases*)
  case *1* **show** *?case* **by**(*rule less_prod_def*)
next
  case *2* **show** *?case* **by**(*simp add*: *less_eq_prod_def*)
next
  case *3* **thus** *?case* **unfolding** *less_eq_prod_def* **by**(*metis order_trans*)
next
  case *4* **thus** *?case* **by**(*simp add*: *less_eq_prod_def*)(*metis eq_iff surjective_pairing*)

qed

end

14.12.1 Extended Framework

subclass (**in** *bounded_lattice*) *semilattice_sup_top* **..**

locale *Val_lattice_gamma* = *Gamma_semilattice* **where** $\gamma = \gamma$
  for γ :: *'av::bounded_lattice* ⇒ *val set* +
assumes *inter_gamma_subset_gamma_inf*:
  γ *a1* ∩ γ *a2* ⊆ γ(*a1* ⊓ *a2*)
and *gamma_bot*[*simp*]: γ ⊥ = {}
begin

lemma *in_gamma_inf*: *x* ∈ γ *a1* ⟹ *x* ∈ γ *a2* ⟹ *x* ∈ γ(*a1* ⊓ *a2*)
by (*metis IntI inter_gamma_subset_gamma_inf subsetD*)

lemma *gamma_inf*: γ(*a1* ⊓ *a2*) = γ *a1* ∩ γ *a2*
by(*rule equalityI*[*OF _ inter_gamma_subset_gamma_inf*])
  (*metis inf_le1 inf_le2 le_inf_iff mono_gamma*)

end

locale *Val_inv* = *Val_lattice_gamma* **where** $\gamma = \gamma$
   for γ :: *'av::bounded_lattice* ⇒ *val set* +
fixes *test_num'* :: *val* ⇒ *'av* ⇒ *bool*
and *inv_plus'* :: *'av* ⇒ *'av* ⇒ *'av* ⇒ *'av* ∗ *'av*
and *inv_less'* :: *bool* ⇒ *'av* ⇒ *'av* ⇒ *'av* ∗ *'av*
assumes *test_num'*: *test_num' i a* = (*i* ∈ γ *a*)
and *inv_plus'*: *inv_plus' a a1 a2* = (a_1',a_2') ⟹
  i1 ∈ γ *a1* ⟹ *i2* ∈ γ *a2* ⟹ *i1*+*i2* ∈ γ *a* ⟹ *i1* ∈ γ a_1' ∧ *i2* ∈ γ a_2'
and *inv_less'*: *inv_less'* (*i1*<*i2*) *a1 a2* = (a_1',a_2') ⟹
  i1 ∈ γ *a1* ⟹ *i2* ∈ γ *a2* ⟹ *i1* ∈ γ a_1' ∧ *i2* ∈ γ a_2'

locale *Abs_Int_inv* = *Val_inv* **where** $\gamma = \gamma$
  for γ :: *'av::bounded_lattice* ⇒ *val set*
begin

lemma *in_gamma_sup_UpI*:
  s ∈ γ_o *S1* ∨ *s* ∈ γ_o *S2* ⟹ *s* ∈ γ_o(*S1* ⊔ *S2*)
by (*metis* (*opaque_lifting*, *no_types*) *sup_ge1 sup_ge2 mono_gamma_o*

subsetD)

fun *aval''* :: *aexp* ⇒ *'av st option* ⇒ *'av* **where**
aval'' e None = ⊥ |
aval'' e (*Some S*) = *aval' e S*

lemma *aval''_correct*: $s \in \gamma_o$ *S* ⟹ *aval a s* ∈ γ(*aval'' a S*)
by(*cases S*)(*auto simp add*: *aval'_correct split*: *option.splits*)

14.12.2 Backward analysis

fun *inv_aval'* :: *aexp* ⇒ *'av* ⇒ *'av st option* ⇒ *'av st option* **where**
inv_aval' (*N n*) *a S* = (*if test_num' n a then S else None*) |
inv_aval' (*V x*) *a S* = (*case S of None* ⇒ *None* | *Some S* ⇒
  let a' = *fun S x* ⊓ *a in*
  if a' = ⊥ *then None else Some*(*update S x a'*)) |
inv_aval' (*Plus e1 e2*) *a S* =
 (*let* (*a1*,*a2*) = *inv_plus' a* (*aval'' e1 S*) (*aval'' e2 S*)
  in inv_aval' e1 a1 (*inv_aval' e2 a2 S*))

The test for *bot* in the *V*-case is important: *bot* indicates that a variable has no possible values, i.e. that the current program point is unreachable. But then the abstract state should collapse to *None*. Put differently, we maintain the invariant that in an abstract state of the form *Some s*, all variables are mapped to non-*bot* values. Otherwise the (pointwise) sup of two abstract states, one of which contains *bot* values, may produce too large a result, thus making the analysis less precise.

fun *inv_bval'* :: *bexp* ⇒ *bool* ⇒ *'av st option* ⇒ *'av st option* **where**
inv_bval' (*Bc v*) *res S* = (*if v=res then S else None*) |
inv_bval' (*Not b*) *res S* = *inv_bval' b* (¬ *res*) *S* |
inv_bval' (*And b1 b2*) *res S* =
  (*if res then inv_bval' b1 True* (*inv_bval' b2 True S*)
   else inv_bval' b1 False S ⊔ *inv_bval' b2 False S*) |
inv_bval' (*Less e1 e2*) *res S* =
  (*let* (*a1*,*a2*) = *inv_less' res* (*aval'' e1 S*) (*aval'' e2 S*)
   in inv_aval' e1 a1 (*inv_aval' e2 a2 S*))

lemma *inv_aval'_correct*: $s \in \gamma_o$ *S* ⟹ *aval e s* ∈ γ *a* ⟹ $s \in \gamma_o$ (*inv_aval' e a S*)
proof(*induction e arbitrary*: *a S*)
  case *N* **thus** *?case* **by** *simp* (*metis test_num'*)
next
  case (*V x*)
  obtain *S'* **where** *S* = *Some S'* **and** $s \in \gamma_s$ *S'* **using** ‹$s \in \gamma_o$ *S*›

by(*auto simp*: *in_gamma_option_iff*)
moreover hence $s\ x \in \gamma$ (*fun* S' *x*)
by(*simp add*: γ*_st_def*)
moreover have $s\ x \in \gamma\ a$ **using** *V*(*2*) **by** *simp*
ultimately show *?case*
by(*simp add*: *Let_def* γ*_st_def*)
(*metis mono_gamma emptyE in_gamma_inf gamma_bot subset_empty*)
next
case (*Plus e1 e2*) **thus** *?case*
using *inv_plus*′[*OF _ aval*′′*_correct aval*′′*_correct*]
by (*auto split*: *prod.split*)
qed

lemma *inv_bval*′*_correct*: $s \in \gamma_o\ S \Longrightarrow bv = bval\ b\ s \Longrightarrow s \in \gamma_o$(*inv_bval*′ *b bv S*)
proof(*induction b arbitrary*: *S bv*)
case *Bc* **thus** *?case* **by** *simp*
next
case (*Not b*) **thus** *?case* **by** *simp*
next
case (*And b1 b2*) **thus** *?case*
by *simp* (*metis And*(*1*) *And*(*2*) *in_gamma_sup_UpI*)
next
case (*Less e1 e2*) **thus** *?case*
apply *hypsubst_thin*
apply (*auto split*: *prod.split*)
apply (*metis* (*lifting*) *inv_aval*′*_correct aval*′′*_correct inv_less*′)
done
qed

definition *step*′ = *Step*
($\lambda x\ e\ S$. *case S of None* ⇒ *None* | *Some S* ⇒ *Some*(*update S x* (*aval*′ *e S*)))
($\lambda b\ S$. *inv_bval*′ *b True S*)

definition *AI* :: *com* ⇒ ′*av st option acom option* **where**
AI c = *pfp* (*step*′ ⊤) (*bot c*)

lemma *strip_step*′[*simp*]: *strip*(*step*′ *S c*) = *strip c*
by(*simp add*: *step*′*_def*)

lemma *top_on_inv_aval*′: ⟦ *top_on_opt S X*; *vars e* ⊆ −*X* ⟧ ⟹ *top_on_opt* (*inv_aval*′ *e a S*) *X*
by(*induction e arbitrary*: *a S*) (*auto simp*: *Let_def split*: *option.splits prod.split*)

lemma *top_on_inv_bval′*: ⟦*top_on_opt S X*; *vars b* ⊆ −*X*⟧ ⟹ *top_on_opt* (*inv_bval′ b r S*) *X*
by(*induction b arbitrary*: *r S*) (*auto simp*: *top_on_inv_aval′ top_on_sup split*: *prod.split*)

lemma *top_on_step′*: *top_on_acom C* (− *vars C*) ⟹ *top_on_acom* (*step′* ⊤ *C*) (− *vars C*)
unfolding *step′_def*
by(*rule top_on_Step*)
(*auto simp add*: *top_on_top top_on_inv_bval′ split*: *option.split*)

14.12.3 Correctness

lemma *step_step′*: *step* (γ_o *S*) (γ_c *C*) ≤ γ_c (*step′ S C*)
unfolding *step_def step′_def*
by(*rule gamma_Step_subcomm*)
(*auto simp*: *intro*!: *aval′_correct inv_bval′_correct in_gamma_update split*: *option.splits*)

lemma *AI_correct*: *AI c* = *Some C* ⟹ *CS c* ≤ γ_c *C*
proof(*simp add*: *CS_def AI_def*)
assume *1*: *pfp* (*step′* ⊤) (*bot c*) = *Some C*
have *pfp′*: *step′* ⊤ *C* ≤ *C* **by**(*rule pfp_pfp*[*OF 1*])
have *2*: *step* (γ_o ⊤) (γ_c *C*) ≤ γ_c *C* — transfer the pfp’
proof(*rule order_trans*)
show *step* (γ_o ⊤) (γ_c *C*) ≤ γ_c (*step′* ⊤ *C*) **by**(*rule step_step′*)
show ... ≤ γ_c *C* **by** (*metis mono_gamma_c*[*OF pfp′*])
qed
have *3*: *strip* (γ_c *C*) = *c* **by**(*simp add*: *strip_pfp*[*OF _ 1*] *step′_def*)
have *lfp c* (*step* (γ_o ⊤)) ≤ γ_c *C*
by(*rule lfp_lowerbound*[*simplified*,**where** *f*=*step* (γ_o ⊤), *OF 3 2*])
thus *lfp c* (*step UNIV*) ≤ γ_c *C* **by** *simp*
qed

end

14.12.4 Monotonicity

locale *Abs_Int_inv_mono* = *Abs_Int_inv* +
assumes *mono_plus′*: *a1* ≤ *b1* ⟹ *a2* ≤ *b2* ⟹ *plus′ a1 a2* ≤ *plus′ b1 b2*
and *mono_inv_plus′*: *a1* ≤ *b1* ⟹ *a2* ≤ *b2* ⟹ *r* ≤ *r′* ⟹
inv_plus′ r a1 a2 ≤ *inv_plus′ r′ b1 b2*
and *mono_inv_less′*: *a1* ≤ *b1* ⟹ *a2* ≤ *b2* ⟹

inv_less′ bv a1 a2 ≤ *inv_less′ bv b1 b2*
begin

lemma *mono_aval′*:
  S1 ≤ *S2* ⟹ *aval′ e S1* ≤ *aval′ e S2*
by(*induction e*) (*auto simp*: *mono_plus′ mono_fun*)

lemma *mono_aval′′*:
  S1 ≤ *S2* ⟹ *aval′′ e S1* ≤ *aval′′ e S2*
apply(*cases S1*)
 apply *simp*
apply(*cases S2*)
 apply *simp*
by (*simp add*: *mono_aval′*)

lemma *mono_inv_aval′*: *r1* ≤ *r2* ⟹ *S1* ≤ *S2* ⟹ *inv_aval′ e r1 S1* ≤ *inv_aval′ e r2 S2*
apply(*induction e arbitrary*: *r1 r2 S1 S2*)
  apply(*auto simp*: *test_num′ Let_def inf_mono split*: *option.splits prod.splits*)
   apply (*metis mono_gamma subsetD*)
  apply (*metis le_bot inf_mono le_st_iff*)
 apply (*metis inf_mono mono_update le_st_iff*)
apply(*metis mono_aval′′ mono_inv_plus′*[*simplified less_eq_prod_def*] *fst_conv snd_conv*)
done

lemma *mono_inv_bval′*: *S1* ≤ *S2* ⟹ *inv_bval′ b bv S1* ≤ *inv_bval′ b bv S2*
apply(*induction b arbitrary*: *bv S1 S2*)
   apply(*simp*)
  apply(*simp*)
 apply *simp*
 apply(*metis order_trans*[*OF _ sup_ge1*] *order_trans*[*OF _ sup_ge2*])
apply (*simp split*: *prod.splits*)
apply(*metis mono_aval′′ mono_inv_aval′ mono_inv_less′*[*simplified less_eq_prod_def*] *fst_conv snd_conv*)
done

theorem *mono_step′*: *S1* ≤ *S2* ⟹ *C1* ≤ *C2* ⟹ *step′ S1 C1* ≤ *step′ S2 C2*
unfolding *step′_def*
by(*rule mono2_Step*) (*auto simp*: *mono_aval′ mono_inv_bval′ split*: *option.split*)

lemma *mono_step′_top*: *C1* ≤ *C2* ⟹ *step′* ⊤ *C1* ≤ *step′* ⊤ *C2*
by (*metis mono_step′ order_refl*)

end

end

14.13 Interval Analysis

theory *Abs_Int2_ivl*
imports *Abs_Int2*
begin

type_synonym *eint* = *int extended*
type_synonym *eint2* = *eint* ∗ *eint*

definition *γ_rep* :: *eint2* ⇒ *int set* **where**
γ_rep p = (*let* (*l*,*h*) = *p in* {*i*. *l* ≤ *Fin i* ∧ *Fin i* ≤ *h*})

definition *eq_ivl* :: *eint2* ⇒ *eint2* ⇒ *bool* **where**
eq_ivl p1 p2 = (*γ_rep p1* = *γ_rep p2*)

lemma *refl_eq_ivl*[*simp*]: *eq_ivl p p*
by(*auto simp*: *eq_ivl_def*)

quotient_type *ivl* = *eint2* / *eq_ivl*
by(*rule equivpI*)(*auto simp*: *reflp_def symp_def transp_def eq_ivl_def*)

abbreviation *ivl_abbr* :: *eint* ⇒ *eint* ⇒ *ivl* ([_, _]) **where**
[*l*,*h*] == *abs_ivl*(*l*,*h*)

lift_definition *γ_ivl* :: *ivl* ⇒ *int set* **is** *γ_rep*
by(*simp add*: *eq_ivl_def*)

lemma *γ_ivl_nice*: *γ_ivl*[*l*,*h*] = {*i*. *l* ≤ *Fin i* ∧ *Fin i* ≤ *h*}
by *transfer* (*simp add*: *γ_rep_def*)

lift_definition *num_ivl* :: *int* ⇒ *ivl* **is** *λi*. (*Fin i*, *Fin i*) **.**

lift_definition *in_ivl* :: *int* ⇒ *ivl* ⇒ *bool*
  is *λi* (*l*,*h*). *l* ≤ *Fin i* ∧ *Fin i* ≤ *h*
by(*auto simp*: *eq_ivl_def γ_rep_def*)

lemma *in_ivl_nice*: *in_ivl i* [*l*,*h*] = (*l* ≤ *Fin i* ∧ *Fin i* ≤ *h*)

by *transfer simp*

definition *is_empty_rep* :: *eint2* ⇒ *bool* **where**
is_empty_rep p = (*let* (*l*,*h*) = *p in l*>*h* | *l*=*Pinf* & *h*=*Pinf* | *l*=*Minf* & *h*=*Minf*)

lemma *γ_rep_cases*: *γ_rep p* = (*case p of* (*Fin i*,*Fin j*) => {*i*..*j*} | (*Fin i*,*Pinf*) => {*i*..} |
(*Minf*,*Fin i*) ⇒ {..*i*} | (*Minf*,*Pinf*) ⇒ *UNIV* | _ ⇒ {})
by(*auto simp add*: *γ_rep_def split*: *prod.splits extended.splits*)

lift_definition *is_empty_ivl* :: *ivl* ⇒ *bool* **is** *is_empty_rep*
apply(*auto simp*: *eq_ivl_def γ_rep_cases is_empty_rep_def*)
apply(*auto simp*: *not_less less_eq_extended_case split*: *extended.splits*)
done

lemma *eq_ivl_iff*: *eq_ivl p1 p2* = (*is_empty_rep p1* & *is_empty_rep p2* | *p1* = *p2*)
by(*auto simp*: *eq_ivl_def is_empty_rep_def γ_rep_cases Icc_eq_Icc split*: *prod.splits extended.splits*)

definition *empty_rep* :: *eint2* **where** *empty_rep* = (*Pinf*,*Minf*)

lift_definition *empty_ivl* :: *ivl* **is** *empty_rep* **.**

lemma *is_empty_empty_rep*[*simp*]: *is_empty_rep empty_rep*
by(*auto simp add*: *is_empty_rep_def empty_rep_def*)

lemma *is_empty_rep_iff*: *is_empty_rep p* = (*γ_rep p* = {})
by(*auto simp add*: *γ_rep_cases is_empty_rep_def split*: *prod.splits extended.splits*)

declare *is_empty_rep_iff*[*THEN iffD1*, *simp*]

instantiation *ivl* :: *semilattice_sup_top*
begin

definition *le_rep* :: *eint2* ⇒ *eint2* ⇒ *bool* **where**
le_rep p1 p2 = (*let* (*l1*,*h1*) = *p1*; (*l2*,*h2*) = *p2 in*
if is_empty_rep(*l1*,*h1*) *then True else*
if is_empty_rep(*l2*,*h2*) *then False else l1* ≥ *l2* & *h1* ≤ *h2*)

lemma *le_iff_subset*: *le_rep p1 p2* ⟷ *γ_rep p1* ⊆ *γ_rep p2*

apply *rule*
apply(*auto simp*: *is_empty_rep_def le_rep_def γ_rep_def split*: *if_splits prod.splits*)[*1*]
apply(*auto simp*: *is_empty_rep_def γ_rep_cases le_rep_def*)
apply(*auto simp*: *not_less split*: *extended.splits*)
done

lift_definition *less_eq_ivl* :: *ivl* ⇒ *ivl* ⇒ *bool* **is** *le_rep*
by(*auto simp*: *eq_ivl_def le_iff_subset*)

definition *less_ivl* **where** *i1* < *i2* = (*i1* ≤ *i2* ∧ ¬ *i2* ≤ (*i1*::*ivl*))

lemma *le_ivl_iff_subset*: *iv1* ≤ *iv2* ⟷ *γ_ivl iv1* ⊆ *γ_ivl iv2*
by *transfer* (*rule le_iff_subset*)

definition *sup_rep* :: *eint2* ⇒ *eint2* ⇒ *eint2* **where**
sup_rep p1 p2 = (*if is_empty_rep p1 then p2 else if is_empty_rep p2 then p1*
else let (*l1*,*h1*) = *p1*; (*l2*,*h2*) = *p2 in* (*min l1 l2*, *max h1 h2*))

lift_definition *sup_ivl* :: *ivl* ⇒ *ivl* ⇒ *ivl* **is** *sup_rep*
by(*auto simp*: *eq_ivl_iff sup_rep_def*)

lift_definition *top_ivl* :: *ivl* **is** (*Minf*,*Pinf*) **.**

lemma *is_empty_min_max*:
¬ *is_empty_rep* (*l1*,*h1*) ⟹ ¬ *is_empty_rep* (*l2*, *h2*) ⟹ ¬ *is_empty_rep* (*min l1 l2*, *max h1 h2*)
by(*auto simp add*: *is_empty_rep_def max_def min_def split*: *if_splits*)

instance
proof (*standard*, *goal_cases*)
case *1* **show** *?case* **by** (*rule less_ivl_def*)
next
case *2* **show** *?case* **by** *transfer* (*simp add*: *le_rep_def split*: *prod.splits*)
next
case *3* **thus** *?case* **by** *transfer* (*auto simp*: *le_rep_def split*: *if_splits*)
next
case *4* **thus** *?case* **by** *transfer* (*auto simp*: *le_rep_def eq_ivl_iff split*: *if_splits*)
next
case *5* **thus** *?case* **by** *transfer* (*auto simp add*: *le_rep_def sup_rep_def is_empty_min_max*)
next

  case *6* **thus** *?case* **by** *transfer* (*auto simp add*: *le_rep_def sup_rep_def is_empty_min_max*)
next
  case *7* **thus** *?case* **by** *transfer* (*auto simp add*: *le_rep_def sup_rep_def*)
next
  case *8* **show** *?case* **by** *transfer* (*simp add*: *le_rep_def is_empty_rep_def*)
qed

end

Implement (naive) executable equality:

instantiation *ivl* :: *equal*
begin

definition *equal_ivl* **where**
equal_ivl i1 (*i2*::*ivl*) = (*i1*≤*i2* ∧ *i2* ≤ *i1*)

instance
proof (*standard*, *goal_cases*)
  case *1* **show** *?case* **by**(*simp add*: *equal_ivl_def eq_iff*)
qed

end

lemma [*simp*]: **fixes** *x* :: *'a*::*linorder extended* **shows** (¬ *x* < *Pinf*) = (*x* = *Pinf*)
by(*simp add*: *not_less*)
lemma [*simp*]: **fixes** *x* :: *'a*::*linorder extended* **shows** (¬ *Minf* < *x*) = (*x* = *Minf*)
by(*simp add*: *not_less*)

instantiation *ivl* :: *bounded_lattice*
begin

definition *inf_rep* :: *eint2* ⇒ *eint2* ⇒ *eint2* **where**
inf_rep p1 p2 = (*let* (*l1*,*h1*) = *p1*; (*l2*,*h2*) = *p2 in* (*max l1 l2*, *min h1 h2*))

lemma *γ_inf_rep*: *γ_rep*(*inf_rep p1 p2*) = *γ_rep p1* ∩ *γ_rep p2*
by(*auto simp*:*inf_rep_def γ_rep_cases split*: *prod.splits extended.splits*)

lift_definition *inf_ivl* :: *ivl* ⇒ *ivl* ⇒ *ivl* **is** *inf_rep*
by(*auto simp*: *γ_inf_rep eq_ivl_def*)

lemma *γ_inf*: *γ_ivl* (*iv1* ⊓ *iv2*) = *γ_ivl iv1* ∩ *γ_ivl iv2*

```
by transfer (rule γ_inf_rep)

definition ⊥ = empty_ivl

instance
proof (standard, goal_cases)
  case 1 thus ?case by (simp add: γ_inf le_ivl_iff_subset)
next
  case 2 thus ?case by (simp add: γ_inf le_ivl_iff_subset)
next
  case 3 thus ?case by (simp add: γ_inf le_ivl_iff_subset)
next
  case 4 show ?case
    unfolding bot_ivl_def by transfer (auto simp: le_iff_subset)
qed

end


lemma eq_ivl_empty: eq_ivl p empty_rep = is_empty_rep p
by (metis eq_ivl_iff is_empty_empty_rep)

lemma le_ivl_nice: [l1,h1] ≤ [l2,h2] ⟷
  (if [l1,h1] = ⊥ then True else
   if [l2,h2] = ⊥ then False else l1 ≥ l2 & h1 ≤ h2)
unfolding bot_ivl_def by transfer (simp add: le_rep_def eq_ivl_empty)

lemma sup_ivl_nice: [l1,h1] ⊔ [l2,h2] =
  (if [l1,h1] = ⊥ then [l2,h2] else
   if [l2,h2] = ⊥ then [l1,h1] else [min l1 l2,max h1 h2])
unfolding bot_ivl_def by transfer (simp add: sup_rep_def eq_ivl_empty)

lemma inf_ivl_nice: [l1,h1] ⊓ [l2,h2] = [max l1 l2,min h1 h2]
by transfer (simp add: inf_rep_def)

lemma top_ivl_nice: ⊤ = [−∞,∞]
by (simp add: top_ivl_def)


instantiation ivl :: plus
begin

definition plus_rep :: eint2 ⇒ eint2 ⇒ eint2 where
plus_rep p1 p2 =
```

(*if is_empty_rep p1* ∨ *is_empty_rep p2 then empty_rep else*
let (*l1*,*h1*) = *p1*; (*l2*,*h2*) = *p2 in* (*l1*+*l2*, *h1*+*h2*))

lift_definition *plus_ivl* :: *ivl* ⇒ *ivl* ⇒ *ivl* **is** *plus_rep*
by(*auto simp*: *plus_rep_def eq_ivl_iff*)

instance ..
end

lemma *plus_ivl_nice*: [*l1*,*h1*] + [*l2*,*h2*] =
(*if* [*l1*,*h1*] = ⊥ ∨ [*l2*,*h2*] = ⊥ *then* ⊥ *else* [*l1*+*l2* , *h1*+*h2*])
unfolding *bot_ivl_def* **by** *transfer* (*auto simp*: *plus_rep_def eq_ivl_empty*)

lemma *uminus_eq_Minf*[*simp*]: −*x* = *Minf* ⟷ *x* = *Pinf*
by(*cases x*) *auto*
lemma *uminus_eq_Pinf*[*simp*]: −*x* = *Pinf* ⟷ *x* = *Minf*
by(*cases x*) *auto*

lemma *uminus_le_Fin_iff*: − *x* ≤ *Fin*(−*y*) ⟷ *Fin y* ≤ (*x*::′*a*::*ordered_ab_group_add extended*)
by(*cases x*) *auto*
lemma *Fin_uminus_le_iff*: *Fin*(−*y*) ≤ −*x* ⟷ *x* ≤ ((*Fin y*)::′*a*::*ordered_ab_group_add extended*)
by(*cases x*) *auto*

instantiation *ivl* :: *uminus*
begin

definition *uminus_rep* :: *eint2* ⇒ *eint2* **where**
uminus_rep p = (*let* (*l*,*h*) = *p in* (−*h*, −*l*))

lemma *γ_uminus_rep*: *i* ∈ *γ_rep p* ⟹ −*i* ∈ *γ_rep*(*uminus_rep p*)
by(*auto simp*: *uminus_rep_def γ_rep_def image_def uminus_le_Fin_iff Fin_uminus_le_iff*
split: *prod.split*)

lift_definition *uminus_ivl* :: *ivl* ⇒ *ivl* **is** *uminus_rep*
by (*auto simp*: *uminus_rep_def eq_ivl_def γ_rep_cases*)
(*auto simp*: *Icc_eq_Icc split*: *extended.splits*)

instance ..
end

lemma *γ_uminus*: *i* ∈ *γ_ivl iv* ⟹ −*i* ∈ *γ_ivl*(− *iv*)

by *transfer* (*rule* γ_*uminus_rep*)

lemma *uminus_nice*: $-[l,h] = [-h,-l]$
by *transfer* (*simp add*: *uminus_rep_def*)

instantiation *ivl* :: *minus*
begin

definition *minus_ivl* :: *ivl* ⇒ *ivl* ⇒ *ivl* **where**
(*iv1*::*ivl*) − *iv2* = *iv1* + −*iv2*

instance ..
end

definition *inv_plus_ivl* :: *ivl* ⇒ *ivl* ⇒ *ivl* ⇒ *ivl*∗*ivl* **where**
inv_plus_ivl iv iv1 iv2 = (*iv1* ⊓ (*iv* − *iv2*), *iv2* ⊓ (*iv* − *iv1*))

definition *above_rep* :: *eint2* ⇒ *eint2* **where**
above_rep p = (*if is_empty_rep p then empty_rep else let* (*l*,*h*) = *p in* (*l*,∞))

definition *below_rep* :: *eint2* ⇒ *eint2* **where**
below_rep p = (*if is_empty_rep p then empty_rep else let* (*l*,*h*) = *p in* (−∞,*h*))

lift_definition *above* :: *ivl* ⇒ *ivl* **is** *above_rep*
by(*auto simp*: *above_rep_def eq_ivl_iff*)

lift_definition *below* :: *ivl* ⇒ *ivl* **is** *below_rep*
by(*auto simp*: *below_rep_def eq_ivl_iff*)

lemma γ_*aboveI*: *i* ∈ γ_*ivl iv* ⟹ *i* ≤ *j* ⟹ *j* ∈ γ_*ivl*(*above iv*)
by *transfer*
  (*auto simp add*: *above_rep_def* γ_*rep_cases is_empty_rep_def*
    split: *extended.splits*)

lemma γ_*belowI*: *i* ∈ γ_*ivl iv* ⟹ *j* ≤ *i* ⟹ *j* ∈ γ_*ivl*(*below iv*)
by *transfer*
  (*auto simp add*: *below_rep_def* γ_*rep_cases is_empty_rep_def*
    split: *extended.splits*)

definition *inv_less_ivl* :: *bool* ⇒ *ivl* ⇒ *ivl* ⇒ *ivl* ∗ *ivl* **where**
inv_less_ivl res iv1 iv2 =

```
(if res
  then (iv1 ⊓ (below iv2 − [1,1]),
        iv2 ⊓ (above iv1 + [1,1]))
  else (iv1 ⊓ above iv2, iv2 ⊓ below iv1))

lemma above_nice: above[l,h] = (if [l,h] = ⊥ then ⊥ else [l,∞])
unfolding bot_ivl_def by transfer (simp add: above_rep_def eq_ivl_empty)

lemma below_nice: below[l,h] = (if [l,h] = ⊥ then ⊥ else [−∞,h])
unfolding bot_ivl_def by transfer (simp add: below_rep_def eq_ivl_empty)

lemma add_mono_le_Fin:
  ⟦x1 ≤ Fin y1; x2 ≤ Fin y2⟧ ⟹ x1 + x2 ≤ Fin (y1 + (y2::'a::ordered_ab_group_add))
by(drule (1) add_mono) simp

lemma add_mono_Fin_le:
  ⟦Fin y1 ≤ x1; Fin y2 ≤ x2⟧ ⟹ Fin(y1 + y2::'a::ordered_ab_group_add)
≤ x1 + x2
by(drule (1) add_mono) simp

global_interpretation Val_semilattice
where γ = γ_ivl and num' = num_ivl and plus' = (+)
proof (standard, goal_cases)
  case 1 thus ?case by transfer (simp add: le_iff_subset)
next
  case 2 show ?case by transfer (simp add: γ_rep_def)
next
  case 3 show ?case by transfer (simp add: γ_rep_def)
next
  case 4 thus ?case
    apply transfer
   apply(auto simp: γ_rep_def plus_rep_def add_mono_le_Fin add_mono_Fin_le)
    by(auto simp: empty_rep_def is_empty_rep_def)
qed


global_interpretation Val_lattice_gamma
where γ = γ_ivl and num' = num_ivl and plus' = (+)
defines aval_ivl = aval'
proof (standard, goal_cases)
  case 1 show ?case by(simp add: γ_inf)
next
  case 2 show ?case unfolding bot_ivl_def by transfer simp
qed
```

global_interpretation *Val_inv*
where γ = γ*_ivl* **and** *num'* = *num_ivl* **and** *plus'* = (+)
and *test_num'* = *in_ivl*
and *inv_plus'* = *inv_plus_ivl* **and** *inv_less'* = *inv_less_ivl*
proof (*standard*, *goal_cases*)
  case *1* **thus** *?case* **by** *transfer* (*auto simp*: γ*_rep_def*)
next
  case (*2* _ _ _ _ _ _ *i1 i2*) **thus** *?case*
    unfolding *inv_plus_ivl_def minus_ivl_def*
    apply(*clarsimp simp add*: γ*_inf*)
    using *gamma_plus'*[*of i1*+*i2* _ −*i1*] *gamma_plus'*[*of i1*+*i2* _ −*i2*]
    by(*simp add*: γ*_uminus*)
next
  case (*3 i1 i2*) **thus** *?case*
    unfolding *inv_less_ivl_def minus_ivl_def one_extended_def*
    apply(*clarsimp simp add*: γ*_inf split*: *if_splits*)
    using *gamma_plus'*[*of i1*+*1* _ −*1*] *gamma_plus'*[*of i2* − *1* _ *1*]
    apply(*simp add*: γ*_belowI*[*of i2*] γ*_aboveI*[*of i1*]
      uminus_ivl.abs_eq uminus_rep_def γ*_ivl_nice*)
    apply(*simp add*: γ*_aboveI*[*of i2*] γ*_belowI*[*of i1*])
    done
qed

global_interpretation *Abs_Int_inv*
where γ = γ*_ivl* **and** *num'* = *num_ivl* **and** *plus'* = (+)
and *test_num'* = *in_ivl*
and *inv_plus'* = *inv_plus_ivl* **and** *inv_less'* = *inv_less_ivl*
defines *inv_aval_ivl* = *inv_aval'*
and *inv_bval_ivl* = *inv_bval'*
and *step_ivl* = *step'*
and *AI_ivl* = *AI*
and *aval_ivl'* = *aval''*
..

Monotonicity:

lemma *mono_plus_ivl*: *iv1* ≤ *iv2* ⟹ *iv3* ≤ *iv4* ⟹ *iv1*+*iv3* ≤ *iv2*+(*iv4*::*ivl*)
apply *transfer*
apply(*auto simp*: *plus_rep_def le_iff_subset split*: *if_splits*)
by(*auto simp*: *is_empty_rep_iff* γ*_rep_cases split*: *extended.splits*)

lemma *mono_minus_ivl*: *iv1* ≤ *iv2* ⟹ −*iv1* ≤ −(*iv2*::*ivl*)
apply *transfer*
apply(*auto simp*: *uminus_rep_def le_iff_subset split*: *if_splits prod.split*)

by(*auto simp*: *γ_rep_cases split*: *extended.splits*)

lemma *mono_above*: *iv1* ≤ *iv2* ⟹ *above iv1* ≤ *above iv2*
apply *transfer*
apply(*auto simp*: *above_rep_def le_iff_subset split*: *if_splits prod.split*)
by(*auto simp*: *is_empty_rep_iff γ_rep_cases split*: *extended.splits*)

lemma *mono_below*: *iv1* ≤ *iv2* ⟹ *below iv1* ≤ *below iv2*
apply *transfer*
apply(*auto simp*: *below_rep_def le_iff_subset split*: *if_splits prod.split*)
by(*auto simp*: *is_empty_rep_iff γ_rep_cases split*: *extended.splits*)

global_interpretation *Abs_Int_inv_mono*
where *γ* = *γ_ivl* **and** *num′* = *num_ivl* **and** *plus′* = (+)
and *test_num′* = *in_ivl*
and *inv_plus′* = *inv_plus_ivl* **and** *inv_less′* = *inv_less_ivl*
proof (*standard*, *goal_cases*)
  case *1* **thus** *?case* **by** (*rule mono_plus_ivl*)
next
  case *2* **thus** *?case*
    unfolding *inv_plus_ivl_def minus_ivl_def less_eq_prod_def*
    by (*auto simp*: *le_infI1 le_infI2 mono_plus_ivl mono_minus_ivl*)
next
  case *3* **thus** *?case*
    unfolding *less_eq_prod_def inv_less_ivl_def minus_ivl_def*
   by (*auto simp*: *le_infI1 le_infI2 mono_plus_ivl mono_above mono_below*)
qed

14.13.1 Tests

value *show_acom_opt* (*AI_ivl test1_ivl*)

Better than *AI_const*:

value *show_acom_opt* (*AI_ivl test3_const*)
value *show_acom_opt* (*AI_ivl test4_const*)
value *show_acom_opt* (*AI_ivl test6_const*)

definition *steps c i* = (*step_ivl* ⊤ ^^ *i*) (*bot c*)

value *show_acom_opt* (*AI_ivl test2_ivl*)
value *show_acom* (*steps test2_ivl 0*)
value *show_acom* (*steps test2_ivl 1*)
value *show_acom* (*steps test2_ivl 2*)
value *show_acom* (*steps test2_ivl 3*)

Fixed point reached in 2 steps. Not so if the start value of x is known:

value *show_acom_opt* (*AI_ivl test3_ivl*)
value *show_acom* (*steps test3_ivl 0*)
value *show_acom* (*steps test3_ivl 1*)
value *show_acom* (*steps test3_ivl 2*)
value *show_acom* (*steps test3_ivl 3*)
value *show_acom* (*steps test3_ivl 4*)
value *show_acom* (*steps test3_ivl 5*)

Takes as many iterations as the actual execution. Would diverge if loop did not terminate. Worse still, as the following example shows: even if the actual execution terminates, the analysis may not. The value of y keeps increasing as the analysis is iterated, no matter how long:

value *show_acom* (*steps test4_ivl 50*)

Relationships between variables are NOT captured:

value *show_acom_opt* (*AI_ivl test5_ivl*)

Again, the analysis would not terminate:

value *show_acom* (*steps test6_ivl 50*)

end

14.14 Widening and Narrowing

theory *Abs_Int3*
imports *Abs_Int2_ivl*
begin

class *widen* =
fixes *widen* :: $'a \Rightarrow 'a \Rightarrow 'a$ (**infix** ∇ *65*)

class *narrow* =
fixes *narrow* :: $'a \Rightarrow 'a \Rightarrow 'a$ (**infix** $\triangle$ *65*)

class *wn* = *widen* + *narrow* + *order* +
assumes *widen1*: $x \le x \nabla y$
assumes *widen2*: $y \le x \nabla y$
assumes *narrow1*: $y \le x \Longrightarrow y \le x \triangle y$
assumes *narrow2*: $y \le x \Longrightarrow x \triangle y \le x$
begin

lemma *narrowid*[*simp*]: $x \triangle x = x$
by (*rule order.antisym*) (*simp_all add*: *narrow1 narrow2*)

end

lemma *top_widen_top*[*simp*]: ⊤ ∇ ⊤ = (⊤::_::{*wn*,*order_top*})
by (*metis eq_iff top_greatest widen2*)

instantiation *ivl* :: *wn*
begin

definition *widen_rep p1 p2* =
  (*if is_empty_rep p1 then p2 else if is_empty_rep p2 then p1 else*
   let (*l1*,*h1*) = *p1*; (*l2*,*h2*) = *p2*
   in (*if l2* < *l1 then Minf else l1*, *if h1* < *h2 then Pinf else h1*))

lift_definition *widen_ivl* :: *ivl* ⇒ *ivl* ⇒ *ivl* **is** *widen_rep*
by(*auto simp*: *widen_rep_def eq_ivl_iff*)

definition *narrow_rep p1 p2* =
  (*if is_empty_rep p1* ∨ *is_empty_rep p2 then empty_rep else*
   let (*l1*,*h1*) = *p1*; (*l2*,*h2*) = *p2*
   in (*if l1* = *Minf then l2 else l1*, *if h1* = *Pinf then h2 else h1*))

lift_definition *narrow_ivl* :: *ivl* ⇒ *ivl* ⇒ *ivl* **is** *narrow_rep*
by(*auto simp*: *narrow_rep_def eq_ivl_iff*)

instance
proof
qed (*transfer*, *auto simp*: *widen_rep_def narrow_rep_def le_iff_subset γ_rep_def subset_eq is_empty_rep_def empty_rep_def eq_ivl_def split*: *if_splits extended.splits*)+

end

instantiation *st* :: ({*order_top*,*wn*})*wn*
begin

lift_definition *widen_st* :: ′*a st* ⇒ ′*a st* ⇒ ′*a st* **is** *map2_st_rep* (∇)
by(*auto simp*: *eq_st_def*)

lift_definition *narrow_st* :: ′*a st* ⇒ ′*a st* ⇒ ′*a st* **is** *map2_st_rep* (△)
by(*auto simp*: *eq_st_def*)

instance
proof (*standard*, *goal_cases*)

  case *1* **thus** *?case* **by** *transfer* (*simp add*: *less_eq_st_rep_iff widen1*)
next
  case *2* **thus** *?case* **by** *transfer* (*simp add*: *less_eq_st_rep_iff widen2*)
next
  case *3* **thus** *?case* **by** *transfer* (*simp add*: *less_eq_st_rep_iff narrow1*)
next
  case *4* **thus** *?case* **by** *transfer* (*simp add*: *less_eq_st_rep_iff narrow2*)
qed

end

instantiation *option* :: (*wn*)*wn*
begin

fun *widen_option* **where**
None ∇ *x* = *x* |
x ∇ *None* = *x* |
(*Some x*) ∇ (*Some y*) = *Some*(*x* ∇ *y*)

fun *narrow_option* **where**
None △ *x* = *None* |
x △ *None* = *None* |
(*Some x*) △ (*Some y*) = *Some*(*x* △ *y*)

instance
proof (*standard*, *goal_cases*)
  case (*1 x y*) **thus** *?case*
    by(*induct x y rule*: *widen_option.induct*)(*simp_all add*: *widen1*)
next
  case (*2 x y*) **thus** *?case*
    by(*induct x y rule*: *widen_option.induct*)(*simp_all add*: *widen2*)
next
  case (*3 x y*) **thus** *?case*
    by(*induct x y rule*: *narrow_option.induct*) (*simp_all add*: *narrow1*)
next
  case (*4 y x*) **thus** *?case*
    by(*induct x y rule*: *narrow_option.induct*) (*simp_all add*: *narrow2*)
qed

end

definition *map2_acom* :: (′*a* ⇒ ′*a* ⇒ ′*a*) ⇒ ′*a acom* ⇒ ′*a acom* ⇒ ′*a acom*
where

map2_acom f C1 C2 = annotate (λp. f (anno C1 p) (anno C2 p)) (strip C1)

instantiation *acom* :: (*widen*)*widen*
begin
definition *widen_acom = map2_acom* (∇)
instance ..
end

instantiation *acom* :: (*narrow*)*narrow*
begin
definition *narrow_acom = map2_acom* (△)
instance ..
end

lemma *strip_map2_acom*[*simp*]:
strip C1 = strip C2 ⟹ *strip*(*map2_acom f C1 C2*) = *strip C1*
by(*simp add*: *map2_acom_def*)

lemma *strip_widen_acom*[*simp*]:
strip C1 = strip C2 ⟹ *strip*(*C1* ∇ *C2*) = *strip C1*
by(*simp add*: *widen_acom_def*)

lemma *strip_narrow_acom*[*simp*]:
strip C1 = strip C2 ⟹ *strip*(*C1* △ *C2*) = *strip C1*
by(*simp add*: *narrow_acom_def*)

lemma *narrow1_acom*: *C2* ≤ *C1* ⟹ *C2* ≤ *C1* △ (*C2*::′*a*::*wn acom*)
by(*simp add*: *narrow_acom_def narrow1 map2_acom_def less_eq_acom_def size_annos*)

lemma *narrow2_acom*: *C2* ≤ *C1* ⟹ *C1* △ (*C2*::′*a*::*wn acom*) ≤ *C1*
by(*simp add*: *narrow_acom_def narrow2 map2_acom_def less_eq_acom_def size_annos*)

14.14.1 Pre-fixpoint computation

definition *iter_widen* :: (′*a* ⇒ ′*a*) ⇒ ′*a* ⇒ (′*a*::{*order*,*widen*})*option*
where *iter_widen f = while_option* (λ*x*. ¬ *f x* ≤ *x*) (λ*x*. *x* ∇ *f x*)

definition *iter_narrow* :: (′*a* ⇒ ′*a*) ⇒ ′*a* ⇒ (′*a*::{*order*,*narrow*})*option*
where *iter_narrow f = while_option* (λ*x*. *x* △ *f x* < *x*) (λ*x*. *x* △ *f x*)

definition *pfp_wn* :: ('*a*::{*order*,*widen*,*narrow*} ⇒ '*a*) ⇒ '*a* ⇒ '*a option*
where *pfp_wn f x* =
  (*case iter_widen f x of None* ⇒ *None* | *Some p* ⇒ *iter_narrow f p*)

lemma *iter_widen_pfp*: *iter_widen f x* = *Some p* ⟹ *f p* ≤ *p*
by(*auto simp add*: *iter_widen_def dest*: *while_option_stop*)

lemma *iter_widen_inv*:
assumes !!*x*. *P x* ⟹ *P*(*f x*) !!*x1 x2*. *P x1* ⟹ *P x2* ⟹ *P*(*x1* ∇ *x2*) **and**
P x
and *iter_widen f x* = *Some y* **shows** *P y*
using *while_option_rule*[**where** *P* = *P*, *OF _ assms*(*4*)[*unfolded iter_widen_def*]]
by (*blast intro*: *assms*(*1*−*3*))

lemma *strip_while*: **fixes** *f* :: '*a acom* ⇒ '*a acom*
assumes ∀ *C*. *strip* (*f C*) = *strip C* **and** *while_option P f C* = *Some C'*
shows *strip C'* = *strip C*
using *while_option_rule*[**where** *P* = λ*C'*. *strip C'* = *strip C*, *OF _ assms*(*2*)]
by (*metis assms*(*1*))

lemma *strip_iter_widen*: **fixes** *f* :: '*a*::{*order*,*widen*} *acom* ⇒ '*a acom*
assumes ∀ *C*. *strip* (*f C*) = *strip C* **and** *iter_widen f C* = *Some C'*
shows *strip C'* = *strip C*
proof−
  have ∀ *C*. *strip*(*C* ∇ *f C*) = *strip C*
    by (*metis assms*(*1*) *strip_map2_acom widen_acom_def*)
  from *strip_while*[*OF this*] *assms*(*2*) **show** *?thesis* **by**(*simp add*: *iter_widen_def*)
qed

lemma *iter_narrow_pfp*:
assumes *mono*: !!*x1 x2*::_::*wn acom*. *P x1* ⟹ *P x2* ⟹ *x1* ≤ *x2* ⟹ *f*
x1 ≤ *f x2*
and *Pinv*: !!*x*. *P x* ⟹ *P*(*f x*) !!*x1 x2*. *P x1* ⟹ *P x2* ⟹ *P*(*x1* △ *x2*)
and *P p0* **and** *f p0* ≤ *p0* **and** *iter_narrow f p0* = *Some p*
shows *P p* ∧ *f p* ≤ *p*
proof−
  let *?Q* = %*p*. *P p* ∧ *f p* ≤ *p* ∧ *p* ≤ *p0*
  have *?Q* (*p* △ *f p*) **if** *Q*: *?Q p* **for** *p*
  proof *auto*
    note *P* = *conjunct1*[*OF Q*] **and** *12* = *conjunct2*[*OF Q*]
    note *1* = *conjunct1*[*OF 12*] **and** *2* = *conjunct2*[*OF 12*]
    let *?p'* = *p* △ *f p*

show *P ?p'* **by** (*blast intro*: *P Pinv*)
have *f ?p'* ≤ *f p* **by**(*rule mono*[*OF* ‹*P* (*p* △ *f p*)› *P narrow2_acom*[*OF* *1*]])
also have ... ≤ *?p'* **by**(*rule narrow1_acom*[*OF 1*])
finally show *f ?p'* ≤ *?p'* **.**
have *?p'* ≤ *p* **by** (*rule narrow2_acom*[*OF 1*])
also have *p* ≤ *p0* **by**(*rule 2*)
finally show *?p'* ≤ *p0* **.**
qed
thus *?thesis*
using *while_option_rule*[**where** *P* = *?Q*, *OF* _ *assms*(*6*)[*simplified* *iter_narrow_def*]]
by (*blast intro*: *assms*(*4*,*5*) *le_refl*)
qed

lemma *pfp_wn_pfp*:
assumes *mono*: !!*x1 x2*::_::*wn acom*. *P x1* ⟹ *P x2* ⟹ *x1* ≤ *x2* ⟹ *f* *x1* ≤ *f x2*
and *Pinv*: *P x* !!*x*. *P x* ⟹ *P*(*f x*)
!!*x1 x2*. *P x1* ⟹ *P x2* ⟹ *P*(*x1* ∇ *x2*)
!!*x1 x2*. *P x1* ⟹ *P x2* ⟹ *P*(*x1* △ *x2*)
and *pfp_wn*: *pfp_wn f x* = *Some p* **shows** *P p* ∧ *f p* ≤ *p*
proof−
from *pfp_wn* **obtain** *p0*
where *its*: *iter_widen f x* = *Some p0* *iter_narrow f p0* = *Some p*
by(*auto simp*: *pfp_wn_def split*: *option.splits*)
have *P p0* **by** (*blast intro*: *iter_widen_inv*[**where** *P*=*P*] *its*(*1*) *Pinv*(*1*−*3*))
thus *?thesis*
by − (*assumption* |
rule iter_narrow_pfp[**where** *P*=*P*] *mono Pinv*(*2*,*4*) *iter_widen_pfp* *its*)+
qed

lemma *strip_pfp_wn*:
⟦ ∀ *C*. *strip*(*f C*) = *strip C*; *pfp_wn f C* = *Some C'* ⟧ ⟹ *strip C'* = *strip* *C*
by(*auto simp add*: *pfp_wn_def iter_narrow_def split*: *option.splits*)
(*metis* (*mono_tags*) *strip_iter_widen strip_narrow_acom strip_while*)

locale *Abs_Int_wn* = *Abs_Int_inv_mono* **where** γ=γ
for γ :: *'av*::{*wn*,*bounded_lattice*} ⇒ *val set*
begin

definition *AI_wn* :: *com* ⇒ *'av st option acom option* **where**
AI_wn c = *pfp_wn* (*step'* ⊤) (*bot c*)

lemma *AI_wn_correct*: *AI_wn c* = *Some C* ⟹ *CS c* ≤ γ_c *C*
proof(*simp add*: *CS_def AI_wn_def*)
  assume *1*: *pfp_wn* (*step'* ⊤) (*bot c*) = *Some C*
  have *2*: *strip C* = *c* ∧ *step'* ⊤ *C* ≤ *C*
   by(*rule pfp_wn_pfp*[**where** *x*=*bot c*]) (*simp_all add*: *1 mono_step'_top*)
  have *pfp*: *step* (γ_o ⊤) (γ_c *C*) ≤ γ_c *C*
  proof(*rule order_trans*)
    show *step* (γ_o ⊤) (γ_c *C*) ≤ γ_c (*step'* ⊤ *C*)
      by(*rule step_step'*)
    show ... ≤ γ_c *C*
      by(*rule mono_gamma_c*[*OF conjunct2*[*OF 2*]])
  qed
  have *3*: *strip* (γ_c *C*) = *c* **by**(*simp add*: *strip_pfp_wn*[*OF _ 1*])
  have *lfp c* (*step* (γ_o ⊤)) ≤ γ_c *C*
    by(*rule lfp_lowerbound*[*simplified*,**where** *f*=*step* (γ_o ⊤), *OF 3 pfp*])
  thus *lfp c* (*step UNIV*) ≤ γ_c *C* **by** *simp*
qed

end

global_interpretation *Abs_Int_wn*
where γ = *γ_ivl* **and** *num'* = *num_ivl* **and** *plus'* = (+)
and *test_num'* = *in_ivl*
and *inv_plus'* = *inv_plus_ivl* **and** *inv_less'* = *inv_less_ivl*
defines *AI_wn_ivl* = *AI_wn*
..

14.14.2 Tests

definition *step_up_ivl n* = ((λ*C*. *C* ∇ *step_ivl* ⊤ *C*)^^*n*)
definition *step_down_ivl n* = ((λ*C*. *C* △ *step_ivl* ⊤ *C*)^^*n*)

For *test3_ivl*, *AI_ivl* needed as many iterations as the loop took to execute. In contrast, *AI_wn_ivl* converges in a constant number of steps:

value *show_acom* (*step_up_ivl 1* (*bot test3_ivl*))
value *show_acom* (*step_up_ivl 2* (*bot test3_ivl*))
value *show_acom* (*step_up_ivl 3* (*bot test3_ivl*))
value *show_acom* (*step_up_ivl 4* (*bot test3_ivl*))
value *show_acom* (*step_up_ivl 5* (*bot test3_ivl*))
value *show_acom* (*step_up_ivl 6* (*bot test3_ivl*))
value *show_acom* (*step_up_ivl 7* (*bot test3_ivl*))

value *show_acom* (*step_up_ivl 8* (*bot test3_ivl*))
value *show_acom* (*step_down_ivl 1* (*step_up_ivl 8* (*bot test3_ivl*)))
value *show_acom* (*step_down_ivl 2* (*step_up_ivl 8* (*bot test3_ivl*)))
value *show_acom* (*step_down_ivl 3* (*step_up_ivl 8* (*bot test3_ivl*)))
value *show_acom* (*step_down_ivl 4* (*step_up_ivl 8* (*bot test3_ivl*)))
value *show_acom_opt* (*AI_wn_ivl test3_ivl*)

Now all the analyses terminate:

value *show_acom_opt* (*AI_wn_ivl test4_ivl*)
value *show_acom_opt* (*AI_wn_ivl test5_ivl*)
value *show_acom_opt* (*AI_wn_ivl test6_ivl*)

14.14.3 Generic Termination Proof

lemma *top_on_opt_widen*:
top_on_opt o1 X ⟹ *top_on_opt o2 X* ⟹ *top_on_opt* (*o1* ∇ *o2* :: _ *st option*) *X*
apply(*induct o1 o2 rule*: *widen_option.induct*)
apply (*auto*)
by *transfer simp*

lemma *top_on_opt_narrow*:
top_on_opt o1 X ⟹ *top_on_opt o2 X* ⟹ *top_on_opt* (*o1* △ *o2* :: _ *st option*) *X*
apply(*induct o1 o2 rule*: *narrow_option.induct*)
apply (*auto*)
by *transfer simp*

lemma *annos_map2_acom*[*simp*]: *strip C2* = *strip C1* ⟹
annos(*map2_acom f C1 C2*) = *map* (%(*x*,*y*).*f x y*) (*zip* (*annos C1*) (*annos C2*))
by(*simp add*: *map2_acom_def list_eq_iff_nth_eq size_annos anno_def*[*symmetric*] *size_annos_same*[*of C1 C2*])

lemma *top_on_acom_widen*:
⟦*top_on_acom C1 X*; *strip C1* = *strip C2*; *top_on_acom C2 X*⟧
⟹ *top_on_acom* (*C1* ∇ *C2* :: _ *st option acom*) *X*
by(*auto simp add*: *widen_acom_def top_on_acom_def*)(*metis top_on_opt_widen in_set_zipE*)

lemma *top_on_acom_narrow*:
⟦*top_on_acom C1 X*; *strip C1* = *strip C2*; *top_on_acom C2 X*⟧
⟹ *top_on_acom* (*C1* △ *C2* :: _ *st option acom*) *X*

by(*auto simp add*: *narrow_acom_def top_on_acom_def*)(*metis top_on_opt_narrow in_set_zipE*)

The assumptions for widening and narrowing differ because during narrowing we have the invariant $y \leq x$ (where y is the next iterate), but during widening there is no such invariant, there we only have that not yet $y \leq x$. This complicates the termination proof for widening.

locale *Measure_wn* = *Measure1* **where** *m=m*
  for *m* :: *'av*::{*order_top*,*wn*} ⇒ *nat* +
fixes *n* :: *'av* ⇒ *nat*
assumes *m_anti_mono*: $x \leq y \Longrightarrow m\ x \geq m\ y$
assumes *m_widen*: $\sim\ y \leq x \Longrightarrow m(x \nabla y) < m\ x$
assumes *n_narrow*: $y \leq x \Longrightarrow x \triangle y < x \Longrightarrow n(x \triangle y) < n\ x$

begin

lemma *m_s_anti_mono_rep*: **assumes** $\forall x.\ S1\ x \leq S2\ x$
shows $(\sum x{\in}X.\ m\ (S2\ x)) \leq (\sum x{\in}X.\ m\ (S1\ x))$
proof−
  from *assms* **have** $\forall x.\ m(S1\ x) \geq m(S2\ x)$ **by** (*metis m_anti_mono*)
  thus $(\sum x{\in}X.\ m\ (S2\ x)) \leq (\sum x{\in}X.\ m\ (S1\ x))$ **by** (*metis sum_mono*)
qed

lemma *m_s_anti_mono*: $S1 \leq S2 \Longrightarrow m\_s\ S1\ X \geq m\_s\ S2\ X$
unfolding *m_s_def*
apply (*transfer fixing*: *m*)
apply(*simp add*: *less_eq_st_rep_iff eq_st_def m_s_anti_mono_rep*)
done

lemma *m_s_widen_rep*: **assumes** *finite X S1* = *S2 on* −*X* ¬ *S2 x* ≤ *S1 x*
  shows $(\sum x{\in}X.\ m\ (S1\ x \nabla S2\ x)) < (\sum x{\in}X.\ m\ (S1\ x))$
proof−
  have *1*: $\forall x{\in}X.\ m(S1\ x) \geq m(S1\ x \nabla S2\ x)$
    by (*metis m_anti_mono wn_class.widen1*)
  have $x \in X$ **using** *assms*(*2*,*3*)
    by(*auto simp add*: *Ball_def*)
  hence *2*: $\exists x{\in}X.\ m(S1\ x) > m(S1\ x \nabla S2\ x)$
    using *assms*(*3*) *m_widen* **by** *blast*
  from *sum_strict_mono_ex1*[*OF* ‹*finite X*› *1 2*]
  show *?thesis* **.**
qed

lemma *m_s_widen*: *finite X* ⟹ *fun S1* = *fun S2 on* −*X* ==>

~ *S2* ≤ *S1* ⟹ *m_s* (*S1* ∇ *S2*) *X* < *m_s S1 X*
apply(*auto simp add*: *less_st_def m_s_def*)
apply (*transfer fixing*: *m*)
apply(*auto simp add*: *less_eq_st_rep_iff m_s_widen_rep*)
done

lemma *m_o_anti_mono*: *finite X* ⟹ *top_on_opt o1* (−*X*) ⟹ *top_on_opt o2* (−*X*) ⟹
o1 ≤ *o2* ⟹ *m_o o1 X* ≥ *m_o o2 X*
proof(*induction o1 o2 rule*: *less_eq_option.induct*)
case *1* **thus** *?case* **by** (*simp add*: *m_o_def*)(*metis m_s_anti_mono*)
next
case *2* **thus** *?case*
by(*simp add*: *m_o_def le_SucI m_s_h split*: *option.splits*)
next
case *3* **thus** *?case* **by** *simp*
qed

lemma *m_o_widen*: ⟦ *finite X*; *top_on_opt S1* (−*X*); *top_on_opt S2* (−*X*); ¬ *S2* ≤ *S1* ⟧ ⟹
m_o (*S1* ∇ *S2*) *X* < *m_o S1 X*
by(*auto simp*: *m_o_def m_s_h less_Suc_eq_le m_s_widen split*: *option.split*)

lemma *m_c_widen*:
strip C1 = *strip C2* ⟹ *top_on_acom C1* (−*vars C1*) ⟹ *top_on_acom C2* (−*vars C2*)
⟹ ¬ *C2* ≤ *C1* ⟹ *m_c* (*C1* ∇ *C2*) < *m_c C1*
apply(*auto simp*: *m_c_def widen_acom_def map2_acom_def size_annos*[*symmetric*] *anno_def*[*symmetric*]*sum_list_sum_nth*)
apply(*subgoal_tac length*(*annos C2*) = *length*(*annos C1*))
prefer *2* **apply** (*simp add*: *size_annos_same2*)
apply (*auto*)
apply(*rule sum_strict_mono_ex1*)
apply(*auto simp add*: *m_o_anti_mono vars_acom_def anno_def top_on_acom_def top_on_opt_widen widen1 less_eq_acom_def listrel_iff_nth*)
apply(*rule_tac x*=*p* **in** *bexI*)
apply (*auto simp*: *vars_acom_def m_o_widen top_on_acom_def*)
done

definition *n_s* :: *'av st* ⇒ *vname set* ⇒ *nat* (n_s) **where**
n_s *S X* = (∑ *x*∈*X*. *n*(*fun S x*))

lemma *n_s_narrow_rep*:

assumes *finite X* $S1 = S2$ *on* $-X$ $\forall x.\ S2\ x \le S1\ x$ $\forall x.\ S1\ x \triangle S2\ x \le S1\ x$
$S1\ x \ne S1\ x \triangle S2\ x$
shows $(\sum x{\in}X.\ n\ (S1\ x \triangle S2\ x)) < (\sum x{\in}X.\ n\ (S1\ x))$
proof−
have *1*: $\forall x.\ n(S1\ x \triangle S2\ x) \le n(S1\ x)$
by (*metis assms*(*3*) *assms*(*4*) *eq_iff less_le_not_le n_narrow*)
have $x \in X$ **by** (*metis Compl_iff assms*(*2*) *assms*(*5*) *narrowid*)
hence *2*: $\exists x{\in}X.\ n(S1\ x \triangle S2\ x) < n(S1\ x)$
by (*metis assms*(*3*−*5*) *eq_iff less_le_not_le n_narrow*)
show *?thesis*
apply(*rule sum_strict_mono_ex1*[*OF* ‹*finite X*›]) **using** *1 2* **by** *blast*+
qed

lemma *n_s_narrow*: *finite X* $\Longrightarrow$ *fun S1* = *fun S2 on* $-X \Longrightarrow S2 \le S1$
$\Longrightarrow S1 \triangle S2 < S1$
$\Longrightarrow n_s\ (S1 \triangle S2)\ X < n_s\ S1\ X$
apply(*auto simp add*: *less_st_def n_s_def*)
apply (*transfer fixing*: *n*)
apply(*auto simp add*: *less_eq_st_rep_iff eq_st_def fun_eq_iff n_s_narrow_rep*)
done

definition *n_o* :: ′*av st option* ⇒ *vname set* ⇒ *nat* (n_o) **where**
n_o *opt X* = (*case opt of None* ⇒ *0* | *Some S* ⇒ n_s *S X* + *1*)

lemma *n_o_narrow*:
top_on_opt S1 ($-X$) $\Longrightarrow$ *top_on_opt S2* ($-X$) $\Longrightarrow$ *finite X*
$\Longrightarrow S2 \le S1 \Longrightarrow S1 \triangle S2 < S1 \Longrightarrow n_o\ (S1 \triangle S2)\ X < n_o\ S1\ X$
apply(*induction S1 S2 rule*: *narrow_option.induct*)
apply(*auto simp*: *n_o_def n_s_narrow*)
done

definition *n_c* :: ′*av st option acom* ⇒ *nat* (n_c) **where**
n_c *C* = *sum_list* (*map* ($\lambda a.\ n_o\ a$ (*vars C*)) (*annos C*))

lemma *less_annos_iff*: $(C1 < C2) = (C1 \le C2\ \wedge$
($\exists i{<}$*length* (*annos C1*). *annos C1* ! *i* < *annos C2* ! *i*))
by(*metis* (*opaque_lifting*, *no_types*) *less_le_not_le le_iff_le_annos size_annos_same2*)

lemma *n_c_narrow*: *strip C1* = *strip C2*
$\Longrightarrow$ *top_on_acom C1* (− *vars C1*) $\Longrightarrow$ *top_on_acom C2* (− *vars C2*)
$\Longrightarrow C2 \le C1 \Longrightarrow C1 \triangle C2 < C1 \Longrightarrow n_c\ (C1 \triangle C2) < n_c\ C1$
apply(*auto simp*: *n_c_def narrow_acom_def sum_list_sum_nth*)

apply(*subgoal_tac length*(*annos C2*) = *length*(*annos C1*))
prefer *2* **apply** (*simp add*: *size_annos_same2*)
apply (*auto*)
apply(*simp add*: *less_annos_iff le_iff_le_annos*)
apply(*rule sum_strict_mono_ex1*)
apply (*auto simp*: *vars_acom_def top_on_acom_def*)
apply (*metis n_o_narrow nth_mem finite_cvars less_imp_le le_less order_refl*)
apply(*rule_tac x=i* **in** *bexI*)
prefer *2* **apply** *simp*
apply(*rule n_o_narrow*[**where** *X* = *vars*(*strip C2*)])
apply (*simp_all*)
done

end

lemma *iter_widen_termination*:
fixes *m* :: *'a::wn acom* ⇒ *nat*
assumes *P_f*: $\bigwedge C.\ P\ C \Longrightarrow P(f\ C)$
and *P_widen*: $\bigwedge C1\ C2.\ P\ C1 \Longrightarrow P\ C2 \Longrightarrow P(C1\ \nabla\ C2)$
and *m_widen*: $\bigwedge C1\ C2.\ P\ C1 \Longrightarrow P\ C2 \Longrightarrow \sim\ C2 \le C1 \Longrightarrow m(C1\ \nabla\ C2) < m\ C1$
and *P C* **shows** ∃ *C'*. *iter_widen f C* = *Some C'*
proof(*simp add*: *iter_widen_def*,
rule measure_while_option_Some[**where** *P* = *P* **and** *f*=*m*])
show *P C* **by**(*rule* ‹*P C*›)
next
fix *C* **assume** *P C* ¬ *f C* ≤ *C* **thus** $P\ (C\ \nabla\ f\ C) \wedge m\ (C\ \nabla\ f\ C) < m\ C$
by(*simp add*: *P_f P_widen m_widen*)
qed

lemma *iter_narrow_termination*:
fixes *n* :: *'a::wn acom* ⇒ *nat*
assumes *P_f*: $\bigwedge C.\ P\ C \Longrightarrow P(f\ C)$
and *P_narrow*: $\bigwedge C1\ C2.\ P\ C1 \Longrightarrow P\ C2 \Longrightarrow P(C1\ \triangle\ C2)$
and *mono*: $\bigwedge C1\ C2.\ P\ C1 \Longrightarrow P\ C2 \Longrightarrow C1 \le C2 \Longrightarrow f\ C1 \le f\ C2$
and *n_narrow*: $\bigwedge C1\ C2.\ P\ C1 \Longrightarrow P\ C2 \Longrightarrow C2 \le C1 \Longrightarrow C1\ \triangle\ C2 < C1 \Longrightarrow n(C1\ \triangle\ C2) < n\ C1$
and *init*: *P C f C* ≤ *C* **shows** ∃ *C'*. *iter_narrow f C* = *Some C'*
proof(*simp add*: *iter_narrow_def*,
rule measure_while_option_Some[**where** *f*=*n* **and** *P* = %*C*. $P\ C \wedge f\ C \le C$])

show *P C* ∧ *f C* ≤ *C* **using** *init* **by** *blast*
next
fix *C* **assume** *1*: *P C* ∧ *f C* ≤ *C* **and** *2*: *C* △ *f C* < *C*
hence *P* (*C* △ *f C*) **by**(*simp add*: *P_f P_narrow*)
moreover then have *f* (*C* △ *f C*) ≤ *C* △ *f C*
by (*metis narrow1_acom narrow2_acom 1 mono order_trans*)
moreover have *n* (*C* △ *f C*) < *n C* **using** *1 2* **by**(*simp add*: *n_narrow P_f*)
ultimately show (*P* (*C* △ *f C*) ∧ *f* (*C* △ *f C*) ≤ *C* △ *f C*) ∧ *n*(*C* △ *f C*) < *n C*
by *blast*
qed

locale *Abs_Int_wn_measure* = *Abs_Int_wn* **where** *γ*=*γ* + *Measure_wn* **where** *m*=*m*
for *γ* :: *'av*::{*wn*,*bounded_lattice*} ⇒ *val set* **and** *m* :: *'av* ⇒ *nat*

14.14.4 Termination: Intervals

definition *m_rep* :: *eint2* ⇒ *nat* **where**
m_rep p = (*if is_empty_rep p then 3 else*
let (*l*,*h*) = *p in* (*case l of Minf* ⇒ *0* | _ ⇒ *1*) + (*case h of Pinf* ⇒ *0* | _ ⇒ *1*))

lift_definition *m_ivl* :: *ivl* ⇒ *nat* **is** *m_rep*
by(*auto simp*: *m_rep_def eq_ivl_iff*)

lemma *m_ivl_nice*: *m_ivl*[*l*,*h*] = (*if* [*l*,*h*] = ⊥ *then 3 else*
(*if l* = *Minf then 0 else 1*) + (*if h* = *Pinf then 0 else 1*))
unfolding *bot_ivl_def*
by *transfer* (*auto simp*: *m_rep_def eq_ivl_empty split*: *extended.split*)

lemma *m_ivl_height*: *m_ivl iv* ≤ *3*
by *transfer* (*simp add*: *m_rep_def split*: *prod.split extended.split*)

lemma *m_ivl_anti_mono*: *y* ≤ *x* ⟹ *m_ivl x* ≤ *m_ivl y*
by *transfer*
(*auto simp*: *m_rep_def is_empty_rep_def γ_rep_cases le_iff_subset*
split: *prod.split extended.splits if_splits*)

lemma *m_ivl_widen*:
~ *y* ≤ *x* ⟹ *m_ivl*(*x* ∇ *y*) < *m_ivl x*
by *transfer*
(*auto simp*: *m_rep_def widen_rep_def is_empty_rep_def γ_rep_cases*

le_iff_subset
    split: *prod.split extended.splits if_splits*)

definition *n_ivl* :: *ivl* ⇒ *nat* **where**
n_ivl iv = *3* − *m_ivl iv*

lemma *n_ivl_narrow*:
  $x \triangle y < x \Longrightarrow n\_ivl(x \triangle y) < n\_ivl\ x$
unfolding *n_ivl_def*
apply(*subst* (*asm*) *less_le_not_le*)
apply *transfer*
by(*auto simp add*: *m_rep_def narrow_rep_def is_empty_rep_def empty_rep_def γ_rep_cases le_iff_subset*
    split: *prod.splits if_splits extended.split*)

global_interpretation *Abs_Int_wn_measure*
where γ = γ*_ivl* **and** *num'* = *num_ivl* **and** *plus'* = (+)
and *test_num'* = *in_ivl*
and *inv_plus'* = *inv_plus_ivl* **and** *inv_less'* = *inv_less_ivl*
and *m* = *m_ivl* **and** *n* = *n_ivl* **and** *h* = *3*
proof (*standard*, *goal_cases*)
  case *2* **thus** *?case* **by**(*rule m_ivl_anti_mono*)
next
  case *1* **thus** *?case* **by**(*rule m_ivl_height*)
next
  case *3* **thus** *?case* **by**(*rule m_ivl_widen*)
next
  case *4* **from** *4*(*2*) **show** *?case* **by**(*rule n_ivl_narrow*)
  — note that the first assms is unnecessary for intervals
qed

lemma *iter_winden_step_ivl_termination*:
  ∃ *C*. *iter_widen* (*step_ivl* ⊤) (*bot c*) = *Some C*
apply(*rule iter_widen_termination*[**where** *m* = *m_c* **and** *P* = %*C*. *strip C* = *c* ∧ *top_on_acom C* (− *vars C*)])
apply (*auto simp add*: *m_c_widen top_on_bot top_on_step'*[*simplified comp_def vars_acom_def*]
  vars_acom_def top_on_acom_widen)
done

lemma *iter_narrow_step_ivl_termination*:
  top_on_acom C (− *vars C*) ⟹ *step_ivl* ⊤ *C* ≤ *C* ⟹
  ∃ *C'*. *iter_narrow* (*step_ivl* ⊤) *C* = *Some C'*

apply(*rule iter_narrow_termination*[**where** $n = n\_c$ **and** $P = \%C'.$ *strip* $C =$ *strip* $C' \wedge$ *top_on_acom* C' $(-$*vars* $C')$])
apply(*auto simp*: *top_on_step'*[*simplified comp_def vars_acom_def*]
    mono_step'_top n_c_narrow vars_acom_def top_on_acom_narrow)
done

theorem *AI_wn_ivl_termination*:
  $\exists$ $C.$ *AI_wn_ivl* $c =$ *Some* C
apply(*auto simp*: *AI_wn_def pfp_wn_def iter_winden_step_ivl_termination*
    split: *option.split*)
apply(*rule iter_narrow_step_ivl_termination*)
apply(*rule conjunct2*)
apply(*rule iter_widen_inv*[**where** $f =$ *step'* $\top$ **and** $P = \%C.$ $c =$ *strip* C & *top_on_acom* C $(-$ *vars* $C)$]])
apply(*auto simp*: *top_on_acom_widen top_on_step'*[*simplified comp_def vars_acom_def*]
  iter_widen_pfp top_on_bot vars_acom_def)
done

14.14.5 Counterexamples

Widening is increasing by assumption, but $x \le f\ x$ is not an invariant of widening. It can already be lost after the first step:

lemma assumes $!!x\ y::'a::wn.\ x \le y \Longrightarrow f\ x \le f\ y$
and $x \le f\ x$ **and** $\neg\ f\ x \le x$ **shows** $x \nabla f\ x \le f(x \nabla f\ x)$
nitpick[*card* = *3*, *expect* = *genuine*, *show_consts*, *timeout* = *120*]

oops

Widening terminates but may converge more slowly than Kleene iteration. In the following model, Kleene iteration goes from 0 to the least pfp in one step but widening takes 2 steps to reach a strictly larger pfp:

lemma assumes $!!x\ y::'a::wn.\ x \le y \Longrightarrow f\ x \le f\ y$
and $x \le f\ x$ **and** $\neg\ f\ x \le x$ **and** $f(f\ x) \le f\ x$
shows $f(x \nabla f\ x) \le x \nabla f\ x$
nitpick[*card* = *4*, *expect* = *genuine*, *show_consts*, *timeout* = *120*]

oops

end

15 Extensions and Variations of IMP

theory *Procs* **imports** *BExp* **begin**

15.1 Procedures and Local Variables

```
type_synonym pname = string

datatype
  com = SKIP
      | Assign vname aexp       (_ ::= _ [1000, 61] 61)
      | Seq    com  com         (_;;/ _  [60, 61] 60)
      | If     bexp com com     ((IF _/ THEN _/ ELSE _)  [0, 0, 61] 61)
      | While  bexp com         ((WHILE _/ DO _)  [0, 61] 61)
      | Var    vname com        ((1{VAR _;/ _}))
      | Proc   pname com com    ((1{PROC _ = _;/ _}))
      | CALL   pname

definition test_com =
{VAR ''x'';
 {PROC ''p'' = ''x'' ::= N 1;
  {PROC ''q'' = CALL ''p'';
   {VAR ''x'';
    ''x'' ::= N 2;;
    {PROC ''p'' = ''x'' ::= N 3;
     CALL ''q'';; ''y'' ::= V ''x''}}}}}

end
theory Procs_Dyn_Vars_Dyn imports Procs
begin
```

15.1.1 Dynamic Scoping of Procedures and Variables

```
type_synonym penv = pname ⇒ com

inductive
  big_step :: penv ⇒ com × state ⇒ state ⇒ bool (_ ⊢ _ ⇒ _ [60,0,60]
55)
where
Skip:    pe ⊢ (SKIP,s) ⇒ s |
Assign:  pe ⊢ (x ::= a,s) ⇒ s(x := aval a s) |
Seq:     ⟦ pe ⊢ (c₁,s₁) ⇒ s₂;  pe ⊢ (c₂,s₂) ⇒ s₃ ⟧ ⟹
         pe ⊢ (c₁;;c₂, s₁) ⇒ s₃ |

IfTrue:  ⟦ bval b s;  pe ⊢ (c₁,s) ⇒ t ⟧ ⟹
         pe ⊢ (IF b THEN c₁ ELSE c₂, s) ⇒ t |
IfFalse: ⟦ ¬bval b s;  pe ⊢ (c₂,s) ⇒ t ⟧ ⟹
         pe ⊢ (IF b THEN c₁ ELSE c₂, s) ⇒ t |
```

WhileFalse: ¬*bval b s* ⟹ *pe* ⊢ (*WHILE b DO c*,*s*) ⇒ *s* |
WhileTrue:
⟦ *bval b* s_1; *pe* ⊢ (*c*,s_1) ⇒ s_2; *pe* ⊢ (*WHILE b DO c*, s_2) ⇒ s_3 ⟧ ⟹
pe ⊢ (*WHILE b DO c*, s_1) ⇒ s_3 |

Var: *pe* ⊢ (*c*,*s*) ⇒ *t* ⟹ *pe* ⊢ ({*VAR x*; *c*}, *s*) ⇒ *t*(*x* := *s x*) |

Call: *pe* ⊢ (*pe p*, *s*) ⇒ *t* ⟹ *pe* ⊢ (*CALL p*, *s*) ⇒ *t* |

Proc: *pe*(*p* := *cp*) ⊢ (*c*,*s*) ⇒ *t* ⟹ *pe* ⊢ ({*PROC p* = *cp*; *c*}, *s*) ⇒ *t*

code_pred *big_step* **.**

values {*map t* [''x'',''y''] |*t*. (λ*p*. *SKIP*) ⊢ (*test_com*, <>) ⇒ *t*}

end
theory *Procs_Stat_Vars_Dyn* **imports** *Procs*
begin

15.1.2 Static Scoping of Procedures, Dynamic of Variables

type_synonym *penv* = (*pname* × *com*) *list*

inductive
big_step :: *penv* ⇒ *com* × *state* ⇒ *state* ⇒ *bool* (_ ⊢ _ ⇒ _ [*60*,*0*,*60*] *55*)
where
Skip: *pe* ⊢ (*SKIP*,*s*) ⇒ *s* |
Assign: *pe* ⊢ (*x* ::= *a*,*s*) ⇒ *s*(*x* := *aval a s*) |
Seq: ⟦ *pe* ⊢ (c_1,s_1) ⇒ s_2; *pe* ⊢ (c_2,s_2) ⇒ s_3 ⟧ ⟹
pe ⊢ (c_1;;c_2, s_1) ⇒ s_3 |

IfTrue: ⟦ *bval b s*; *pe* ⊢ (c_1,*s*) ⇒ *t* ⟧ ⟹
pe ⊢ (*IF b THEN* c_1 *ELSE* c_2, *s*) ⇒ *t* |
IfFalse: ⟦ ¬*bval b s*; *pe* ⊢ (c_2,*s*) ⇒ *t* ⟧ ⟹
pe ⊢ (*IF b THEN* c_1 *ELSE* c_2, *s*) ⇒ *t* |

WhileFalse: ¬*bval b s* ⟹ *pe* ⊢ (*WHILE b DO c*,*s*) ⇒ *s* |
WhileTrue:
⟦ *bval b* s_1; *pe* ⊢ (*c*,s_1) ⇒ s_2; *pe* ⊢ (*WHILE b DO c*, s_2) ⇒ s_3 ⟧ ⟹
pe ⊢ (*WHILE b DO c*, s_1) ⇒ s_3 |

Var: *pe* ⊢ (*c*,*s*) ⇒ *t* ⟹ *pe* ⊢ ({*VAR x*; *c*}, *s*) ⇒ *t*(*x* := *s x*) |

Call1: (p,c)#pe ⊢ (c, s) ⇒ t ⟹ (p,c)#pe ⊢ (CALL p, s) ⇒ t |
Call2: ⟦ p′ ≠ p; pe ⊢ (CALL p, s) ⇒ t ⟧ ⟹
    (p′,c)#pe ⊢ (CALL p, s) ⇒ t |

Proc: (p,cp)#pe ⊢ (c,s) ⇒ t ⟹ pe ⊢ ({PROC p = cp; c}, s) ⇒ t

code_pred *big_step* **.**

values {map t [''x'', ''y''] |t. [] ⊢ (test_com, <>) ⇒ t}

end
theory *Procs_Stat_Vars_Stat* **imports** *Procs*
begin

15.1.3 Static Scoping of Procedures and Variables

type_synonym *addr = nat*
type_synonym *venv = vname ⇒ addr*
type_synonym *store = addr ⇒ val*
type_synonym *penv = (pname × com × venv) list*

fun *venv :: penv × venv × nat ⇒ venv* **where**
venv(_,ve,_) = ve

inductive
  big_step :: penv × venv × nat ⇒ com × store ⇒ store ⇒ bool
  (_ ⊢ _ ⇒ _ [60,0,60] 55)
where
Skip: e ⊢ (SKIP,s) ⇒ s |
Assign: (pe,ve,f) ⊢ (x ::= a,s) ⇒ s(ve x := aval a (s o ve)) |
Seq: ⟦ e ⊢ (c_1,s_1) ⇒ s_2; e ⊢ (c_2,s_2) ⇒ s_3 ⟧ ⟹
    e ⊢ (c_1;;c_2, s_1) ⇒ s_3 |

IfTrue: ⟦ bval b (s ∘ venv e); e ⊢ (c_1,s) ⇒ t ⟧ ⟹
    e ⊢ (IF b THEN c_1 ELSE c_2, s) ⇒ t |
IfFalse: ⟦ ¬bval b (s ∘ venv e); e ⊢ (c_2,s) ⇒ t ⟧ ⟹
    e ⊢ (IF b THEN c_1 ELSE c_2, s) ⇒ t |

WhileFalse: ¬bval b (s ∘ venv e) ⟹ e ⊢ (WHILE b DO c,s) ⇒ s |
WhileTrue:
  ⟦ bval b (s_1 ∘ venv e); e ⊢ (c,s_1) ⇒ s_2;
    e ⊢ (WHILE b DO c, s_2) ⇒ s_3 ⟧ ⟹
  e ⊢ (WHILE b DO c, s_1) ⇒ s_3 |

```
Var: (pe,ve(x:=f),f+1) ⊢ (c,s) ⇒ t  ⟹
      (pe,ve,f) ⊢ ({VAR x; c}, s) ⇒ t |

Call1: ((p,c,ve)#pe,ve,f) ⊢ (c, s) ⇒ t  ⟹
        ((p,c,ve)#pe,ve′,f) ⊢ (CALL p, s) ⇒ t |
Call2: ⟦ p′ ≠ p;  (pe,ve,f) ⊢ (CALL p, s) ⇒ t ⟧ ⟹
       ((p′,c,ve′)#pe,ve,f) ⊢ (CALL p, s) ⇒ t |

Proc: ((p,cp,ve)#pe,ve,f) ⊢ (c,s) ⇒ t
      ⟹  (pe,ve,f) ⊢ ({PROC p = cp; c}, s) ⇒ t

code_pred big_step .


values {map t [10,11] |t.
  ([], <''x'' := 10, ''y'' := 11>, 12)
  ⊢ (test_com, <>) ⇒ t}

end
theory C_like imports Main begin
```

15.2 A C-like Language

```
type_synonym state = nat ⇒ nat

datatype aexp = N nat | Deref aexp (!) | Plus aexp aexp

fun aval :: aexp ⇒ state ⇒ nat where
aval (N n) s = n |
aval (!a) s = s(aval a s) |
aval (Plus a₁ a₂) s = aval a₁ s + aval a₂ s

datatype bexp = Bc bool | Not bexp | And bexp bexp | Less aexp aexp

primrec bval :: bexp ⇒ state ⇒ bool where
bval (Bc v) _ = v |
bval (Not b) s = (¬ bval b s) |
bval (And b₁ b₂) s = (if bval b₁ s then bval b₂ s else False) |
bval (Less a₁ a₂) s = (aval a₁ s < aval a₂ s)


datatype
  com = SKIP
```

```
    | Assign aexp aexp        (_ ::= _ [61, 61] 61)
    | New    aexp aexp
    | Seq    com  com         (_;/ _  [60, 61] 60)
    | If     bexp com com     ((IF _/ THEN _/ ELSE _)  [0, 0, 61] 61)
    | While  bexp com         ((WHILE _/ DO _)  [0, 61] 61)
```

inductive

```
  big_step :: com × state × nat ⇒ state × nat ⇒ bool  (infix ⇒ 55)
where
Skip:    (SKIP,sn) ⇒ sn |
Assign:  (lhs ::= a,s,n) ⇒ (s(aval lhs s := aval a s),n) |
New:     (New lhs a,s,n) ⇒ (s(aval lhs s := n), n+aval a s)  |
Seq:   ⟦ (c1,sn1) ⇒ sn2;  (c2,sn2) ⇒ sn3 ⟧ ⟹
         (c1;c2, sn1) ⇒ sn3 |

IfTrue:  ⟦ bval b s;  (c1,s,n) ⇒ tn ⟧ ⟹
         (IF b THEN c1 ELSE c2, s,n) ⇒ tn |
IfFalse: ⟦ ¬bval b s;  (c2,s,n) ⇒ tn ⟧ ⟹
         (IF b THEN c1 ELSE c2, s,n) ⇒ tn |

WhileFalse: ¬bval b s ⟹ (WHILE b DO c,s,n) ⇒ (s,n) |
WhileTrue:
  ⟦ bval b s1;  (c,s1,n) ⇒ sn2;  (WHILE b DO c, sn2) ⇒ sn3 ⟧ ⟹
   (WHILE b DO c, s1,n) ⇒ sn3
```

code_pred *big_step* **.**

declare [[*values_timeout = 3600*]]

Examples:

definition

```
array_sum =
 WHILE Less (!(N 0)) (Plus (!(N 1)) (N 1))
 DO ( N 2 ::= Plus (!(N 2)) (!(!(N 0)));
      N 0 ::= Plus (!(N 0)) (N 1) )
```

To show the first n variables in a *nat* ⇒ *nat* state:

definition

```
  list t n = map t [0 ..< n]
```

values {*list t n* |*t n*. (*array_sum*, *nth*[*3*,*4*,*0*,*3*,*7*],*5*) ⇒ (*t*,*n*)}

definition

linked_list_sum =

WHILE Less (*N 0*) (!(*N 0*))
DO (*N 1* ::= *Plus*(!(*N 1*)) (!(!(*N 0*)));
N 0 ::= !(*Plus*(!(*N 0*))(*N 1*)))

values {*list t n* |*t n*. (*linked_list_sum*, *nth*[*4*,*0*,*3*,*0*,*7*,*2*],*6*) ⇒ (*t*,*n*)}

definition
array_init =
New (*N 0*) (!(*N 1*)); *N 2* ::= !(*N 0*);
WHILE Less (!(*N 2*)) (*Plus* (!(*N 0*)) (!(*N 1*)))
DO (!(*N 2*) ::= !(*N 2*);
N 2 ::= *Plus* (!(*N 2*)) (*N 1*))

values {*list t n* |*t n*. (*array_init*, *nth*[*5*,*2*,*7*],*3*) ⇒ (*t*,*n*)}

definition
linked_list_init =
WHILE Less (!(*N 1*)) (!(*N 0*))
DO (*New* (*N 3*) (*N 2*);
N 1 ::= *Plus* (!(*N 1*)) (*N 1*);
!(*N 3*) ::= !(*N 1*);
Plus (!(*N 3*)) (*N 1*) ::= !(*N 2*);
N 2 ::= !(*N 3*))

values {*list t n* |*t n*. (*linked_list_init*, *nth*[*2*,*0*,*0*,*0*],*4*) ⇒ (*t*,*n*)}

end
theory *OO* **imports** *Main* **begin**

15.3 Towards an OO Language: A Language of Records

abbreviation *fun_upd2* :: (′*a* ⇒ ′*b* ⇒ ′*c*) ⇒ ′*a* ⇒ ′*b* ⇒ ′*c* ⇒ ′*a* ⇒ ′*b* ⇒ ′*c*
(_/′((*2*_,_ :=/ _)′) [*1000*,*0*,*0*,*0*] *900*)
where *f*(*x*,*y* := *z*) == *f*(*x* := (*f x*)(*y* := *z*))

type_synonym *addr* = *nat*
datatype *ref* = *null* | *Ref addr*

type_synonym *obj* = *string* ⇒ *ref*
type_synonym *venv* = *string* ⇒ *ref*
type_synonym *store* = *addr* ⇒ *obj*

datatype *exp* =

```
  Null |
  New |
  V string |
  Faccess exp string        (_•/_ [63,1000] 63) |
  Vassign string exp        ((_ ::=/ _) [1000,61] 62) |
  Fassign exp string exp    ((_•_ ::=/ _) [63,0,62] 62) |
  Mcall exp string exp      ((_•/_<_>) [63,0,0] 63) |
  Seq exp exp               (_;/ _ [61,60] 60) |
  If bexp exp exp           (IF _/ THEN (2_)/ ELSE (2_) [0,0,61] 61)
and bexp = B bool | Not bexp | And bexp bexp | Eq exp exp

type_synonym menv = string ⇒ exp
type_synonym config = venv × store × addr

inductive
  big_step :: menv ⇒ exp × config ⇒ ref × config ⇒ bool
    ((_ ⊢/ (_/ ⇒ _)) [60,0,60] 55) and
  bval ::  menv ⇒ bexp × config ⇒ bool × config ⇒ bool
    (_ ⊢ _ → _ [60,0,60] 55)
where
Null:
me ⊢ (Null,c) ⇒ (null,c) |
New:
me ⊢ (New,ve,s,n) ⇒ (Ref n,ve,s(n := (λf. null)),n+1) |
Vaccess:
me ⊢ (V x,ve,sn) ⇒ (ve x,ve,sn) |
Faccess:
me ⊢ (e,c) ⇒ (Ref a,ve',s',n') ⟹
 me ⊢ (e•f,c) ⇒ (s' a f,ve',s',n') |
Vassign:
me ⊢ (e,c) ⇒ (r,ve',sn') ⟹
 me ⊢ (x ::= e,c) ⇒ (r,ve'(x:=r),sn') |
Fassign:
⟦ me ⊢ (oe,c1) ⇒ (Ref a,c2);  me ⊢ (e,c2) ⇒ (r,ve3,s3,n3) ⟧ ⟹
 me ⊢ (oe•f ::= e,c1) ⇒ (r,ve3,s3(a,f := r),n3) |
Mcall:
⟦ me ⊢ (oe,c1) ⇒ (or,c2);  me ⊢ (pe,c2) ⇒ (pr,ve3,sn3);
   ve = (λx. null)(''this'' := or, ''param'' := pr);
   me ⊢ (me m,ve,sn3) ⇒ (r,ve',sn4) ⟧
  ⟹
 me ⊢ (oe•m<pe>,c1) ⇒ (r,ve3,sn4) for or |
Seq:
⟦ me ⊢ (e1,c1) ⇒ (r,c2);  me ⊢ (e2,c2) ⇒ c3 ⟧ ⟹
 me ⊢ (e1; e2,c1) ⇒ c3 |
```

IfTrue:
$\llbracket\ me \vdash (b,c_1) \rightarrow (True,c_2);\ \ me \vdash (e_1,c_2) \Rightarrow c_3\ \rrbracket \Longrightarrow$
$me \vdash (IF\ b\ THEN\ e_1\ ELSE\ e_2,c_1) \Rightarrow c_3$ |
IfFalse:
$\llbracket\ me \vdash (b,c_1) \rightarrow (False,c_2);\ \ me \vdash (e_2,c_2) \Rightarrow c_3\ \rrbracket \Longrightarrow$
$me \vdash (IF\ b\ THEN\ e_1\ ELSE\ e_2,c_1) \Rightarrow c_3$ |

$me \vdash (B\ bv,c) \rightarrow (bv,c)$ |

$me \vdash (b,c_1) \rightarrow (bv,c_2) \Longrightarrow me \vdash (Not\ b,c_1) \rightarrow (\neg bv,c_2)$ |

$\llbracket\ me \vdash (b_1,c_1) \rightarrow (bv_1,c_2);\ \ me \vdash (b_2,c_2) \rightarrow (bv_2,c_3)\ \rrbracket \Longrightarrow$
$me \vdash (And\ b_1\ b_2,c_1) \rightarrow (bv_1 \wedge bv_2,c_3)$ |

$\llbracket\ me \vdash (e_1,c_1) \Rightarrow (r_1,c_2);\ \ me \vdash (e_2,c_2) \Rightarrow (r_2,c_3)\ \rrbracket \Longrightarrow$
$me \vdash (Eq\ e_1\ e_2,c_1) \rightarrow (r_1{=}r_2,c_3)$

code_pred (*modes*: *i => i => o => bool*) *big_step* **.**

Example: natural numbers encoded as objects with a predecessor field. Null is zero. Method succ adds an object in front, method add adds as many objects in front as the parameter specifies.

First, the method bodies:

definition
m_succ = (''s'' ::= New)•''pred'' ::= V ''this''; V ''s''

definition *m_add =*
IF Eq (V ''param'') Null
THEN V ''this''
ELSE V ''this''•''succ''<Null>•''add''<V ''param''•''pred''>

The method environment:

definition
menv = (λm. Null)(''succ'' := m_succ, ''add'' := m_add)

The main code, adding 1 and 2:

definition *main =*
''1'' ::= Null•''succ''<Null>;
''2'' ::= V ''1''•''succ''<Null>;
V ''2'' • ''add'' <V ''1''>

Execution of semantics. The final variable environment and store are converted into lists of references based on given lists of variable and field names to extract.

values
$\{(r,\ map\ ve'\ [''1'',''2''],\ map\ (\lambda n.\ map\ (s'\ n)[''pred''])\ [0..<n])|$
$\quad r\ ve'\ s'\ n.\ menv \vdash (main,\ \lambda x.\ null,\ nth[],\ 0) \Rightarrow (r,ve',s',n)\}$

end

References

[1] T. Nipkow. Winskel is (almost) right: Towards a mechanized semantics textbook. In V. Chandru and V. Vinay, editors, *Foundations of Software Technology and Theoretical Computer Science*, volume 1180 of *Lect. Notes in Comp. Sci.*, pages 180–192. Springer-Verlag, 1996.

[2] T. Nipkow and G. Klein. *Concrete Semantics with Isabelle/HOL*. Springer, 2014. http://concrete-semantics.org.